MICROCOMPUTER RAYTHEON COMPONENTS

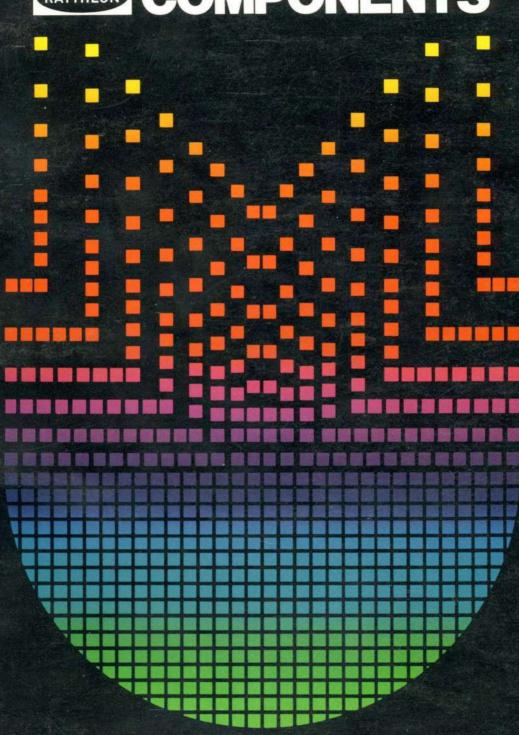


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INTRODUCTION

THE THIRD REVOLUTION

Over the past three years, the semiconductor industry has experienced the third revolutionary period of its relatively short history. The first of these occurred in the fifties with the introduction of mass-produceable transistors, which not only revolutionized the computer and communications industries, but also spawned a number of new companies and expanded the horizons of a number of older ones who elected to participate in this infant technology.

The second revolution occurred in the sixties when the first integrated circuits appeared. These first I.C.'s, featuring one or two gates per chip and selling for \$20 to \$30, created an even greater impact than the transistor itself. From the users' standpoint, I.C.'s allowed an order-of-magnitude increase in the complexity of the equipment they were designing (or an order-of-magnitude decrease in size, depending on the nature of that equipment) and, simultaneously, an order-of-magnitude improvement in cost-effectiveness. Coupled with this was a change in the engineering disciplines required by the user, since circuit designers, as such, were being replaced by logic designers and systems architects.

The second revolution caused a similar change in the semiconductor industry since it was now hiring the circuit designers who were leaving the end-equipment field. In addition, manufacturing, reliability and marketing skills began to expand, converting the industry from one dominated by the semi-conductor "black-art" specialists to one somewhat more balanced in its approach to business. Those companies that recognized the revolution (including many new ones) survived and prospered while those that didn't dropped out.

Similar changes are occurring today during the third revolution—this one created by the development of microprocessors. Once again order-of-magnitude improvements are being enjoyed by those users who have recognized the significance of microcomputer technology, and once again the engineering disciplines employed by the end user are changing, the emphasis now being placed on programming and other software development activities.

From the semiconductor manufacturer's standpoint, the changes are equally significant. Logic design and system architecture are becoming necessary skills for those who hope to succeed in the seventies. Marketing and sales will also be upgraded as semiconductor manufacturers begin selling "systems" rather than "components." Applications, training, software and hardware development tools and development systems are all new demands placed on the semiconductor manufacturer who hopes to survive the third revolution.

As with the first two revolutions, the inception of the microprocessor was characterized by the introduction of many different device types from many different manufacturers using many different technologies. Some of these will become high volume, multi-sourced, industry standard products (like TTL) and others will survive at lower volumes because they have certain performance advantages

which the "standard" devices cannot match in selected applications (similar to the role played by CMOS and ECL).

Most, however, will fall by the wayside as manufacturing efficiency and competitive pressures force the price to a point where the non-standard types become unprofitable. This point is rapidly approaching in the microprocessor area. The selection of the right microprocessor, therefore, is one of the current challenges facing today's system designer.

There are, today, basically two markets for microprocessors and two basic approaches being employed by semiconductor manufacturers to satisfy those markets. The two markets can best be defined as the *logic replacement market* and the *performance market*. The logic replacement market is best characterized by a certain cost-performance inelasticity which states that once a certain minimal performance level is achieved, no additional cost (and presumably selling price) can be justified in terms of additional units sold. Thus, the emphasis in the logic replacement market is either cost reduction at a given performance level or increased performance at no increase in cost. (This market, for example, would use TTL logic in applications where TTL performance was not required, but where TTL was the lowest cost solution.)

In the performance market, cost-performance is relatively elastic; i.e. additional end-equipment performance can demand a higher selling price and still provide the manufacturer with added volume and/or revenue. Here, the tendency would be to use the highest performance technology available in an effort to gain a performance edge over competition.

The two basic technologies which today compete for the total microprocessor market are Metal-Oxide-Semiconductor (MOS) and bipolar. While some overlap does exist, the market generally has accepted MOS in the logic replacement segment and bipolar in the performance segment.

With the large number of MOS alternatives available, the system designer has no simple chore in determining the best solution for his application. Not only does he have a choice of 4, 8, 12, and 16 bit devices, but once settling on the size of his data word, he'll have up to a dozen different n-bit designs to choose from. As with most alternatives of this nature, any one of a number of different designs will meet his system requirements and the designer will be apt to choose one of the two or three types within each category that appear to be "industry standards."

The situation is not quite as confusing in the bipolar arena simply because there are not as many alternatives as with the MOS approach. This is partially explained by the fact that most bipolar designs are bit-slice oriented thus eliminating the need to fix the size of the data word prior to selecting a particular bit slice design.

Perhaps even more significant in explaining the relatively small number of bipolar alternatives is the nature of the performance market which uses the technology. Since this



market tends to be smaller than the logic replacement market in terms of the number of users, the breadth of applications and the total amount of end-user equipment shipped, there is not the need for nor the incentive to create a large number of different device types. Thus, the semiconductor industry—at least to this point in time—has elected to concentrate their efforts on a relatively small number of device types.

One of these types (in fact the only one currently supported by more than two manufacturers) is the 2900 family, a bit slice approach introduced to the market in 1975. This state-of-the art family consists of a series of MSI/LSI devices designed to complement and enhance one another in typical microcomputer applications.

Recognizing the need for multi-sourcing in the market served by bipolar microprocessors, Raytheon, in 1975, became the first alternate source for 2900 family products and today has more 2900 family parts than any other semi-conductor manufacturer in the industry except for AMD who originated the series. In addition, Raytheon is committed to the continued support of the family, not only to the extent of providing all 2900 series devices, but also in the equally important areas of bipolar memory, logic, and the development tools required for both hardware and software design.

THE 2900 FAMILY

All microcomputers (in fact all practical digital computers) consist of four basic functions—central processing unit (CPU), control, memory, and input/output (I/O), (See Figure 1) Connecting the various functions are the address and data busses plus various control lines. While all microcomputers have essentially the same general architecture (i.e., perform the four functions shown in Figure 1), there are a number of differences in the way the various functions are implemented. The most obvious differences occur between MOS and bipolar technologies.

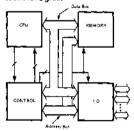


Figure 1, Typical Computer Functions.

In the MOS approach, for example, the most common implementation places the CPU and control function both on one chip. This chip—which was the original "microprocessor"—inputs instructions and data, performs the operation required by the instruction, and outputs the resulting data, It also provides the addressing function to locate the instructions in memory, and to read or write data in memory and/or i/O. Because the decoding circuitry for the instruc-

tions (the main portion of the control section) is hardwired on the chip, the same instruction (set of "1's" and "0's") will always cause the same operation (Add, Subtract, Compare, etc.) to be performed, and therefore, the instruction set is said to be "fixed". This is one of the major distinctions between MOS and bipolar microprocessors.

The microprocessor, however, is not a computer until one adds the memory and I/O circuits. In the case of MOS, this normally consists of ROM or PROM for program storage, RAM for data storage, and either discrete logic or special purpose LSI chips for I/O.*

*Some newer MOS devices include some memory and I/O capability or the same chip as the CPU and control. While somewhat limited in capability, these new chips can realistically be called microcomputers without the necessity of external memory or I/O.

With the more popular bipolar microprocessors, such as the 2900 family, the CPU and control function are implemented differently. The CPU, for example, rather than being a fixed number of data bits wide (4, 8, 12, or 16 as in MOS) is actually a 2 or 4 bit slice that can be cascaded to any width that is a multiple of the basic slice. A four bit slice such as the 2901, for example, can be used to implement 4, 8, 12, 16...... 32, etc. bit microcomputers.

There are some obvious advantages and disadvantages to a bit slice approach when compared to the fixed data word, MOS approach. The disadvantage is, of course, that more chips are required per CPU.

Not only does the CPU function usually require more chips in bipolar designs, but also the control function is not even included on the chip as it is with the MOS devices. There are, however, advantages to bit slices over fixed data word machines; one of the most obvious being the capability of building larger, more powerful machines. In MOS, there simply isn't anything available over 16 bits, while 24, 32, and even 48 bit bipolar microcomputers have been designed and manufactured. Since wider data paths usually mean higher throughput, the bipolar approach can usually provide higher performance simply by building a bigger machine.

The other major architectural difference between MOS and bipolar is in the implementation of the control function. As mentioned previously, the MOS controller shares the same chip with the CPU and features a fixed instruction decoder. The bit slice controller, on the other hand, is constructed using a number of devices including PROMs, registers, and sequencers. (See Figure 2)

As in the case of the MOS controller, the main function of the bipolar equivalent is to accept instructions from memory via the data bus, decode the instruction, and then provide the individual control signals to the CPU, memory, and I/O. These control signals are generated by combining the decoded instruction with a multi-phase clock to provide a properly timed sequence of control words to the various system elements.



As shown in Figure 2, (*) the control signals for the CPU, memory, and I/O are actually outputs from an array of PROMs (3). Each set of these control signals will cause a specific action to occur during a given clock period. Since it normally takes a series of these signals to actually accomplish what is normally considered to be a machine instruction (add, subtract, etc.), the individual sets are referred to as microinstructions. Thus, a machine instruction (or macroinstruction) consists of a series of microinstructions emanating from the PROM array.

*Figure 2 shows only the major system elements for clarity. Various register and other logic elements will be required in an actual design.

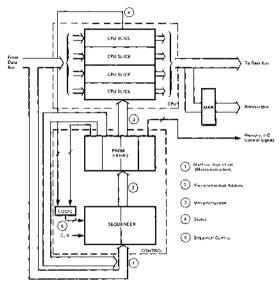


Figure 2. Typical Bit-Slice Implementation for CPU and Control.

As one might expect, the series of microinstructions to accomplish a common machine instruction would be stored in consecutive address locations in the PROMs. Thus to accomplish a machine instruction—say an ADD—one would simply set a counter with the address of the first microinstruction of the ADD sequence and then simply increment the counter as each microinstruction is completed. This would continue until the final microinstruction of the ADD series was executed at which time the control would be set to accept the next machine instruction.

The circuitry that provides the addressing and counting functions is called a sequencer and most bipolar microprocessor families provide special LSI chips to accomplish this function. For example, the 2900 family currently provides two such circuits—the 2909 and 2911—and a third is in development. As with the CPU chips, sequencers are normally provided as slices to allow for a variety of microprogram sizes.

In addition to the addressing and counting functions, most sequencers provide circuitry to accommodate skips, jumps, and branches in the microprogram. This is accomplished by combining CPU status information (4) with a portion of the microcode to exercize control over the source of the next microinstruction address.

As in any microcomputer implementation, the machine instruction is stored in main memory (usually PROM) and is presented to the control circuit via the data bus ①. In most bipolar architectures, the machine instruction is simply the starting address of the sequence of microinstructions which will be used to actually accomplish the machine instruction.

As with the bipolar CPU, there are advantages and disadvantages in the manner in which the control function is implemented. One obvious disadvantage is chip count. It normally requires in excess of 10 packages to construct a bipolar microprocessor controller whereas the MOS controller is provided as part of the CPU chip.

The other factor which is sometimes considered to be a disadvantage is the necessity of generating the microinstructions. While this does, in essence, require twice the programming effort compared with MOS microprocessor programming, it is, at the same time, one of the features which enhances the performance aspects of the bipolar approach. By designing his own instruction set—which is precisely what the microprogrammer is doing—the user can tailor the instructions to his unique application, which improves performance and minimizes macroprogram storage requirements.

The ability to microprogram his machine also allows the user to emulate virtually any computer architecture (including many of today's popular mini's) thus making it possible to utilize existing applications software which often is more costly to generate than the hardware.

SUMMARY

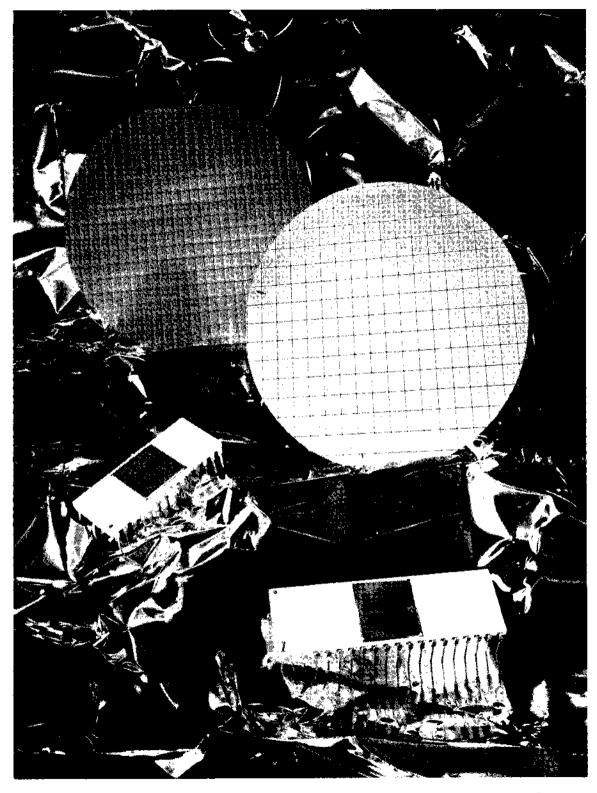
In the following pages of this catalog, a detailed description of each member of the 2900 family is presented in the form of a data sheet. Where additional information is required to fully explain the function of a given device, a users guide is also provided.

In addition, data sheet and applications information is provided for Raytheon's bipolar memory products (PROMs and RAMs) since they will form an integral part of any bipolar microcomputer design.

Finally, a simple system design is presented to further clarify the basic operation of the 2900 family and to give a typical example of the procedure used to design bipolar microcomputer systems.

If there are any questions regarding the material presented in this catalog or for any additional information or applications assistance, simply contact Raytheon's Application Staff by calling 415-968-9211 and asking for Microprocessor Applications.





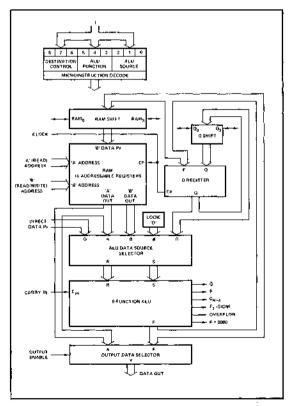


GENERAL DESCRIPTION

The four-bit bipolar microprocessor slice is designed as a high-speed cascadable element intended for use in CPU's, peripheral controllers, programmable microprocessors and numerous other applications. The microinstruction flexibility of the Am2901 will allow efficient emulation of almost any digital computing machine.

The device, as shown in the block diagram below, consists of a 16-word by 4-bit two-port RAM, a high-speed ALU, and the associated shifting, decoding and multiplexing circuitry. The nine-bit microinstruction word is organized into three groups of three bits each and selects the ALU source operands, the ALU function, and the ALU destination register. The microprocessor is cascadable with full look-ahead or with ripple carry, has three-state outputs, and provides various status flag outputs from the ALU. Advanced low-power Schottky processing is used to fabricate this 40-lead LSI chip.

MICROPROCESSOR SLICE BLOCK DIAGRAM



DISTINCTIVE CHARACTERISTICS

- Two-address architecture
 - Independent simultaneous access to two working registers saves machine cycles.
- Eight-function ALU
 - Performs addition, two subtraction operations, and five logic functions on two source operands.
- Flexible data source selection -
 - ALU data is selected from five source ports for a total of 203 source operand pairs for every ALU function.
- Left/right shift independent of ALU
 - Add and shift operations take only one cycle.
- Four status flags
 - Carry, overflow, zero, and negative.
- Expandable
 - Connect any number of Am2901's together for longer word lengths.
- Microprogrammable
 - Three groups of three bits each for source operand, ALU function, and destination control.

ORDERING INFORMATION

Package Type	Temperature Range	Order Number
Molded DIP	0°C to +70°C	AM2901PC
Hermetic DIP	0°C to +70°C	AM2901DC
Hermetic DIP	–55°C to +125°C	AM2901DM
Hermetic Flat Pack	–55°C to +125°C	AM2901FM
Dice	0°C to +70°C	AM2901XC



MAXIMUM RATINGS (Above which the useful life may be impaired)

Storage Temperature	-65°C to +150°C
Temperature (Ambient) Under Bias	-55°C to +125°C
Supply Voltage to Ground Potential	-0.5 V to +6.3 V
DC Voltage Applied to Outputs for HIGH Output State	-0.5 V to +V _{CC} max.
DC Input Voltage	-0.5 V to +5.5 V
DC Output Current, Into Outputs	30 mA
DC Input Current	-30 mA to +5.0 mA

OPERATING RANGE

P/N	Ambient Temperature	V _{CC}
Am2901PC, DC	0°C to +70°C	4.75 V to 5.25 V
Am2901DM, FM	-55°C to +125°C	4,50 V to 5,50 V

STANDARD SCREENING

(Conforms to MIL-STD-883 for Class C Parts)

	MIL-STD-883		-
Step	Method	Conditions	Am2901DM,FM
Pre-Seal Visual Inspection	2010	В	100%
Stabilization Bake	1008	C 24-hour 150°C	100%
Temperature Cycle	1010	C65°C to +150°C 10 cycles	100%
Centrifuge	2001	B 10,000G	100%
Fine Leak	1014	A 5 x 10 ⁻⁸ atm-cc/cm ³	100%
Gross Leak	1014	C2 Fluorocarbon	100%
Electrical Test	5004	See below for	100%
Subgroups 1 and 7	5004	definitions of subgroups	
Insert Additional Screening he	re for Class B Parts		
Group A Sample Tests	"" ""	I	
Subgroup 1			LTPD = 5
Subgroup 2	!		LTPD = 7
Subgroup 3	FROE	See below for	LTPD = 7
Subgroup 7	5005	definitions of subgroups	LTPD = 7
Subgroup 8			LTPD = 7
Subgroup 9	1		LTPO = 7

ADDITIONAL SCREENING FOR CLASS B PARTS

_	MIL-STD-883	0. 15.7	Level
Step	Method	Conditions	Am2901DMB, FMI
Burn- n	1015	D 125°C 160 hours min	100%
Electrical Test	5004		i
Subgroup 1		İ	100%
Subgroup 2			100%
Subgroup 3			100%
Subgroup 7		1	100%
Subgroup 9			100%

GROUP A SUBGROUPS (as defined in MIL-STD-883, method 5005)

Subgroup	Parameter	Temperature				
1	DC	25°C				
2	DC	Maximum rated temperature				
3	DC	Minimum rated temperature				
7	Function	25°C				
8	Function	Maximum and minimum rated temperature				
9	Switching	25°C				
10	Switching	Maximum Rated Temeperature				
11	Switching	Minimum Rated Temperature				



ELECTRICAL CHARACTERISTICS OVER OPERATING RANGE (Unless Otherwise Noted)

(Group A, Subgroups 1, 2 and 3)

rameters	Description	Test Cond	Test Conditions (Note 1)			Typ. (Note 2)	Max.	Units								
]		1 _{OH} =1 Y ₀ , Y ₁ , Y		2.4		j									
						 	}									
				.0mA, C _{n+4}	2,4	1	——-									
VOH	Output HIGH Voltage	V _{CC} = MIN. VIN = VIH or VIL		00µA, OVR, <u>Р</u>	2.4			Volts								
		114 11H 22 11E	10H * −6		2.4											
		I _{OH} = −6 RAM _{0, 3}	, C _{0, 3}	2.4												
			10H = ~1	.6mA, G	2.4											
ICEX	Output Leakage Current for F = 0 Output	V _{CC} = MIN., V _{OH} = V _{IN} = V _{IH} or V _{IL}	5.5V				250	μА								
			IOL * 16				0.5									
		V _{CC} = MIN.,		mA, C _{n+4} , F = 0			0.5									
VOL	Output LOW Voltage	VIN = VIH or VIL		mA, OVFI, P			0.5	Volts								
			I _{OL} = 6.0	lmA, F3			0.5									
V _{IH}	Input HIGH Level	Guaranteed input log voltage for all inputs			2.0			Volts								
۷۱۲	Input LOW Level	Guaranteed input log	ical LOW	Military	· · · · · - · -	 	0.7	Volts								
* L	input 2011 CEVE	voltage for all inputs		Commercial			8.0	VOITS								
v _t	Input Clamp Voltage	V _{CC} = M1N., 1 _{IN} = -	18mA				-1.5	Volts								
		Clock, O	Clock, OE			-0.36										
			Ag, A1,	A2, A3			-0.36	1								
		V _{CC} = MAX. V _{IN} = 0.5V	B ₀ , B ₁ , B ₂ , B ₃				-0.36	mA								
116	Input LOW Current		D ₀ , D ₁ , D ₂ , D ₃			<u> </u>	-0.72									
			10, 11, 12			<u> </u>	-0.36									
	}		13, 14, 15			ļ	-0.72									
				, Q _{0, 3} (Note 4)		ļ i	0.8									
			c _n				-3.6									
	1		Clock, OE				20									
			A ₀ , A ₁ ,	A ₂ , A ₃			20	μΑ								
		V _{CC} = MAX.	8 ₀ , B ₁ , 6	3 ₂ , B ₃			20									
Iн	Input HIGH Current		D ₀ , D ₁ ,	· ·			40									
	į l	V _{IN} = 2.7V	10, 11, 12				20									
			13, 14, 15				40									
				, Q _{0, 3} (Note 4)			100									
1.	Input HIGH Current	V _{CC} * MAX., V _{IN} =	C _n				1.0									
<u>'lı</u>	inpocition content	*CC - MXV-1 * IN =	1	Vo = 2 411			50	MA								
			Y ₀ , Y ₁ , Y ₂ , Y ₃			 										
tan:			- 2, 3	V _O = 0.5V			-50									
	Off State (High Impedance) Output Current	V _{CC} = MAX.	V _{CC} = MAX.	VCC = MAX.	VCC = MAX.	VCC = MAX.	V _{CC} = MAX.	VCC = MAX.	V _{CC} = MAX.	VCC = MAX.	RAMo, s				100	μА
			Ω _{0,3}	V _O = 0.5V (Note 4)			-800									
			Y ₀ , Y ₁ ,	Y ₂ , Y ₃ , \overline{G}	15		-40									
		V+0 - 6.7 5V	Cn+4		-15	I	-40									
IOS	Output Short Circuit Current (Note 3)	V _{CC} = 5.75V V _O = 0.5V	OVR, P		-15		-40	mA								
	1.3545 67	, U 4.04	F3		-15	L	-40									
	<u> </u>		RAM _{0, 3}	, Q _{0, 3}	-15		_ 4 0									
	,		Military			185	280									

Notes: 1. For conditions shown as MIN, or MAX,, use the appropriate value specified under Electrical Characteristics for the applicable device type.

These are three-state outputs internally connected to TTL inputs. Input characteristics are measured with 1678 in a state such that the
three-state output is OFF.



^{2.} Typical limits are at V_{CC} = 5.0V, 25°C ambient and maximum loading.

^{3.} Not more than one output should be shorted at a time. Duration of the short circuit test should not exceed one second.

TABLE I

Tables 1, 11, and 111 below define the timing requirements of the Am2901 in a system. The Am2901 is guaranteed to function correctly over the operating range when used within the delay and set-up time constraints of these tables for the appropriate device type. The tables are divided into three types of parameters; clock characteristics, combinational delays from inputs to outputs, and set-up and hold time requirements. The latter table defines the time prior to the end of the cycle (i.e., clock LOW-to-HIGH transition) that each input must be stable to guarantee that the correct data is written into one of the internal registers.

The performance of the Am2901 within the limits of these tables is guaranteed by the testing defined as "Group A, Subgroup 9" Electrical Testing. For a copy of the tests and limits used for subgroup 9, contact Product Marketing.

CYCLE TIME AND CLOCK CHARACTERISTICS

TIME	Am2901DC,PC	Am2901DM,FM
Read-Modify-Write Cycle (time from selection of A, B registers to end of cycle)	105 ns	120ns
Maximum Clock Frequency to Shift Q Register (50% duty cycle)	9.5MHz	8.3MHz
Minimum Clock LOW Time	30ns	30 ns
Minimum Clock HIGH Time	30ns	30 ns
Minimum Clock Period	105ns	120ns

TABLE 11 MAXIMUM COMBINATIONAL PROPAGATION DELAYS (all in ns, CL \leq 15 pF)

	1	Am2901DC, PC (0°C to +70°C; 5 V ±5%)							Am	29010	M, FM	(-55°	C to +	125°C;	5 V ±1	0%)
To Output	v			G, ₽	F≖0	OVD	Shift Outputs			F.		Ĝ, P	F=0		Shift Outputs	
From Input	Y	F3	C _{n+4}	G, F	470	OVR	RAM _O RAM ₃		Y	F3	Cn+4	G, P	RL≃ 470	i	RAM ₀ RAM ₃	
A, B	110	85	80	80	110	75	110	_	120	95	90	90	120	85	120	_
D (arithmetic mode)	100	70	70	70	100	60	95	_	110	80	75	75	110	65	105	<u> </u>
D (I = X37) (Note 5)	60	50		_	60		60		65	55	-		65	-	65	
Cn	55	35	30		50	40	55		60	40	30	-	55	45	60	-
l ₀₁₂	85	65	65	65	80	65	80	_	90	70	70	70	85	70	85	
1345	70	55	60	60	70	60	65	<u> </u>	75	60	65	65	75	65	70	
1678	55		-	-	_		45	45	60	_	_	_	_	_	50	50
OE Enable/Disable	40/25	_	_		_	-	-	_	40/25	_	_	_	-	. –	-	-
A bypassing ALU (I = 2xx).	60	-	_		_	-	_		65		-	-		_	_	-
Clock 🍜 (Note 6)	115	85	100	100	110	95	105	60	125	95	110	110	120	105	115	65

SET-UP AND HOLD TIMES (all in ns) (Note 1)

TABLE III

From Input	Notes	Am2901DC, PC (0°C	to +70°C, 5V ±5%)	Am2901DM,FM (-55°C to +125°C,5V, ±10%			
	IAD(62	Set-Up Time	Hold Time	Set-Up Time	Hold Time		
A, B Source	2, 4 105 3, 5 t _{pw} L + 30		0	120 t _{pw} L + 30	0		
B Dest.	2, 4	tpwL + 15	0	t _{pw} L +15	0		
D (arithmetic mode)		100	0	110	0		
D (I = X37) (Note 5)		60	0	65	0		
Cn		55	0	60	0		
I ₀₁₂		85	0	90	0		
¹ 345		70	0	75	0		
¹ 678	4	t _{pw} L + 15	0	t _{pw} L + 15	0		
RAM _{0, 3} , Q _{0, 3}		30	0	30	0		

Notes: 1. See Figure 11 and 12.

- 2. If the B address is used as a source operand, allow for the "A, B source" set-up time; If it is used only for the destination address, use the "B dest." set-up time.
- 3. Where two numbers are shown, both must be met.
- 4. "tpwL" is the clock LOW time.
- 5. DV0 is the fastest way to load the RAM from the D inputs. This function is obtained with I = 337.
- 6. Using Q register as source operand in arithmetic mode. Clock is not normally in critical speed path when Q is not a source.

SET-UP AND HOLD TIMES (minimum cycles from each input)

Set-up and hold times are defined relative to the clock LOW-to-HIGH edge. Inputs must be steady at all times from the set-up time prior to the clock until the hold time after the clock. The set-up times allow sufficient time to perform the correct operation on the correct data so that the correct ALU data can be written into one of the registers.

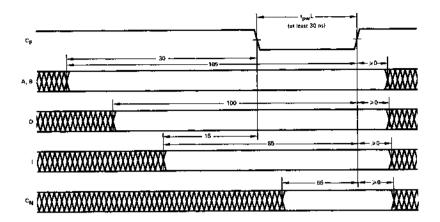
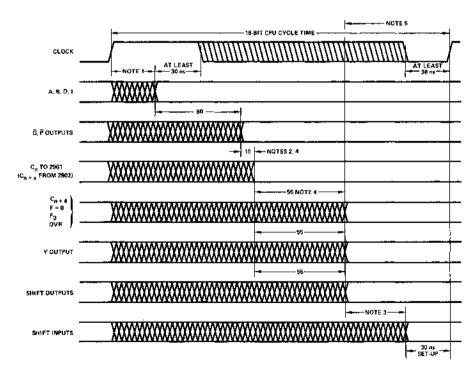


Figure 1. Minimum Cycle Times from Inputs. Numbers Shown are Minimum Data Stable Times for Am2901 DC, in ns. See Table III for Detailed Information.



- Notes: 1. This delay is the max, t_{pd} of the register containing A, B, D, and I. For the Am2918, use 13ns.
 - 2. 10ns for look-ahead carry. For ripple carry over 16 bits use $2 \times (C_n \rightarrow C_{n+4})$, or 60ns.
 - This is the delay associated with the multiplexer between the shift outputs and shift inputs on the Am2901s. See Figure 19.
 - 4. Not applicable for logic operations.
 - 5. Clock rising edge may occur here if add and shift do not occur on same cycle.

Figure 2. Switching Waveforms for 16-Bit System Assuming A, B, D and I are all Driven from Registers with the same Propagation Delay, Clocked by the Am2901 Clock.



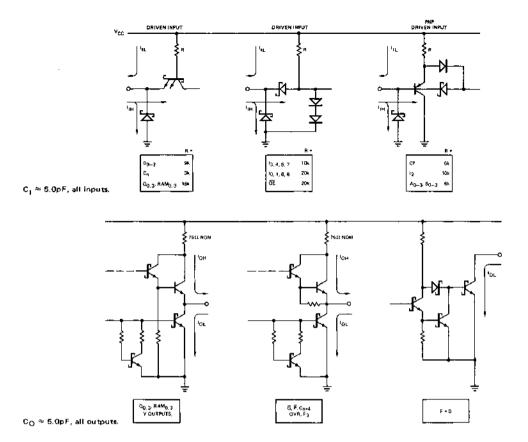


Figure 3. Am2901 Input/Output Current Interface Conditions.

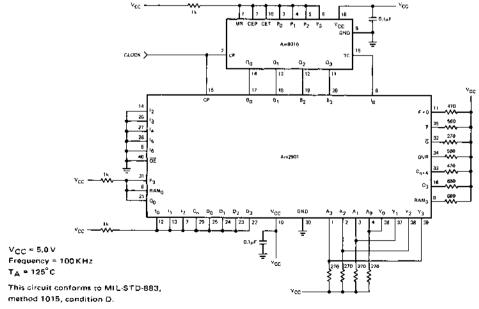


Figure 4. Am2901 Burn-In Circuit.



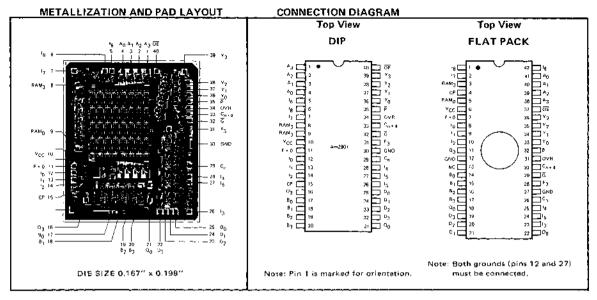


Figure 5.

DEFINITIONS

- A₀₋₃ The four address inputs to the register stack used to select one register whose contents are displayed through the A-port.
- B₀₋₃ The four address inputs to the register stack used to select one register whose contents are displayed through the B-port and into which new data can be written when the clock goes LOW.
- I₀₋₈ The nine instruction control lines to the Am2901, used to determine what data sources will be applied to the ALU (I₀₁₂), what function the ALU will perform (I₃₄₅), and what data is to be deposited in the Q-register or the register stack (I₆₇₈).
- Q₃ A shift line at the MSB of the Q register (Q₃) and the register stack (RAM₃). Electrically these lines are three-state outputs connected to TTL inputs internal to the Am2901. When the destination code on I₆₇₈ indicates an up shift (octal 6 or 7) the three-state outputs are enabled and the MSB of the Q register is available on the Q₃ pin and the MSB of the ALU output is available on the RAM₃ pin. Otherwise, the three-state outputs are OFF (high-impedance) and the pins are electrically LS-TTL inputs. When the destination code calls for a down shift, the pins are used as the data inputs to the MSB of the Q register (octal 4) and RAM (octal 4 or 5).
- Q₀ Shift lines like Q₃ and RAM₃, but at the LSB of the Q-register and RAM. These pins are tied to the Q₃ and RAM₃ pins of the adjacent device to transfer data between devices for up and down shifts of the Q-register and ALU data.
- D₀₋₃ Direct data inputs. A four-bit data field which may be selected as one of the ALU data sources for entering data into the Am2901. D₀ is the LSB.

- Y₀₋₃ The four data outputs of the Am2901. These are three-state output lines. When enabled, they display either the four outputs of the ALU or the data on the A-port of the register stack, as determined by the destination code I₆₇₈.
- Output Enable. When OE is HIGH, the Y outputs are OFF; when OE is LOW, the Y outputs are active (HIGH or LOW).
- P, G The carry generate and propagate outputs of the Am2901's ALU. These signals are used with the Am2902 for carry-lookahead. See Figure 8 for the logic equations.
- OVR Overflow. This pin is logically the Exclusive-OR of the carry-in and carry-out of the MSB of the AŁU. At the most significant end of the word, this pin indicates that the result of an arithmetic two's complement operation has overflowed into the sign-bit. See Figure 8 for logic equation.
- F=0 This is an open collector output which goes HIGH (OFF) if the data on the four ALU outputs F_{0-3} are all LOW. In positive logic, it indicates the result of an ALU operation is zero.
- C_n The carry-in to the Am2901's ALU.
- C_{n+4} The carry-out of the Am2901's ALU. See Figure 8 for equations.
- CP The clock to the Am2901. The Q register and register stack outputs change on the clock LOW-to-HIGH transition. The clock LOW time is internally the write enable to the 16 x 4 RAM which comprises the "master" latches of the register stack. While the clock is LOW, the "slave" latches on the RAM outputs are closed, storing the data previously on the RAM outputs. This allows synchronous master-slave operation of the register stack.

A detailed block diagram of the bipolar microprogrammable microprocessor structure is shown in Figure 1. The circuit is a four-bit slice cascadable to any number of bits. Therefore, all data paths within the circuit are four bits wide. The two key elements in the Figure 1 block diagram are the 16-word by 4-bit 2-port RAM and the high-speed ALU.

Data in any of the 16 words of the Random Access Memory (RAM) can be read from the A-port of the RAM as controlled by the 4-bit A address field input. Likewise, data in any of the 16 words of the RAM as defined by the B address field input can be simultaneously read from the B-port of the RAM. The same code can be applied to the A select field and B select field in which case the identical file data will appear at both the RAM A-port and B-port outputs simultaneously.

When enabled by the RAM write enable (RAM EN), new data is always written into the file (word) defined by the B address field of the RAM. The RAM data input field is driven by a 3-input multiplexer. This configuration is used to shift the ALU output data (F) if desired. This three-input multiplexer scheme allows the data to be shifted up one bit position, shifted down one bit position, or not shifted in either direction.

The RAM A-port data outputs and RAM B-port data outputs drive separate 4-bit latches. These latches hold the RAM data while the clock input is LOW. This eliminates any possible race conditions that could occur while new data is being written into the RAM.

The high-speed Arithmetic Logic Unit (ALU) can perform three binary arithmetic and five logic operations on the two 4-bit input words R and S. The R input field is driven from a 2-input multiplexer, while the S input field is driven from a 3-input multiplexer. Both multiplexers also have an inhibit capability; that is, no data is passed. This is equivalent to a "zero" source operand.

Referring to Figure 6, the ALU R-input multiplexer has the RAM A-port and the direct data inputs (D) connected as inputs. Likewise, the ALU S-input multiplexer has the RAM A-port, the RAM B-port and the Q register connected as inputs.

This multiplexer scheme gives the capability of selecting various pairs of the A, B, D, Q and "O" inputs as source operands to the ALU. These five inputs, when taken two at a time, result in ten possible combinations of source operand pairs. These combinations include AB, AD, AQ, AD, BD, BQ, BQ, DQ, DQ and QQ. It is apparent that AD, AQ and AQ are somewhat redundant with BD, BQ and BQ in that if the A address and B address are the same, the identical function results. Thus, there are only seven completely non-redundant source operand pairs for the ALU. The Am2901 microprocessor implements eight of these pairs. The microinstruction inputs used to select the ALU source operands are the IQ, I1, and I2 inputs. The definition of IQ, I1, and I2 for the eight source operand combinations are as shown in Figure 7. Also shown is the octal code for each selection.

The two source operands not fully described as yet are the D input and Q input. The D input is the four-bit wide direct data field input. This port is used to insert all data into the working registers inside the device. Likewise, this input can be used in the ALU to modify any of the internal data files. The Q register is a separate 4-bit file intended primarily for multiplication and division routines but it can also be used as an accumulator or holding register for some applications.

The ALU itself is a high-speed arithmetic/logic operator capable of performing three binary arithmetic and five logic functions. The I₃, I₄, and I₅ microinstruction inputs are used to select the

ALU function. The definition of these inputs is shown in Figure 8. The local code is also shown for reference. The normal technique for cascading the ALU of several devices is in a look-ahead carry mode. Carry generate, \vec{F} , and carry propagate, \vec{F} , are outputs of the device for use with a carry-look-ahead-generator such as the Am2902 ('182). A carry-out, C_{n+4} , is also generated and is available as an output for use as the carry flag in a status register. Both carry-in (C_n) and carry-out (C_{n+4}) are active HIGH.

The ALU has three other status-oriented outputs. These are F_3 , F=0, and overflow (OVR). The F_3 output is the most significant (sign) bit of the ALU and can be used to determine positive or negative results without enabling the three-state data outputs. F_3 is non-inverted with respect to the sign bit output Y_3 . The F=0 output is used for zero detect. It is an open-collector output and can be wire OR'ed between microprocessor stices. F=0 is HIGH when all F outputs are LOW. The overflow output (OVR) is used to flag arithmetic operations that exceed the available two's complement number range. The overflow output (OVR) is HIGH when overflow exists. That is, when C_{n+3} and C_{n+4} are not the same polarity.

The ALU data output is routed to several destinations. It can be a data output of the device and it can also be stored in the RAM or the Q register, Eight possible combinations of ALU destination functions are available as defined by the 16, 17, and 18 microinstruction inputs. These combinations are shown in Figure 9.

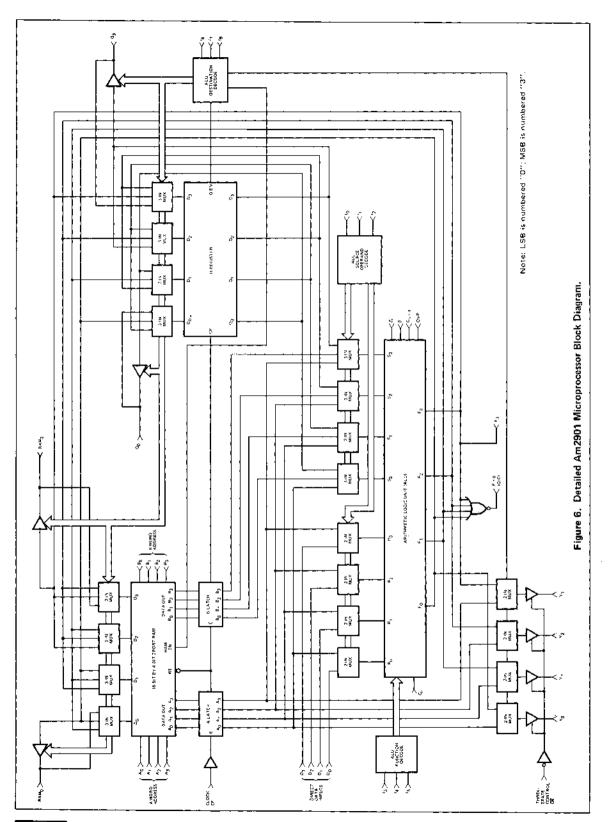
The four-bit data output field (Y) features three-state outputs and can be directly bus organized. An output control (OE) is used to enable the three-state outputs. When OE is HIGH, the Y outputs are in the high-impedance state.

A two-input multiplexer is also used at the data output such that either the A-port of the RAM or the AŁU outputs (F) are selected at the device Y outputs. This selection is controlled by the I₆, I₇, and I₈ microinstruction inputs. Refer to Figure 4 for the selected output for each microinstruction code combination.

As was discussed previously, the RAM inputs are driven from a three-input multiplexer. This allows the ALU outputs to be entered non-shifted, shifted up one position (X2) or shifted down one position ($\dot{\tau}2$). The shifter has two ports; one is labeled RAM0 and the other is labeled RAM3. Both of these ports consist of a buffer-driver with a three-state output and an input to the multiplexer. Thus, in the shift up mode, the RAM3 buffer is enabled and the RAM0 multiplexer input is enabled. Likewise, in the shift down mode, the RAM0 buffer and RAM3 input are enabled. In the no-shift mode, both buffers are in the high-impedance state and the multiplexer inputs are not selected. This shifter is controlled from the $t_{\rm B}$, $t_{\rm T}$ and $t_{\rm B}$ microinstruction inputs as defined in Figure 9.

Similarly, the Q register is driven from a 3-input multiplexer. In the no-shift mode, the multiplexer enters the ALU data into the Q register. In either the shift-up or shift-down mode, the multiplexer selects the Q register data appropriately shifted up or down. The Q shifter also has two ports; one is labeled Q_0 and the other is Q_3 . The operation of these two ports is similar to the RAM shifter and is also controlled from I_6 , I_7 , and I_8 as shown in Figure 9.

The clock input to the Am2901 controls the RAM, the Q register, and the A and B data latches. When enabled, data is clocked into the Q register on the LOW-to-HIGH transition of the clock. When the clock input is HIGH, the A and B latches are open and will pass whatever data is present at the RAM outputs. When the clock input is LOW, the latches are closed and will retain the last data entered. If the RAM-EN is enabled, new data will be written into the RAM file (word) defined by the B address field when the clock input is LOW.



	MICE	o cod		DURCE ANDS		
12	2 I ₁ I _D Octal Code			R	s	
L	L	L	0		Q	
L	L	н	1	A	8	
L	н	L	2	0	Ω	
L	н	н	3	0	6	
н	Ł	L	4	0	Α	
н	Ł	н	5	5	A	
н	н	L	6	D	a	
н	н	н	17 5 0			

Figure 7.	ALU Source	Operand Control.
	, , , , , , , , , , , , , , , , , , , ,	

	MICE	RO COD	E	ALU		
15	l ₄	l ₄ l ₃ Octa		Function	Symbol	
L	L	Ĺ	0	R Plus S	R+5	
L	L	н	1	S Minus R	5 – R	
Ł	H	L.	2	R Minus \$	R - \$	
C	H	H	3	RORS	RVS	
н	L	L	1 4	FI AND S	RA \$	
н	L	н	5	FI AND \$	R∧s	
н	H	L.	6	R EX-OR S	R ¥ 5	
н	н	н	7	R EX-NOR S	R∀S	

Figure 8. ALU Function Control.

	MI		MICRO CODE			RAM FUNCTION		Q-REG. FUNCTION Y			AM FTER	SHIF) TER
l ₈	ı	7	16	Octal Code	Shift	Load	Shift	Load	ОИТРИТ	RAMO	RAM ₃	ao	03
L	ı		L	0	×	NONE	NONE	F→Q	F	×	×	х	×
L	ı	-	н	1	×	NONE	×	NONE	F	×	х	Х	×
L	1	4	L	2	NONE	F→B	×	NONE	А	×	х	х	×
ι	·	1	н	3	NONE	F→B	×	NONE	F	х	×	x	х
Н		L	L	4	DOMN	F/2 → 8	DOWN	0/2→0	F	F ₀	IN ₃	a_0	IN ₃
н	l	-	н	5	DOWN	F/2 → B	×	NONE	F	Fo	IN ₃	ao	x
н	ŀ	-	L	6	UP	2 F → B	UP	20 → 0	F	INO	F ₃	ING	α ₃
н	1	4	н	7	UP	2F → B	×	NONE	F	1N _O	F ₃	×	a ₃

X* Don't care. Electrically, the shift pin is a TTL input internally connected to a three-state output which is in the high-impedance state.

Figure 9, ALU Destination Control.

\sum	210 OCTAL	0	1	2	3 -	4	5	6	7
OCT 4 3	ALU Source ALU Function	A, Q	А, В	0, 0	О, В	O, A	Ð, A	D, Q	D, O
0	Cn≖L RPlesS	A+Q	A+B	a	В	A	D+A	D+Q	D
L	C _n = H	A+Q+1	A+B+1	Q+1	B+1	A+1	D+A+1	D+Q+1	D+1
1	Cn = L S Minus R	Q-A-1	B-A-1	Q1	8-1	A-1	A~D~1	Q-0-1	7-0-
	C _n = H	Q-A	B-A	a	В	A	A-D	Q-D	p
2	Cn = L R Minus S	A-Q-1	A-8-1	-Q-1	-B-1	-A-1	D-A-1	D-Q-1	D-1
	Cn=H	A-O	A-B	Q	-8	-A	D-A	D-G	D
3	RORS	DVA	ΑVΒ	a	6	А	D∨A	ο∨α	D
4	R AND S	ΑΛQ	A ^ 8	0	0	0	D^A	D ^ C	o.
5	R AND S	Ā∧ū	ĀAB	ā	В	A	ō∧a	ō∧q	0
6	R EX-OR S	A 4 Q	AYB	a	8	A	D∀A	D∀Q	۵
7	R EX-NORS	DYA	A V B	ā	Ē	Ā	D∀À	ō∀ō	۵

^{+ =} P1us; -- = Minus; V = OR; A = AND; ₩ = EX-OR

Figure 10. Source Operand and ALU Function Matrix.



B = Register Addressed by B inputs.

Up is toward MSB, Down is toward LSB.

Octal 543, 210	Group	Function
4 0 4 1 4 5 4 6	AND	A \ Q A \ B D \ A D \ Q
3 0 3 1 3 5 3 6	OR	AVQ AV8 DVA DVQ
6 0 6 1 6 5 6 6	EX-OR	A ♥ Q A ♥ B D ♥ A D ♥ Q
7 0 7 1 7 5 7 6	EX-NOR	Ā∀Q Ā∀B □∀Ā □∀Q
7 2 7 3 7 4 7 7	INVERT	B A D
6 2 6 3 6 4 6 7	PASS	Q B A D
3 2 3 3 3 4 3 7	PASS	Q B A D
4 2 4 3 4 4 4 7	"ZERO"	0 0 0
5 0 5 1 5 5 5 6	MASK	Ā∧Q Ā∧B Ō∧A Ō∧Q

Figure 11. ALU Logic Mode Functions. (Cn Irrelevant)

SOURCE OPERANDS AND ALU FUNCTIONS

There are eight source operand pairs available to the ALU as selected by the I_0 , I_1 , and I_2 instruction inputs. The ALU can perform eight functions; five logic and three arithmetic. The I_3 , I_4 , and I_5 instruction inputs control this function selection. The carry input, C_n , also affects the ALU results when in the arithmetic mode. The C_n input has no effect in the logic mode. When I_0 through I_5 and C_n are viewed together, the matrix of Figure 10 results. This matrix fully defines the ALU/source operand function for each state.

The ALU functions can also be examined on a "task" basis, i.e., add, subtract, AND, OR, etc. In the arithmetic mode, the carry will affect the function performed while in the logic mode, the carry will have no bearing on the ALU output. Figure 11 defines the various logic operations that the Am2901 can perform and Figure 12 shows the arithmetic functions of the device. Both carry-in LOW ($C_{\rm R}=0$) and carry-in HIGH ($C_{\rm R}=1$) are defined in these operations.

Octal	C _n = 0	(Low)	C _n = 1	(High)
¹ 543, ¹ 210	Group	Function	Group	Function
0.0		A+Q		A+Q+1
0 1	ADD	A+B	ADD plus	A+B+1
0.5		D+A	опе	D+A+1
0.6		D+Q		D+Q+1
0 2		a		Q+1
0 3	PASS	В	Increment	B+1
0.4		A		A+1
0 7		D		D+1
1 2		Ω-1	•	a
1 3	Decrement	B-1	PASS	B
14		A-1	!	Α
2 7		D-1		D
2 2		Q-1		-0
2 3	1's Comp.	_B—1	2's Comp.	– В
2 4		A-1	(Negate)	-A
1 7		_D_1		–D
1.0		Q-A-1		Q-A
1 1 1	Subtract	B-A-1	Subtract	B-A
1.5	(1's Comp)	A-D-1	(2's Comp)	A-D
1.6		Q-0-1	ļ	Q-D
2 0		A-Q-1		A-Q
2 1		A-8-1		A-B
2.5		D-A-1	l	D-A
26		D-Q-1	i	DQ

Figure 12. ALU Arithmetic Mode Functions.

LOGIC FUNCTIONS FOR G, P, Cn+4, AND OVR

The four signals G, P, C_{n+4} , and OVR are designed to indicate carry and overflow conditions when the Am2901 is in the add or subtract mode. The table below indicates the logic equations for these four signals for each of the eight ALU functions. The R and S inputs are the two inputs selected according to Figure 7.

Definitions (+ = OR)

$$P_0 = R_0 + S_0$$
 $G_0 = R_0 S_0$
 $P_1 = R_1 + S_1$ $G_1 = R_1 S_1$
 $P_2 = R_2 + S_2$ $G_2 = R_2 S_2$
 $P_3 = R_3 + S_3$ $G_3 = R_3 S_3$

$$C_4 = G_3 + P_3G_2 + P_3P_2G_1 + P_3P_2P_1G_0 + P_3P_2P_1P_0C_0$$

$$C_3 = G_2 + P_2G_1 + P_2P_1G_0 + P_2P_1P_0C_0$$



Function	P	<u> </u>	C _{n+4}	OVR
R+S	P ₃ P ₂ P ₁ P ₀	$G_3 + P_3G_2 + P_3P_2G_1 + P_3P_2P_1G_0$	C ₄	.C ₃ ₩ C ₄
S-R		Same as R + S equations, but sub	stitute $\widehat{\mathbf{R}_{i}}$ for \mathbf{R}_{i} in definitions —	
R – S		Same as R + S equations, but sub	estitute $\overline{S_i}$ for S_i in definitions —	
R∨S	LOW	$P_3P_2P_1P_0$ $P_3P_2P_1P_0 + C_n$		$\overline{P_3P_2P_1P_0} + C_0$
R∧s	LOW	$G_3 + G_2 + G_1 + G_0$	G ₃ + G ₂ + G ₁ + G ₀ + C _n	G ₃ + G ₂ + G ₁ + G ₀ + C _r
Ā∧S	LOW	Same as R ∧ S equation	ons, but substitute $\overline{R_i}$ for R_i in def	initions —
R∀S		Same as $\overrightarrow{R} \overrightarrow{\forall} \overrightarrow{S}$, but substitu	te $\overline{R_{i}}$ for R_{i} in definitions	
R∀S	G ₃ + G ₂ + G ₁ + G ₀	G ₃ +P ₃ G ₂ +P ₃ P ₂ G ₁ +P ₃ P ₂ P ₁ G ₀	$\frac{\overline{G_3} + P_3 G_2 + P_3 P_2 G_1}{+ P_3 P_2 P_1 P_0 (G_0 + \overline{C}_0)}$	See note
	R+S S-R R-S R∨S R∧S R√S	R+S	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	$R+S \qquad \overline{P_3P_2P_1P_0} \qquad \overline{G_3+P_3G_2+P_3P_2G_1+P_3P_2P_1G_0} \qquad C_4$ $S-R \qquad \qquad Same as R+S equations, but substitute $\overline{R_i}$ for R_i in definitions} =$ $R-S \qquad \qquad Same as R+S equations, but substitute $\overline{S_i}$ for S_i in definitions} =$ $R\vee S \qquad \qquad DW \qquad P_3P_2P_1P_0 \qquad \overline{P_3P_2P_1P_0}+C_n$ $R\wedge S \qquad LOW \qquad \overline{G_3+G_2+G_1+G_0} \qquad G_3+G_2+G_1+G_0+C_n$ $\overline{R}\wedge S \qquad LOW \qquad Same as R\wedge S equations, but substitute $\overline{R_i}$ for R_i in definitions} =$ $\overline{R}\vee S \qquad \qquad Same as R\wedge S equations, but substitute $\overline{R_i}$ for R_i in definitions} =$ $\overline{R}\vee S \qquad \qquad Same as R\vee S, but substitute $\overline{R_i}$ for R_i in definitions} =$ $\overline{R}\vee S \qquad \qquad Same as R\vee S =$ $\overline{S} = S_1 + S_2 + S_3 + S_4 + S_4 + S_5 + S_4 + S_5 + $

Figure 13,

FUNCTIONAL DESCRIPTION

The Am2902 is a high-speed, look-ahead carry generator which accepts up to four pairs of carry propagate and carry generate signals and a carry input and provides anticipated carries across four groups of binary ALU's. The device also has carry propagate and carry generate outputs which may be used for further levels of look-ahead.

The Am2902 is generally used with the 2901 bipolar microprocessor unit to provide look-ahead over word lengths of more than four bits. The look-ahead carry generator can be used with binary ALU's in an active LOW or active HIGH input operand mode by reinterpreting-the carry functions. The connections to and from the ALU to the look-ahead carry generator are identical in both cases.

The logic equations provided at the outputs are:

$$\begin{split} &C_{n+x} = G_0 + P_0 C_n \\ &C_{n+y} = G_1 + P_1 G_0 + P_1 P_0 C_n \\ &C_{n+z} = G_2 + P_2 G_1 + P_2 P_1 G_0 + P_2 P_1 P_0 C_n \\ &G = G_3 + P_3 G_2 + P_3 P_2 G_1 + P_3 P_2 P_1 G_0 \\ &P = P_3 P_2 P_1 P_0 \end{split}$$

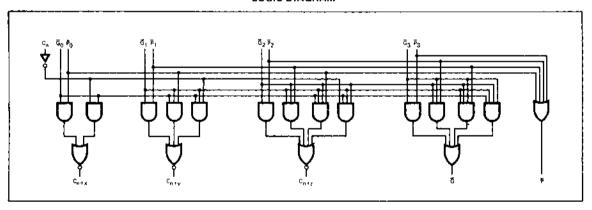
DISTINCTIVE CHARACTERISTICS

- Provides look-ahead carries across a group of four Am2901 microprocessor ALU's
- Capability of multi-level look-ahead for high-speed arithmetic operation over large word lengths
- Typical carry propagation delay of 6 ns
- 100% reliability assurance testing in compliance with MIL-STD-883

ORDERING INFORMATION

Package Type	Temperature Range	Order Number
Molded DIP	0°C to +70°C	AM2902PC
Hermetic DIP	0°C to +70°C	AM2902DC
Dice	0°C to +70°C	AM2902XC
Hermetic DIP	-55°C to +125°C	AM2902DM
Hermetic Flat Pack	-55°C to +125°C	AM2902FM
Dice	-55°C to +125°C	AM2902XM

LOGIC DIAGRAM





MAXIMUM RATINGS (Above which the useful life may be impaired)

Storage Temperature	−65°C to +150°C
Temperature (Ambient) Under Bias	-55°C to +125°C
Supply Voltage to Ground Potential	-0.5V to +7.0V
DC Voltage Applied to Outputs for HIGH Output State	$-0.5V$ to $+V_{CC}$ max.
DC Input Voltage	-0.5V to +5.5V
DC Output Current, Into Outputs	30 mA
DC Input Current	-30 mA to +5,0 mA

ELECTRICAL CHARACTERISTICS OVER OPERATING TEMPERATURE RANGE (Unless Otherwise Noted)

Am2902XC Am2902XM Parameters	T _A = 0°C to +70°C T _A = -55°C to +125°C Description	V _{CC} = 5.0V ±5% (COM'L) MIN. V _{CC} = 5.0V ±10% (MIL) MIN. Test Conditions (No.	= 4.50V	MAX. = 5,25V MAX. = 5,50V Min.	Typ.	Max.	Units
V _{OH}	Output HIGH Voltage	V _{CC} = MIN., I _{OH} = -0.8mA V _{IN} = V _{IH} or V _{IL}	V _{CC} = MIN., I _{OH} = -0.8mA			IVIBA.	Volts
VOL	Output LOW Voltage	V _{CC} = MIN., I _{OL} = 16mA V _{IN} = V _{IH} or V _{IL}			0.2	0.4	Volts
v _{iH}	Input HIGH Level	Guaranteed input logical HIGI voltage for all inputs	4	2.0			Volts
VIL	Input LOW Level	Guaranteed input logical LOW voltage for all inputs	1			0.8	Volts
Vi	Input Clamp Voltage	V _{CC} = MIN., I _{IN} = -12mA				-1.5	Volts
			Cn			-3.2	
			₽3			-4.8	
IIL Input LOW Cui	James I OM Corres	M MAY Nov 0.404	₽ ₂			-6.4	mA
	mput cow current	V _{CC} = MAX., V _{tN} = 0.4V	$\overline{P}_0, \overline{P}_1, \overline{G}_3$			-8.0	
			$\overline{G}_0, \overline{G}_2$			-14,4	
			₫ ₁			-16	
			Cn			80	
			P ₃			120	
1 _{IH}	Input HIGH Current	V _{CC} = MAX., V _{IN} = 2.4V	₽2		l " <u>.</u> _	160	μА
(Note 3)	mpat mon corrent	VCC - MAX., VIN - 2.4V	$\overline{P}_0, \overline{P}_1, \overline{G}_3$			200	μ.
1			$\overline{G}_0, \overline{G}_2$			360]
			G ₁			400	
1)	Input HIGH Current	V _{CC} = MAX., V _{IN} = 5.5V				1.0	mA
^I SC	Output Short Circuit (Note 4)	V _{CC} = MAX., V _{OUT} = 0.0V		40		100	mA
		V _{CC} = MAX.	MIL		62	99	
tcc	Power Supply Current	All Outputs LOW	COM'L		58	94	mA
,,,,	COMBI SUPPLY COTTENT	V _{CC} = MAX.	MIL		37		mA
		All Ouputs HIGH COM'L			35		"'^

Notes: 1. For conditions shown as MIN, or MAX,, use the appropriate value specified under Electrical Characteristics for the applicable device type, 2. Typical limits are at V_{CC} = 5.0V, 25°C ambient and maximum loading, 3. Actual input currents = Unit Load Current X Input Load Factor (see Loading Rules),

4. Not more than one output should be shorted at a time. Duration of the short circuit test should not exceed one second,



Parameter	From (Input)	To (Output)	Test Figure	Test Conditions	Min	Тур	Max	Units
₹PLH	C _s	C _{n+j}	2	$\overline{\overline{P}_0} = \overline{P}_1 = \overline{P}_2 = 0 \text{ V}$ $\overline{G}_0 = \overline{G}_1 = \overline{G}_2 = 4.5 \text{ V}$	i	11	14	ns
tPHL	, , ,	Vn+j		$\overline{\mathbf{G}}_{0} = \overline{\mathbf{G}}_{1} = \overline{\mathbf{G}}_{2} = 4.5 \text{ V}$		11	14	iia.
t _P LH	—		3	$\bar{P}_i = 0 \ V (j > i)$		6.0	8.0	
tPHL	' i	C _{n+i}	3	$\mathbf{C}_{0} = \overline{\mathbf{G}}_{0} = \overline{\mathbf{G}}_{1} = \overline{\mathbf{G}}_{2} = 4.5 \text{ V}$		6.0	8.0	110
tPLH .		C _{n+i}		$\overline{\mathbf{G}}_{i} = 0 \ \forall \ (i > i)$		8.0	10	ns
tPHL	, G, ,	O _{n+j}		$\vec{G}_1 = 0 \text{ V } (i > i)$ $C_n = \vec{P}_0 = \vec{P}_1 = \vec{P}_2 = 4.5 \text{ V}$		8.0	10	113
tp_H	i P ;	G or P	2	$\overline{P}_i = 0 \ \forall \ (i \ge i)$	i	11	14	ns
tpHL	' '; !	461		$\mathbf{C}_{n}^{1} = \overline{\mathbf{G}}_{0} = \overline{\mathbf{G}}_{1} = \overline{\mathbf{G}}_{2} = 4.5 \text{ V}$		11	14	113
t₽LH	G.	G or ₱	2	$\overline{G}_i = 0 \ \forall (i > 1)$	Ī	12	14	ns
tpHL	ı u	G 01 1	_	$\mathbf{C}_{n} = \mathbf{\overline{P}}_{0} = \mathbf{\overline{P}}_{1} = \mathbf{\overline{P}}_{2} = 4.5 \text{ V}$	Γ	12	14	115

 $\mathbf{C}_{\mathbf{n}}$ Carry-in. The carry-in input to the look-ahead generator. Also the carry-in input to the nth Am2901 microprocessor ALU input.

 \mathbf{C}_{n+j} Carry-out, (j = x, y, z). The carry-out output to be used at the carry-in inputs of the n+1, n+2 and n+3 microprocessor ALU slices.

 G_j , P_j Generate and propagate inputs respectively (i = 0, 1, 2, 3). The carry generate and carry propagate inputs from the n, n+1, n+2 and n+3 microprocessor ALU slices.

G, P. Generate and propagate outputs respectively. The carry generate and carry propagate outputs that can be used with the next higher level of carry look-ahead if used.

TRUTH TABLE

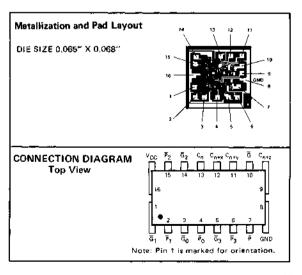
_	Inputs						Outputs		
C,	ā,	Po	G,	P,	Ğ,	P,	Ğ,	P,	C _{n+x} C _{n+y} C _{n+z} G P
	H	P _o H X L							L L H H
XLXH XXLXXH	XHHXLX	XHXXXL	HHHLXX	H X X L L					L L H H
XXXXXXXXX	X X H X X L X	X X X X X X	XHHXLXX	X X X X L L	HHLXXX	H X X X L			L L H H H
	X X X H X X X		X X H X X X L	X X X X X X	X H H X L X X	XXXX	HHLXXX	HXXXX	H H H L L
		H X X L		X X X L		X H X L		X X X	H H H H L

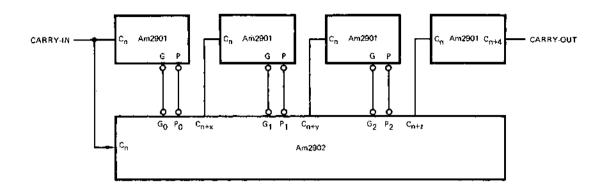
H = HIGH Voltage Level L = LOW Voltage Level X = Don't Care

LOADING RULES (In Unit Loads)

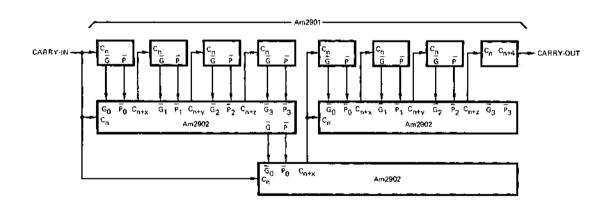
			Fan-out			
Input/Output	Pin No.'s	Input Unit Load	Output HIGH	Output LOW		
<u>G</u> 1	1	8.0	-	-		
Pj	2	4.0	_	_		
<u>G</u> 0	3	7.2				
Po	4	4.0	_			
$\overline{G_3}$	5	4.0	_	_		
P ₃	6	2.4	-	~		
P	7	_	16	8		
GND	8	~	-	-		
C _{p+z}	9		16	8		
G	10		16	8		
C _{n+y}	11		16	8		
C _{n+x}	12	_	16	8		
Cn	13	1.6	-			
G ₂	14	7.2	_			
P ₂	15	3.2	-	_		
Vcc	16	-		_		

A Schottky TTL Unit Load is defined as $50\mu\text{A}$ measured at 2.7V HIGH and -2.0mA measured at 0.6V LOW.





16-BIT CARRY LOOK-AHEAD CONNECTION.



32-BIT ALU, THREE LEVEL CARRY LOOK-AHEAD.



Quad Two-Input OC Bus Transceiver With Three-State Receiver

FUNCTIONAL DESCRIPTION

The Am2905 is a high-performance, low-power Schottky bus transceiver intended for bipolar or MOS microprocessor system applications. The device consists of four D-type edge-triggered flip-flops with a built-in two-input multiplexer on each. The flip-flop outputs are connected to four open-collector bus drivers. Each bus driver is internally connected to one input of a differential amplifier in the receiver. The four receiver differential amplifier outputs drive four D-type latches that feature three-state outputs.

This LSI bus transceiver is fabricated using advanced low-power Schottky processing. All inputs (except the BUS inputs) are one LS unit load. The open-collector bus output can sink up to 100 mA at 0.8V maximum. The BUS input differential amplifier contains disconnect protection diodes such that the bus is fail-safe when power is not applied. The bus enable input $\{\overline{BE}\}$ is used to force the driver outputs to the high-impedance state. When \overline{BE} is HIGH, the driver is disabled. The open-collector structure of the driver allows wired-OR operations to be performed on the bus.

The input register consists of four D-type flip-flops with a buffered common clock and a two-input multiplexer at the input of each flip-flop. A common select input (S) controls the four multiplexers. When S is LOW, the A_i data is stored in the register and when S is HIGH, the B_i data is stored. The buffered common clock (DRCP) enters the data into this driver register on the LOW-to-HIGH transition.

Data from the A or B inputs is inverted at the BUS output. Likewise, data at the BUS input is inverted at the receiver output. Thus, data is non-inverted from driver input to receiver output. The four receivers each feature a built-in

D-type latch that is controlled from the buffered receiver latch enable (\overline{RLE}) input. When the \overline{RLE} input is LOW, the latch is open and the receiver outputs will follow the bus inputs (BUS data inverted and \overline{OE} LOW). When the \overline{RLE} input is HIGH, the latch will close and retain the present data regardless of the bus input. The four latches have three-state outputs and are controlled by a buffered common three-state control (\overline{OE}) input. When \overline{OE} is HIGH, the receiver outputs are in the high-impedance state.

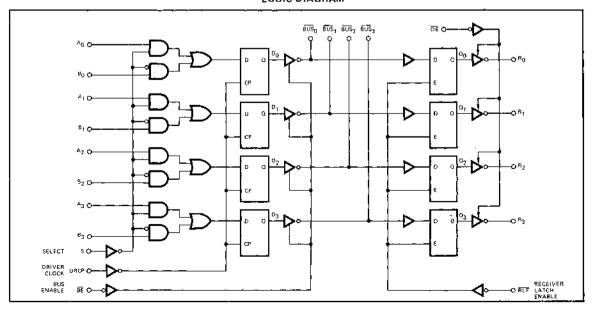
DISTINCTIVE CHARACTERISTICS

- Quad high-speed LSI bus-transceiver
- Open-collector bus driver
- Two-port input to D-type register on driver
- Bus driver output can sink 100 mA at 0.8V max.
- Receiver has output latch for pipeline operation
- Three-state receiver outputs sink 12 mA
- Advanced low-power Schottky processing
- 100% reliability assurance testing in compliance with MIL-STD-883

ORDERING INFORMATION

Package Type	Temperature Range	Order Number
Molded DIP	0°C to +70°C	AM2905PC
Hermetic DIP	0°C to +70°C	AM2905DC
Dice	0°C to +70°C	AM2905XC
Hermetic DIP	-55°C to +125°C	AM2905DM
Hermetic Flat Pak	–55°C to +125°C	AM2905FM
Dice	-55°C to +125°C	AM2905XM

LOGIC DIAGRAM



MAXIMUM RATINGS (Above which the useful life may be impaired)

Storage Temperature	-65°C to +150°C
Temperature (Ambient) Under Bias	-55°C to +125°C
Supply Voltage to Ground Potential	-0.5V to +7V
DC Voltage Applied to Outputs for HIGH Output State	-0.5V to +V _{CC} max.
DC Input Voltage	-0.5V to +5.5V
DC Output Current, Into Outputs (Except Bus)	30mA
DC Output Current, Into Bus	200 mA
DC Input Current	-30mA to +5.0mA

ELECTRICAL CHARACTERISTICS

The following conditions apply unless otherwise noted:

Am2905XC (COM'L) $T_A = 0^{\circ} C \text{ to } +70^{\circ} C$ Am2905XM (M)L) $T_A = -55^{\circ} C \text{ to } +125^{\circ}$

 $T_A = 0^{\circ} C \text{ to } +70^{\circ} C$ $V_{CC}MIN. = 4.75 \text{ V}$ $V_{CC}MAX. = 5.25 \text{ V}$ $T_A = -55^{\circ} C \text{ to } +125^{\circ} C$ $V_{CC}MIN. = 4.50 \text{ V}$ $V_{CC}MAX. = 6.50 \text{ V}$

BUS INPUT/OUTPUT CHARACTERISTICS OVER OPERATING TEMPERATURE RANGE

Parameters	Description	Test Co	Test Conditions (Note 1)				Max.	Units	
			I _{OL} = 40mA			0.32	0.5		
VOL	Bus Output LOW Voltage	$v_{CC} = MIN$.	10L = 70mA			0.41	0.7	Volts	
·			IOL = 100mA	IOL = 100mA			0,8		
			V _O ≈ 0.4∨				-50		
10	Bus Leakage Current	VCC = MAX.	VO = 4.5V	MIL			200	μA	
			. 70 4.37	COM'L			100		
FOFF	Bus Leakage Current (Power OFF)	V _O = 4.5V					100	μА	
VTH	Receiver Input HIGH	Bus enable = 2.4\		MIL	2.4	2.0		14.14.	
''H	Threshold	Bus enable - 2,41	,	Ç0M'L	2.3	2.0		Volts	
VTL	Receiver Input LOW	Bus enable = 2.4\	Pin cooks = 3.4V			2.0	1,5	Volts	
-16	Threshold	50s Glable - 2.41	COM'L			2.0	1.6	1 2016	



ELECTRICAL CHARACTERISTICS

The following conditions apply unless otherwise noted:

Am2905XC (COM'L) $T_A = 0^{\circ} C \text{ to } +70^{\circ} C$ $V_{CC} MIN. = 4.75 V$ $V_{CC} MAX. = 5.25 V$ Am2905XM MIL) $T_A = -55^{\circ} C \text{ to } +125^{\circ} C$ $V_{CC} MIN. = 4.50 V$ $V_{CC} MAX. = 5.50 V$

ameters	Description	FICS OVER OPERATING TEMPERATURE RANGE Test Conditions (Note 1)				Typ. (Note 2)	Max.	Units
VOH	Receiver Output	VCC = VIN	MIL, IQH	-1.0mA	2.4	3.4		
VOH	HIGH Voltage	VIN = VIE or VIH	IN = V _{IL} or V _{IH} COM'L, 1 _{OH} = -2.6mA		2.4	3.4		Volts
		V addh				0.27	0.4	
VOL	Receiver Output LOW Voltage	V _{CC} = MIN. V _{IN} = V _{II} or V _{IH}	IOL = 8mA			0.32	0.45	Volts
		AIM AIT OLAIM	1 _{OL} = 12m	A		0.37	0.5	
VIH	Input HIGH Level (Except Bus)	Guaranteed input logi for all inputs	Guaranteed input logical HIGH for all inputs			i		Volts
.,	Input LOW Level	Guaranteed input logi	ical LOW	MIL		!	0.7	
VIL	(Except Bus)	for all inputs	for all inputs COM'L				0.B	Volts
v,	Input Clamp Voltage (Except Bus)	VCC = MIN., IIN = -			1	-1.5	Volts	
r _{IL}	Input LOW Current (Except Bus)	V _{CC} = MAX., V _{IN} = 1	0.4∨				0.36	mA
чн	Input HIGH Current (Except Bus)	VCC = MAX., VIN =	2.7V				20	μА
I _I	Input HIGH Current (Except Bus)	V _{CC} = MAX., V _{IN} =	5.5V				100	μА
10	Voc = MAX.		V _O = 2.4 V				20	
.0			V _O = 0.4 V			-20	μА	
¹sc	Receiver Output Short Circuit Current	VCC = MAX.			-12		-65	mΑ
lcc	Power Supply Current	V _{CC} = MAX., All inpu	uts = GND			69	105	mA

	IG CHARACTERISTICS OVER NG TEMPERATURE RANGE		,	4m2905XI Тур,	М	,	2			
arameters	Description	Test Conditions	Min.	(Note 2)	Max.	Min.	Typ. (Note 2)	Max.	Units	
tPHL	Driver Clock (DRCP) to Bus			21	40	I	21	36	ns	
tPLH .	Driver Clock (DRCF) to Bus	C_L (BUS) = $50 \rho F$		21	40		21	36	211	
tPHL	Bus Enable (BE) to Bus	AL (BUS) = 501)		13	26		13	23	ns	
tPLH	Bus Evable (RE) to Bus			13	26	Γ	13	23	113	
tş	5		25	1		23			П\$	
th	Data Inputs (A or 8)		8.0			7.0			115	
tş	5-1		33			30			ns	
th	Select Input (\$)		0.8			7.0			l la	
tPW	Driver Clock (DRCP) Pulse Width (HIGH)		28			25			nş	
tPLH	Bus to Receiver Output			18	37	T	18	34	ns	
tPHL .	(Latch Enable)	CL = 15pF		18	37		18	34	ns	
tPLH	Latch Enable to Receiver Output	$R_{\rm L}$ = 2.0 k Ω		21	37		21	34	ns	
tPHL .	Laten Enable to Receiver Output	1		21	37		21	34] ''	
ts	Toward and Franks (D. C .)	ı	21	T		18	T		ns	
th	Bus to Latch Enable (RLE)		7.0			5.0			""	
tZH	00			14	28	1	14	25	ns	
^t ZL	Output Control to Receiver Output			14	28]	14	25	''s	
tHZ	Output Control on Brandway Co.	C _L = 5.0 pF		14	28		14	25	ns	
tLZ	Output Control to Receiver Output	$R_{L} = 2.0 \text{ K}\Omega$		14	28		14	25	''s	

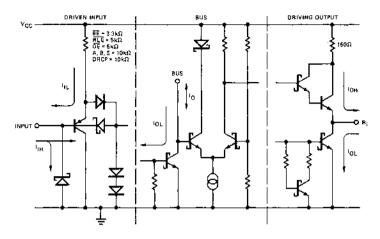
Notes: 1. For conditions shown as MIN, or MAX,, use the appropriate value specified under Electrical Characteristics for the applicable device type.

2. Typical limits are at V_{CC} = 5.0 V, $25^{\circ}C$ ambient and maximum loading.

3. Not more than one output should be shorted at a time. Duration of the short circuit test should not exceed one second.

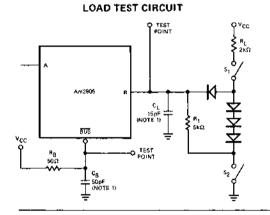


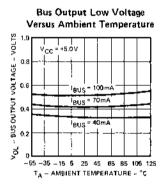
INPUT/OUTPUT CURRENT INTERFACE CONDITIONS

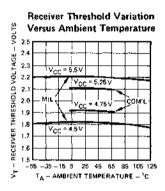


Note: Actual current flow direction shown.

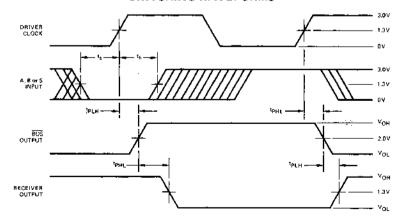
TYPICAL PERFORMANCE CURVES







SWITCHING WAVEFORMS



Note: Bus to Receiver output delay is measured by clocking data into the driver register and measuring the BUS to R combinatorial delay.



FUNCTION TABLE

			INPUT	s				RNAL EVIÇE	BUS	ОШТРИТ	FUNCTION		
S	Ai	Вį	DRCP	ΒĒ	RLE	ŌĒ	Di	α_{i}	BUSi	Ri	, 0.101.011		
Х	Х	Х	Х	Н	×	Х	Х	X	Z	Х	Driver output disable		
X	Х	Х	Х	x	Х	Н	Х	Х	Х	Z	Receiver output disable		
Х	X.	Х	Х	н	L	L	х	L	L	Н	Driver output disable and receive data		
×	х	х	x	н	L	L	×	н	н	L	via Bus input		
Х	х	Х	Х	×	Н	Х	Х	NC	Х	Х	Latch received data		
L	L	Х	†	Х	Х	Х	L,	Х	X	Х			
L	н	х	1	×	х	×	н	х	х	×	Load driver register		
н	х	L	1	Х	×	х	L	х	Х	×	Coad dittel tediziel		
Н	х	Н	1	X	Х	Х	Н	. x	Х	×			
х	х	х	L	Х	х	Х	NC	×	Х	×	No driver clock restrictions		
×	х	х	н	х	×	x	NC	×	Х	×	NO GIVE CIOCK TEST ICTIONS		
х	х	х	×	L	х	х	L	Х	н	×	Daine Bure		
×	х	х	×	L	x	х	Н	Х	L	x	Drive Bus		

H = HIGH L = LOW Z = HIGH Impedance NC = No change X = Don't care

1 ~ LOW-to-HIGH transition

i = 0, 1, 2, 3

DEFINITIONS

A₀, A₁, A₂, A₃ The "A" word data input into the two input multiplexer of the driver register.

B₀, B₁, B₂, B₃ The "B" word data input into the two

input multiplexers of the driver register.

Select. When the select input is LOW, the A data word is applied to the driver register. When the select input is HIGH, the B word is applied to the driver register.

DRCP Driver Clock Pulse, Clock pulse for the

driver register.

BE Bus Enable. When the Bus Enable is HIGH, the four drivers are in the high impedance

state.

BUS₀, BUS₁ The four driver outputs and receiver inputs (data is inverted).

Ro, R1, R2, R3 The four receiver outputs. Data from the

bus is inverted while data from the A or B

inputs is non-inverted,

RLE Receiver Latch Enable. When RLE is

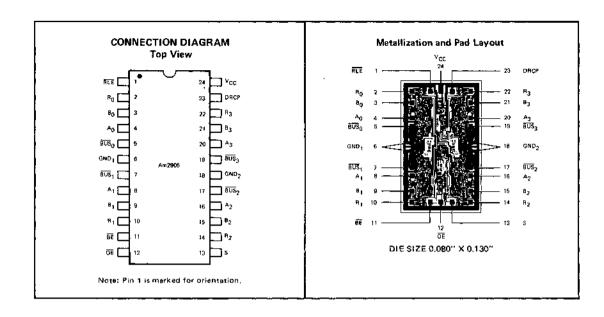
LOW, data on the BUS inputs is passed through the receiver latches. When RLE is HIGH, the receiver latches are closed and will retain the data independent of

all other inputs.

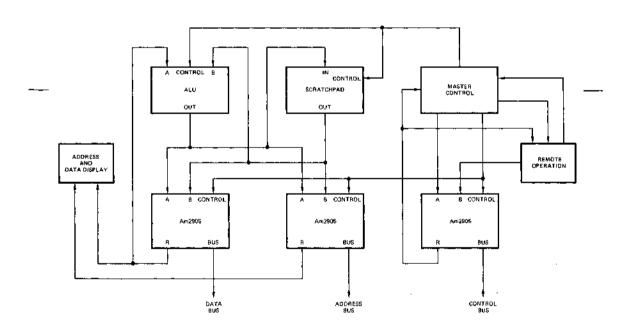
OE Output Enable. When the OE input is

HIGH, the four three state receiver outputs are in the high-impedance state.





APPLICATIONS



The Am2905 is a universal Bus Transceiver useful for many system data, address, control and timing input/output interfaces.



FUNCTIONAL DESCRIPTION

The Am2906 is a high-performance, low-power Schottky bus transceiver intended for bipolar or MOS microprocessor system applications. The device consists of four D-type edge-triggered flip-flops with a built-in two-input multiplexer on each. The flip-flop outputs are connected to four open-collector bus drivers. Each bus driver is internally connected to one input of a differential amplifier in the receiver. The four receiver differential amplifier outputs drive four D-type latches. The device also contains a four-bit odd parity checker/generator.

This LSI bus transceiver is fabricated using advanced low-power Schottky processing. All inputs (except the BUS inputs) are one LS unit load. The open-collector bus output can sink up to 100 mA at 0.8V maximum. The BUS input differential amplifier contains disconnect protection diodes such that the bus is fail-safe when power is not applied. The bus enable input (BE) is used to force the driver outputs to the high-impedance state. When BE is HIGH, the driver is disabled. The open-collector structure of the driver allows wired-OR operations to be performed on the bus.

The input register consists of four D-type flip-flops with a buffered common clock and a two-input multiplexer at the input of each flip-flop. A common select input (S) controls the four multiplexers. When S is LOW, the A_j data is stored in the register and when S is HIGH, the B_j data is stored. The buffered common clock (DRCP) enters the data into this driver register on the LOW-to-HIGH transition.

Data from the A or B input is inverted at the BUS output. Likewise, data at the BUS input is inverted at the receiver output. Thus, data is non-inverted from driver input to

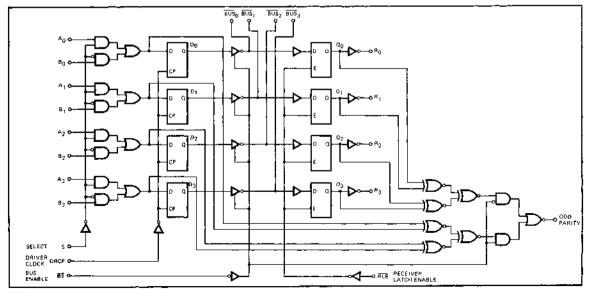
receiver output. The four receivers each feature a built-in D-type latch that is controlled from the buffered receiver latch enable (\overline{RLE}) input, When the \overline{RLE} input is LOW, the latch is open and the receiver outputs will follow the bus inputs (BUS data inverted). When the \overline{RLE} input is HIGH, the latch will close and retain the present data regardless of the bus input.

The Am2906 features a built-in four-bit odd parity checker/generator. The bus enable input (\overline{BE}) controls whether the parity output is in the generate or check mode. When the bus enable is LOW (driver enabled), odd parity is generated based on the A or B field data input to the driver register. When \overline{BE} is HIGH, the parity output is determined by the four latch outputs of the receiver. Thus, if the driver is enabled, parity is generated and if the driver is in the high-impedance state, the BUS parity is checked.

DISTINCTIVE CHARACTERISTICS

- · Quad high-speed LSI bus transceiver.
- Open-collector bus driver.
- Two-port input to D-type register on driver.
- Bus driver output can sink 100 mA at 0.8V max.
- Internal odd 4-bit parity checker/generator.
- Receiver has output latch for pipeline operation.
- Receiver outputs sink 12 mA.
- Advanced low-power Schottky processing.
- 100% reliability assurance testing in compliance with MIL-STD-883.

LOGIC DIAGRAM



ORDERING INFORMATION

Package Type	Temperature Range	Order Number
Molded DIP	0°C to +70°C	AM2906PC
Hermetic DIP	0°C to +70°C	AM2906DC
Dice	0°C to +70°C	AM2906XC
Hermetic DIP	-55°C to +125°C	AM2906DM
Hermetic Flat Pak	–55°C to +125°C	AM2906FM
Dice	-55°C to +125°C	AM2906XM

MAXIMUM RATINGS (Above which the useful life may be impaired)

Storage Temperature	-65°C to +150°C
Temperature (Ambient) Under Bias	−55°C to +125°C
Supply Voltage to Ground Potential	-0.5V to +7V
DC Voltage Applied to Outputs for HIGH Output State	-0.5V to +V _{CC} max.
DC Input Voltage	-0.5V to +5.5V
DC Output Current, Into Outputs (Except Bus)	30mA
DC Output Current, Into Bus	200 mA
DC Input Current	-30mA to +5.0mA

ELECTRICAL CHARACTERISTICS

The following conditions apply unless otherwise noted:

Am2906XC (COM*L) Am2906XM (MIL)

BUS INPUT/OUTPUT CHARACTERISTICS OVER OPERATING TEMPERATURE RANGE

Parameters	Description	Test Co	Test Conditions (Note 1)			Typ. (Note 2)	Max.	Units
			IOL = 40mA			0.32	0.5	
VOL	Bus Output LOW Voltage		IOL = 70mA			0.41	0.7	Volts
			IOL = 100mA	I _{OL} = 100mA			8.0	1
			V _O = 0.4V			I	-50	· ·
Io I	Bus Leakage Current	V _{CC} = MAX.	VO = 4.5V	MIL			200	μА
			VO - 4.54	COM'L			100	
IOFF	Bus Leakage Current (Power OFF)	V ₀ = 4.5V					100	μA
V _{TH}	Receiver Input HIGH	Pur apable = 2.41	1	MIL	2.4	2.0		Volts
	Threshold	Bus enable - 2,41	Bus enable = 2.4V COM'L		2.3	2.0		VOITS
VTL	Receiver Input LOW	Rus enable = 2.4\	Bus enable = 2.4V MIL COM'L			2.0	1,5	Volts
-10	Threshold	503 511dbic 2.41				2.0	1.6	



ELECTRICAL CHARACTERISTICS

The following conditions apply unless otherwise noted:

TA = 0°C to +70°C Am2906XC (COM'L) Am2906XM (MIL)

TA - -55°C to +125°C VCC MIN, 4.5V

V_{CC} MIN, = 4.75V

VCC MAX. = 5.25V VCC MAX. - 5.5V

DC CHARACTERISTICS OVER OPERATING TEMPERATURE RANGE Typ. Description Test Conditions (Note 1) Min. Max. Units **Parameters** (Note 2) I_{OH} = -1mA MIL 24 3.4 VCC = MIN. Receiver Output HIGH Voltage VIN = VIL or VIH I_{OH} = -2.6mA 24 3.4 COM'L V_{OH} Volts $V_{CC} = MIN_{\odot} 1_{OH} = -660 \mu A$ MIL 2.5 3.4 Parity Output HIGH Voltage $V_{IN} = V_{IH} \text{ or } V_{IL}$ COM'L 2.7 3.4 IOL = 4mA 0.27 0.4 VCC = MIN. Output LOW Voltage VOL 0.32 0.45 IOL = 8mA Volts (Except Bus) VIN = VIL or VIH 0.37 0.5 I_{OL} = 12mA Input HIGH Level Guaranteed input logical HIGH v_{iH} 2.0 Volts (Except Bus) for all inputs 0.7 MIL Input LOW Level Guaranteed input logical LOW $v_{\rm ft}$ Volts (Except Bus) for all inputs 8.0 COM'L Input Clamp Voltege ٧ı V_{CC} = MIN., I_{IN} = -18mA -1.2Volts (Except Bus) Input LOW Current -0.36mΑ HΔ $V_{CC} = MAX., V_{IN} = 0.4V$ (Except Bus) Input HIGH Current 'nн $V_{CC} = MAX., V_{1N} = 2.7V$ 20 μА (Except Bus) Input HIGH Current 1 100 μΑ $V_{CC} = MAX_{..} V_{EN} = 5.5V$ (Except Bus) **Output Short Circuit Current** Isc VCC = MAX. -12-65 mΑ (Except Bus) Power Supply Current V_{CC} = MAX., All inputs = GND 105 mΑ **Icc** 72

	G CHARACTERISTICS ERATING TEMPERATURE RA	NCE		\m2906XI	VI.	£	.m2906X	C	
arameters	Description	Test Conditions	Min.	Typ. (Note 2)	Max.	Min.	Typ. (Note 2)	Max.	Units
tPHL				21	40	1	21	36	
tp_H	Driver Clock (DRCP) to Bus	C _L (BUS) = 50pF		21	40		21	36	ns
†PHL	Bus Enable (BE) to Bus	RuiBUSI = 50 Ω		13	26	1	13	23	ПS
tPLH .	Bus chable (BE) to Bus			13	26		13	23] "
t _s	Data Inputs (A or B)		25			23	Ţ		пѕ
t _h	Data inputs (A or B)		8.0			7.0] ''s
t _s	Calcada a 181	1	33		<u>.</u> ,	30			ns
th	Select Inputs (S)		8.0	1		7.0			1 "
tpW	Clock Pulse Width (HtGH)	1	28			25			пş
tp_H	Bus to Receiver Output		"	18	37		18	34	ns
t _{PHL}	(Latch Enabled)			18	37]	18	34] ''`
tPLH :	A seed Food last Books Access			21	37		21	34	
tPHL.	Latch Enable to Receiver Output	C _L = 15pF		21	37		21	34	ns
t _S	Provide the transfer of the Control	$R_L = 2.0 k\Omega$	21			18			
th	Bus to Latch Enable (RLE)		7.0			5.0			ns
tPLH	A or B Data to Odd Parity Output]		21	40		21	36	
^t PHL	(Driver Enabled)			21	40		21	36	лş
tp_H	Bus to Odd Parity Output	1		21	40		21	36	
tPHL	(Driver Inhibited, Latch Enabled)			21	40		21	36	ns
tPLH	Latch Enable (RLE) to	1	-	21	40		21	36	
tPHL	Odd Parity Output			21	40		21	36	ns

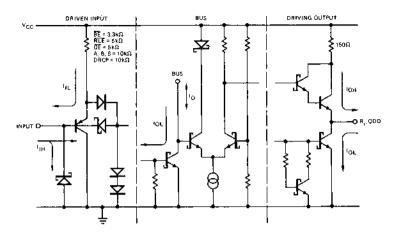
Notes: 1. For conditions shown as MIN, or MAX,, use the appropriate value specified under Electrical Characteristics for the applicable device type.

^{3.} Not more than one output should be shorted at a time. Duration of the short directit test should not exceed one second.



^{2.} Typical limits are at V_{CC} = 5.0 V, 25° C ambient and maximum loading.

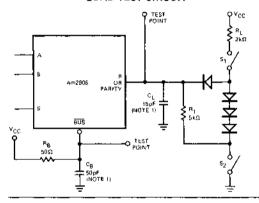
INPUT/OUTPUT CURRENT INTERFACE CONDITIONS

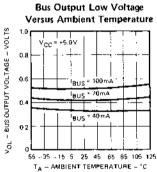


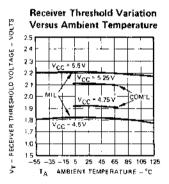
Note: Actual current flow direction shown,

TYPICAL PERFORMANCE CURVES

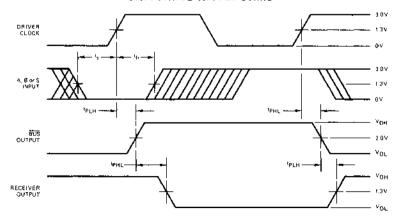
LOAD TEST CIRCUIT







SWITCHING WAVEFORMS



Note: Bus to Receiver output delay is measured by clocking data into the driver register and measuring the BUS to R combinatorial delay.



FUNCTION TABLE

		INP	UTS				ERNAL DEVICE	eus	OUTPUT	FUNCTION		
s	A _i	Вį	DRCP	BE	RLE	D;	Ωį	BUSi	R _i			
×	×	x	×	н	×	x	x	z	×	Driver output disable		
×	X	Х	Х	Н	L	Х	L	L	Н	Driver output disable and receive data		
х	х	×	х	н	L	×	н	н	L	via Bus input Latch received data		
Х	Х	х	×	X	н	×	NC	Х	х			
L	L	Х	†	X	х	L	×	Х	Х			
L	н	х	t	х	x	н	×	х	х	Load driver register		
н	Х	L	†	x	х	L .	×	x	×	Load driver register		
н	X	н	Ť	×	x	н	×	×	×			
Х	X	Х	L	х	×	NC	×	Х	×	No driver clock restrictions		
х	×	×	н	х	×	NC	×	×	×	No driver clock restrictions		
Х	X	х	×	L	×	L	х	н	×	Drive bus		
X	Х	X	х	L	x	н	х	L	x	Drive ous		

H = HIGH

Z = HIGH Impedance

X = Don't care

i = 0, 1, 2, 3

L = LOW

NC = No change

1 = LOW-to-HIGH transition

DEFINITIONS

The "A" word data input into the two A₀, A₁, A₂, A₃ input multiplexer of the driver register.

The "B" word data input into the two B₀, B₁, B₂, B₃ input multiplexers of the driver register.

S Select. When the select input is LOW, the

A data word is applied to the driver register. When the select input is HIGH, the B word is applied to the driver register.

DRCP Driver Clock Pulse, Clock pulse for the

driver register.

BE Bus Enable, When the Bus Enable is HIGH, the four drivers are in the high impedance

BUSO, BUS1 BUS2, BUS2 The four driver outputs and receiver inputs (data is inverted),

Ro, R1, R2, R3 The four receiver outputs. Data from the

bus is inverted while data from the A or B inputs is non-inverted.

RLE Receiver Latch Enable. When RLE is LOW, data on the BUS inputs is passed through the receiver latches. When RLE is HIGH, the receiver latches are closed

> and will retain the data independent of all other inputs.

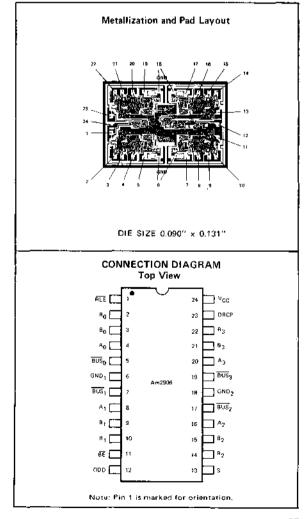
ODD ODD Parity output. Generates parity

with the driver enabled, checks parity with the driver in the HIGH-imped-

ance state.

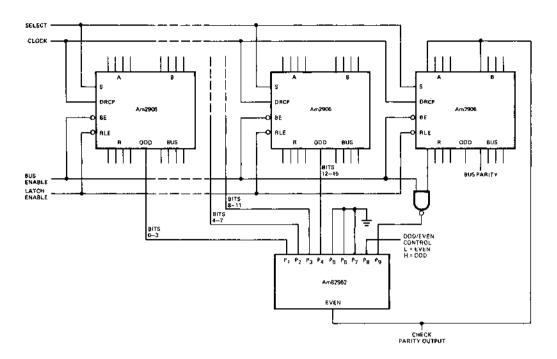
PARITY OUTPUT FUNCTION TABLE

S	BE	ODD PARITY OUTPUT				
L	L	ODD = A0 + A1 + A2 + A3				
Н	L	ODD = B0 + B1 + B2 + B3				
×	н	ODD = Q ₀ + Q ₁ + Q ₂ + Q ₃				





APPLICATIONS



Generating or checking parity for 16 data bits.



FUNCTIONAL DESCRIPTION

RAYTHEON

The Am2907 is a high-performance, low-power Schottky bus transceiver intended for bipolar or MOS microprocessor system applications. The device consists of four D-type edge-triggered flip-flops. The flip-flop outputs are connected to four opencollector bus drivers. Each bus driver is internally connected to one input of a differential amplifier in the receiver. The four receiver differential amplifier outputs drive four D-type latches, that feature three-state outputs. The device also contains a four-bit odd parity checker/generator.

This LSI bus transceiver is fabricated using advanced lowpower Schottky processing. All inputs (except the BUS inputs) are one LS unit load. The open-collector bus output can sink up to 100 mA at 0.8 V maximum. The BUS input differential amplifier contains disconnect protection diodes such that the bus is fail-safe when power is not applied. The bus enable input (BE) is used to force the driver outputs to the high-impedance state. When BE is HIGH, the driver is disabled. The opencollector structure of the driver allows wired-OR operations to be performed on the bus.

The input register consists of four D-type flip-flops with a buffered common clock. The buffered common clock (DRCP) enters the A_i data into this driver register on the LOW-to-HIGH

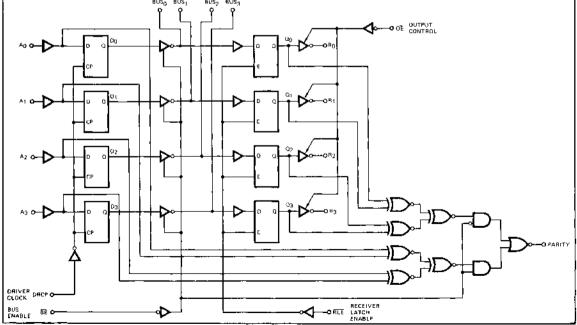
Data from the A input is inverted at the BUS output, Likewise, data at the BUS input is inverted at the receiver output. Thus, data is non-inverted from driver input to receiver output. The four receivers each feature a built-in D-type latch that is controlled from the buffered receiver latch enable (RLE) input, When the RLE input is LOW, the latch is open and the receiver outputs will follow the bus inputs (BUS data inverted and \overline{OE} LOW). When the RLE input is HIGH, the latch will close and retain the present data regardless of the bus input. The four latches have three-state outputs and are controlled by a buffered common three-state control (\overline{OE}) input. When \overline{OE} is HIGH, the receiver outputs are in the high-impedance state.

The Am2907 features a built-in four-bit odd parity checker/ generator. The bus enable input (8E) controls whether the parity output is in the generate or check mode. When the bus enable is LOW (driver enabled), odd parity is generated based on the A field data input to the driver register. When BE is HIGH, the parity output is determined by the four latch outputs of the receiver. Thus, if the driver is enabled, parity is generated and if the driver is in the high-impedance state, the BUS parity is checked.

DISTINCTIVE CHARACTERISTICS

- Quad high-speed LSI bus-transceiver
- Open-collector bus driver.
- D-type register on driver.
- Bus driver output can sink 100 mA at 0.8 V max.
- Internal odd 4-bit parity checker/generator
- Receiver has output latch for pipeline operation
- Three-state receiver outputs sink 12 mA
- Advanced Low-Power Schottky processing
- 100% reliability assurance testing in compliance with MIL-STD-883

LOGIC DIAGRAM BUS_O BUS₁ BUS₂ QUIS.



ORDERING INFORMATION

	Package Type	Temperature Range	Order Number		
	Molded DIP	0°C to +70°C	AM2907PC		
Н	ermetic DIP	0°C to +70°C	AM2907DC		
	Dice	0°C to +70°C	AM2907XC		
Н	ermetic DIP	-55°C to +125°C	AM2907DM		
* Her	metic Flat Pak	-55°C to +125°C	AM2907FM		
	Dice	-55°C to +125°C	AM2907XM		

^{*}Available on special order

MAXIMUM RATINGS (Above which the useful life may be impaired)

Storage Temperature	-65°C to +150°C
Temperature (Ambient) Under Bias	-55°C to +125°C
Supply Voltage to Ground Potential	-0.5 V to +7 V
DC Voltage Applied to Outputs for HIGH Output State	-0.5 V to +V _{CC} max.
DC Input Voltage	-0.5 V to +5.5 V
DC Output Current, Into Outputs (Except BUS)	30 mA
DC Output Current, Into Bus	200 mA
DC Input Current	-30 mA to +5.0 mA

ELECTRICAL CHARACTERISTICS

The following conditions apply unless otherwise noted:

BUS INPUT/OUTPUT CHARACTERISTICS OVER OPERATING TEMPERATURE RANGE

Parameters	Description	Test Conditions (Note 1)			Min,	(Note 2)	Max.	Units
		V _{CC} = MIN.	IOL = 40mA			0.32	0.5	
VOL	Bus Output LOW Voltage		IOL = 70mA			0.41	0.7	Volts
	İ		I _{OL} = 100mA			0.55	0.8	
	Bus Leakage Current	V _{CC} = MAX.	V _O = 0.4 V				50	
10			V _O = 4.5 V	MIL			200	μА
!				COM'L			100	
OFF	Bus Leakage Current (Power Off)	V _O = 4.5 V					100	μА
	Receiver Input HIGH Threshold	Bus Enable = 2.4 V MIL COM		MIL	2.4	2.0	-	Volts
VTH				COM'L	2.3	2.0		
V	Paris I CH Thursday	Bus Enable ≃ 2.4 V		MIL		2,0	1.5	Volts
VTL	Receiver Input LOW Threshold			COM'L		2.0	1.6	



ELECTRICAL CHARACTERISTICS

The following conditions apply unless otherwise noted:

DC CHARACTERISTICS OVER OPERATING TEMPERATURE RANGE

Parameters	Description	Test Con	Min.	Typ. (Note 2)	Max.	Units		
	Receiver	V _{CC} = MIN.	MIL: IOH =	-1mA	2.4	3.4		Malte
ч он	Output HIGH Voltage	VIN = VIL or VIH	COM'L: TOH	COM'L: IOH = -2.6mA		3.4		Volts
V	Parity	V _{CC} = MIN., I _{OH} =	-660µA	MIL	2.5	3.4		
VOH	Output HIGH Voltage	VIN = VIH or VIL		COM'L	2.7	3.4		Volts
ľ	Output LOW Voltage	VCC = MIN.	10L = 4mA	1		0.27	0.4	
v ol	(Except Bus)	VIN = VIL or VIH	1 _{QL} = 8mA			0.32	0.45	Volts
	(Except pas)	AIM - AIT OLAIH	10L = 12mA			0.37	0.5	1
VIH	Input HIGH Level	Guaranteed input to	gical HIGH		2.0			Volts
VIH	(Except Bus)	for all inputs			2.0	l		40.13
V	Input LOW Level	Guaranteed input lo	gical LOW	MrL			0.7	Volts
VIL	(Except Bus)	for all inputs	COM'L			0.8	7 0,115	
VI	Input Clamp Voltage	V _{CC} = MIN., 1 _{IN} = -18mA					-1.2	Volts
*1	(Except Bus)	A GC - MILM" 1 M -	- Ionia				-1.2	Voits
lu.	Input LOW Current	V _{CC} = MAX., V _{IM}	= 0.4 V				-0.36	mA
hL.	(Except Bus)	ACC - MINAY AIM	· 0.4 V	_			~0.30	""
	Input HIGH Current	Vcc = MAX., V _{IM} = 2.7 V				20	μА	
•ин	(Except Bus)	AGG - MINY ALM	- 2.7 V				20	"
1.	Input HIGH Current	V _{CC} = MAX., V _{IAI}	1			1	100	μΑ
1	(Except Bus)	ACC - MINY ATM	- 5.5 V			1	100	μ.Α.
laa	Output Short Circuit	Voc - MAY			-12		~-65	mA
Isc	Current (Except Bus)	V _{CC} = MAX.				 		IIIA
Icc	Power Supply Current	V _{CC} = MAX., All Ir	iputs = GND			75	110	mA
la	Off-State Output Current	V _{CC} = MAX.	Vo = 2.4 V				20	μΑ
'0	(Receiver Outputs)	v _{CC} = MAX.					-20] ~~

	IG CHARACTERISTICS ERATING TEMPERATURE RA Description	ANGE Test Conditions	A Min,	m2907X/ Typ. (Nate 2)	VI Max,	A Min.	m 2907X(Typ. (Note 2)	: Max.	Units
†PHL	······································		T	21	40	Γ	21	36	Ţ
1PLH	Driver Clock (DRCP) to Bus	C ₁ (BUS) = 50pF	 -	21	40		21	36	п\$
TPHL .		R _L (BUS) = 50Ω		13	26		13	23	
t _{PLH}	Bus Enable (BE) to Bus			13	26		13	23	ns
t _s			25	†-··-		23	ļ		1 .
th	A Data Inputs		8.0			7.0			ns ns
tpw	Clock Pulse Width (HIGH)		28			25			nş
tPLH	Bus to Receiver Output			18	37		18	34	
tPHL.	(Latch Enabled)			18	37		18	34	ns
tPLH	Latch Enable to Receiver Output			21	37		21	34	ns
tPHL i	Laten Enable to Receiver Dutput	0 -45-5		21	37	I	21	34] "5
ts	0	- C _L = 15pF R ₁ = 2.0kΩ	21			18	T i		- ns
th	Bus to Latch Enable (RLE)	11 2.0104	7.0			5.0			
tPLH	A Data to Odd Parity Out			21	40		21	36	ns
tPHL	(Driver Enabled)			21	40		21	36] ""
tPLH	Bus to Odd Parity Out			21	40	l	21	36	ns
tPHL	(Oriver Inhibit)			21	40		21	36] ''*
tPLH ;	Latch Enable (RLE) to Odd			21	40		21	36	05
^t PHL	Parity Output			21	40		21	36] ''.
tzн	Output Control to Output			14	28		14	2 5	ns
tzL	Output Control to Output			14	28		14	25	1 "
tHZ	Output Control to Output	CL = 5.0pF		14	28		14	25	ns
tLZ	Output Control to Output	R _L = 2.0kΩ		14	28		14	25	""

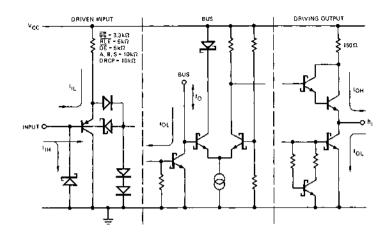
Notes: 1. For conditions shown as MIN, or MAX, use the appropriate value specified under Electrical Characteristics for the applicable davice type.

^{3.} Not more than one output should be shorted at a time. Duration of the short circuit test should not exceed one second.



^{2.} Typical limits are at V_{CC} = 5.0V, $25^{\circ}C$ ambient and maximum loading.

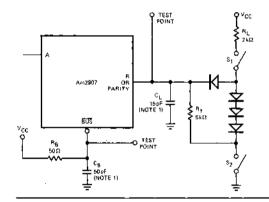
INPUT/OUTPUT CURRENT INTERFACE CONDITIONS

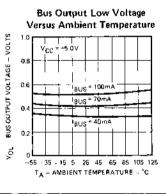


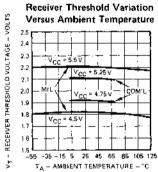
Note: Actual current flow direction shown,

TYPICAL PERFORMANCE CURVES

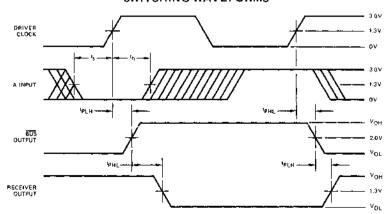
LOAD TEST CIRCUIT







SWITCHING WAVEFORMS



Note: Bus to Reciever output delay is measured by clocking data into the driver register and measuring the BUS to R combinatorial delay.



TRUTH TABLE

	INPUTS				INTE TO DI	RNAL EVICE	BUS	оитрит	FUNCTION
Ai	DRCP	ΒE	RLE	ŌĒ	Di	Qi	Bį	Ri	
X	Х	Н	×	X	×	×	н	X	Driver output disable
×	Х	Х	Х	н	Х	Х	Х	Z	Receiver output disable
Х	Х	н	L	L	×	L	L	н	Driver output disable and receive data
х	x	н	Ļ	L	×	Н	н	L	via Bus input
X	Х	Х	Н	Х	х	NC	Х	×	Latch received data
L	1	Х	х	×	L	X	Х	Х	Look driver engineer
н	1	×	X	×	н	х	×	_x _	Load driver register
X	L	х	X	×	NC	×	х	×	No driver clock restrictions
x	н	х	х	x	NC	х	х	x	NO ditact clock restrictions
х	Х	L	×	×	L	Х	Н	X	6: 8
×	x	Ł	×	х	н	×	L	x	Drive Bus

H = HIGH L = LQW Z = High Impedance NC = No Change X = Don't Care

1 = LOW-to-HIGH Transition

i = 0, 1, 2, 3

DEFINITIONS

DRCP Driver Clock Pulse. Clock pulse for the driver register.

BE Bus Enable. When the Bus Enable is LOW, the four drivers are in the high impedance state.

BUS₀, BUS₁, BUS₂. BUS₃ The four driver outputs and receiver inputs (data is inverted).

 \mathbf{R}_0 , \mathbf{R}_1 , \mathbf{R}_2 , \mathbf{R}_3 . The four receiver outputs. Data from the bus is inverted while data from the A or B inputs is non-inverted.

RLE Receiver Latch Enable. When RLE is LOW, data on the BUS inputs is passed through the receiver latches. When RLE is HIGH, the receiver latches are closed and will retain the data independent of all other inputs.

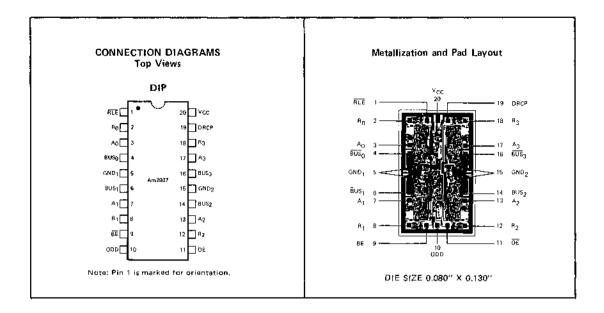
ODD Odd parity output. Generates parity with the driver enabled, checks parity with the driver in the high-impedance state.

OE Output Enable. When the \overrightarrow{OE} input is HIGH, the four three-state receiver outputs are in the high-impedance state.

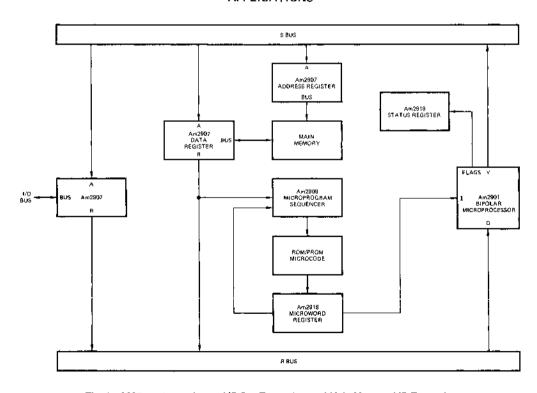
PARITY OUTPUT FUNCTION TABLE

BE	ODD PARITY OUTPUT					
L	ODD = A ₀ + A ₁ + A ₂ + A ₃					
н,	$ODD = \mathbf{Q}_0 \oplus \mathbf{Q}_1 \oplus \mathbf{Q}_2 \oplus \mathbf{Q}_3$					





APPLICATIONS



The Am2907 can be used as an I/O Bus Transceiver and Main Memory I/O Transceiver in high-speed Microprocessor Systems.



GENERAL DESCRIPTION

The Am2909 is a four-bit wide address controller intended for sequencing through a series of microinstructions contained in a ROM or PROM. Two Am2909's may be interconnected to generate an eight-bit address (256 words), and three may be used to generate a twelve-bit address (4K words).

The Am2909 can select an address from any of four sources. They are: 1) a set of external direct inputs (D); 2) external data from the R inputs, stored in an internal register; 3) a four-word deep push/pop stack; or 4) a program counter register (which usually contains the last address plus one). The push/pop stack includes certain control lines so that it can efficiently execute nested sub-routine linkages. Each of the four outputs can be OR'ed with an external input for conditional skip or branch instructions, and a separate line forces the outputs to all zeroes. The outputs are three-state.

The Am2911 is an identical circuit to the Am2909, except the four OR inputs are removed and the D and R inputs are tied together. The Am2911 is in a 20-pin, 0.3" centers package.

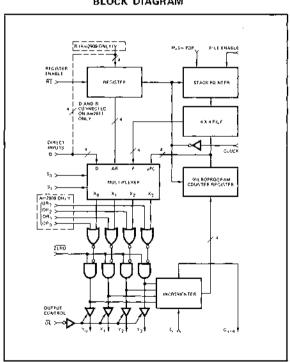
ORDERING INFORMATION

Package Type	Temperature Range	Am2909 Order Number	Am2911 Order Number
Molded DIP	0°C to +70°C	AM2909PC	AM2911PC
Hermetic DIP	0°C to +70°C	AM2909DC	AM2911DC
Hermetic DIP	-55°C to +125°C	AM2909DM	AM2911DM
Hermetic Flat Pa	k -55°C to +125°C	Am2909FM	-
Dice	0°C to +70°C	Am2909XC	_

DISTINCTIVE CHARACTERISTICS

- 4-bit slice cascadable to any number of microwords
- Internal address register
- Branch input for N-way branches
- Cascadable 4-bit microprogram counter
- 4 x 4 file with stack pointer and push pop control for nesting microsubroutines.
- · Zero input for returning to the zero microcode word
- Individual OR input for each bit for branching to higher microinstructions (Am2909 only).
- Three-state outputs
- All internal registers change state on the LOW-to-HIGH transition of the clock
- Am2909 in 28-pin package
- Am2911 in 20-pin package

MICROPROGRAM SEQUENCER BLOCK DIAGRAM



MAXIMUM RATINGS (Above which the useful life may be impaired)

Storage Temperature	65°C to +150°C
Temperature (Ambient) Under Bias	-55°C to +125°C
Supply Voltage to Ground Potential	-0.5 V to +7.0 V
DC Voltage Applied to Outputs for HIGH Output State	-0.5 V to +V _{CC} max.
DC Input Voltage	-0.5 V to +7.0 V
DC Output Current, Into Outputs	30 mA
DC Input Current	-30 mA to +5.0 mA

OPERATING RANGE

P/N	Ambient Temperature	Vcc
Am2909/2911DC, PC	0°C to +70°C	4.75 V to 5.25 V
Am2909/2911DM, FM	-55°C to +125°C	4,50 V to 5,50 V

STANDARD SCREENING

(Conforms to MIL-STD-883 for Class C Parts)

Step	MIL-STD-883 Method	Conditions	Am2909/Am2911DM,FM
Pre-Seal Visual Inspection	2010	В	100%
Stabilization Bake	1008	C 24-hour 150° C	100%
Temperature Cycle	1010	C -65° C to +150° C 10 cycles	100%
Centrifuge	2001	B10,000 G	100%
Fine Leak	1014	A 5 x 10 ⁸ atm-cc/cm ³	100%
Gross Leak	1014	C2 Fluorocarbon	100%
Electrical Test		See below for	100%
Subgroups 1 and 7	5004	definitions of subgroups	
Insert Additional Screening h	ere for Class B Pa	rts	
Group A Sample Tests			
Subgroup 1	1		LTPD = 5
Subgroup 2	!		LTPD = 7
Subgroup 3	1	See below for	LTPD = 7
Subgroup 7	5005	definitions of subgroups	LTPD = 7
Subgroup 8			LTPD = 7
Subgroup 9	!		LTPD = 7

ADDITIONAL SCREENING FOR CLASS B PARTS

MIL-STO-883			Level
Step	Method	Conditions	Am2909/Am2911DMB, FM6
Burn-In	1015	D 125°C 160 hours	
Electrical Test	5004		
Subgroup 1			100%
Subgroup 2		į.	100%
Subgroup 3		!	100%
Subgroup 7			100%
Subgroup 9			100%

GROUP A SUBGROUPS

(as defined in MiL-STD-883, method 5005)

Subgroup	Parameter	Temperature
1	DC	25°C
2	DC	Maximum rated temperature
3	DC	Minimum rated temperature
7	Function	25°C
8	Function	Maximum and minimum rated temperature
9	Switching	25°C
10	Switching	Maximum Rated Temeperature
11	Switching	Minimum Rated Temperature



ELECTRICAL CHARACTERISTICS OVER OPERATING RANGE (Unless Otherwise Noted)

Parameters	Description	Test Conditions (Note 1)				Min.	(Note 2)	Max.	Units
V		VCC = MIN.,	MIL	¹0H =	10H = -1.0mA				
voн	Output HIGH Voltage	VIN = VIH or VIL	COM'L	Тон =	-2.6mA	2.4			Volts
1			I _{OL} = 4.0	0mA				0.4	
	Overse LOW Value	V _{CC} = MIN.,	I _{OL} = 8.0	0mA				0.45	
VOL	Output LOW Voltage	VIN = VIH or VIL	I _{OL} = 12mA (Note 5)				0.5	Volts	
V _{IH}	Input HIGH Level	Guaranteed input for voltage for all inputs	*			2.0			Volt:
.,		Guaranteed input logical LOW MIL voltage for all inputs COM'L				0.7	Volts		
VIL	Input LOW Level					0.8			
Vi	Input Clamp Voltage	VCC " MIN., IIN =	= -18mA					-1.5	Volt
1 _{IL}	Input LOW Current	1 M = = MAX	Cn				-1.08		
		V _{CC} = MAX., V _{IN} = 0.4 V	Push/Pop, OE				0.72	mA	
				Others (Note 6)					-0.36
	· · · · · · · · · · · · · · · · · · ·		Cn					40	
ч	Input HIGH Current	V _{CC} = MAX., V _{IN} = 2.7 V	Push/Pop			·		40	μA
		V [4] - 2,7 4	Others (Note 6)				20		
	1	V _{CC} = MAX., C _n , Push/Pop				0.2			
11	Input HIGH Current	V _{IN} = 7,0 V	Others (N	Note 6)			1	0.1	mΑ
IOS	Output Short Circuit Current (Note 3)	V _{CC} = MAX.			- 40		-100	mA	
lcc	Power Supply Current	V _{CC} = MAX. (Note	lote 4)				80	130	mA
tozL	2	V _{CC} = MAX.,	Vou⊤=	0,4∨		•		-20	
fozh	Output OFF Current	ŌĒ = 2,7 V	VOUT =	2.7 V				20	μА

Notes: 1. For conditions shown as MIN, or MAX,, use the appropriate value specified under Electrical Characteristics for the applicable device type.

Typical limits are at V_{CC} = 5.0V, 25° C ambient and maximum loading.
 Not more than one output should be shorted at a time. Duration of the short circuit test should not exceed one second.

Apply GND to C₀, R₀, R₁, R₂, R₃, OR₀, OR₁, OR₂, OR₃, D₀, D₁, D₂, and D₃. Other inputs open, All outputs open, Measured after a LOW-to-HIGH clock transition.

The 12mA guarantee applies only to Y₀, Y₁, Y₂ and Y₃.
 For the Am2911, D₁ and B₁ are internally connected. Loading is doubled (to same values as Push/Pop).



SWITCHING CHARACTERISTICS OVER OPERATING RANGE

All parameters are guaranteed worst case over the operating voltage and temperature range for the device type, (Grade C = 0° C to +70°C, 4.75V to 5.25V; Grade M = -55°C to +125°C, 4.5V to 5.5V)

TABLE I MINIMUM CLOCK REQUIREMENTS

Minimum Clock LOW Time	50
Minimum Clock HIGH Time	30

TABLE II
MAXIMUM COMBINATORIAL
PROPAGATION DELAYS

OUTPUTS	Yį	C _{n+4}
Ō€	25	-
ŽERO	35	45
on _i	20	32
s ₀ , s ₁	40	50
Đį	20	32
c _n	_	18

TABLE III
MAXIMUM DELAYS
FROM CLOCK TO OUTPUTS

FUNCTIONAL PATH	GRADE	CLOCK TO Y _i	CLOCK TO C _{n+4}
Register	С	48	58
(S ₁ S ₀ = LH)	М	55	65
μ Program Counter	С	48	58
(S ₁ S ₀ = LL)	М	55	65
File	С	70	80
(S1 S0 = HL)	IV)	80	90

R_L = 2.0kΩ C_L ≤ 15pF

TABLE IV SET-UP AND HOLD TIME REQUIREMENTS

EXTERNAL INPUTS	. t _{\$}	th
FIE	20	5.0
Ri	15	0
PU\$H/POP	20	5.0
FE	20	0
Cn	15	0
D;	20	0
QR _i	20	0
S ₀ , S ₁	40	0
ZERO	40	0

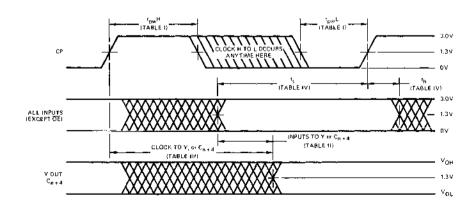


Figure 1. Switching Waveforms. See Tables for Specific Values.



DEFINITIONS

A set of symbols is used in this data sheet to represent various internal and external registers and signals used with the Am2909. Since its principle application is as a controller

for a microprogram store, it is necessary to define some signals associated with the microcode itself. Figure 3 illustrates the basic interconnection of Am2909, memory, and microinstruction register. The definitions here apply to this architecture.

Inputs to Am2909/Am2911

S ₁ , S ₀	Control lines for address source selection
FE, PUP	Control lines for push/pop stack
ŘÉ	Enable line for internal address register
ORi	Logic OR inputs on each address output line
ZERO	Logic AND input on the output lines
ŌĒ	Output Enable, When OE is HIGH, the Y out-
	puts are OFF (high impedance)
Cn	Carry-in to the incrementer
Ri	Inputs to the internal address register
Di	Direct inputs to the multiplexer
CP	Clock input to the AR and µPC register and
	Push-Pop stack

Internal Signals

μPC	Contents	of	the	microprogram	counter

REG Contents of the register

STKO-STK3

Contents of the push/pop stack. By definition, the word in the four-by-four file, addressed by the stack pointer is STKO. Conceptually data is pushed into the stack at STKO; a subsequent push moves STKO to STK1; a pop implies STK3 \rightarrow STK2 \rightarrow STK1 \rightarrow STKO. Physically, only the stack pointer changes when a push or pop is performed. The data does not move. I/O

occurs at STK0.

SP Contents of the stack pointer

External to the Am2909/Am2911

A Address to the control memory

I(A) Instruction in control memory at address A

μWR Contents of the microword register (at output of control memory). The microword register contains the instruction currently being exe-

cuted.

T_n Time period (cycle) n

Outputs from the Am2909/Am2911

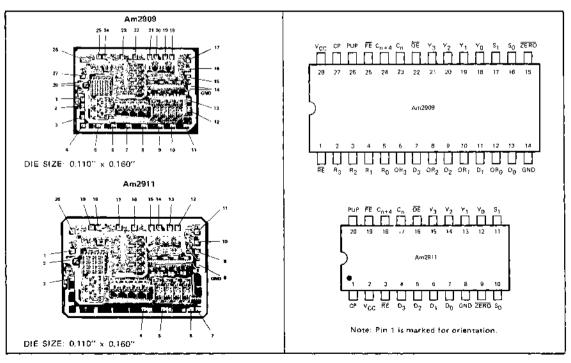
Yi Address outputs from Am2909. (Address inputs

to control memory.)

Cn+4 Carry out from the incrementer

METALLIZATION AND PAD LAYOUT

CONNECTION DIAGRAMS Top Views





ARCHITECTURE OF THE Am2909/Am2911

The Am2909/Am2911 are bipolar microprogram sequencers intended for use in high-speed microprocessor applications. The device is a cascadable 4-bit slice such that two devices allow addressing of up to 256-words of microprogram and three devices allow addressing of up to 4K words of microprogram. A detailed logic diagram is shown in Figure 4.

The device contains a four-input multiplexer that is used to select either the address register, direct inputs, microprogram counter, or file as the source of the next microinstruction address. This multiplexer is controlled by the S₀ and S₁ inputs.

The address register consists of four D-type, edge triggered flip-flops with a common clock enable. When the address register enable is LOW, new data is entered into the register on the clock LOW-to-HIGH transition. The address register is available at the multiplexer as a source for the next micro-instruction address. The direct input is a four-bit field of inputs to the multiplexer and can be selected as the next microinstruction address. On the Am2911, the direct inputs are also used as inputs to the register. This allows an N-way branch where N is any word in the microcode.

The Am2909/Am2911 contains a microprogram counter (μ PC) that is composed of a 4-bit incrementer followed by a 4-bit register. The incrementer has carry-in (C_n) and carry-out (C_{n+4}) such that cascading to larger word lengths is straightforward. The μ PC can be used in either of two ways. When the least significant carry-in to the incrementer is HIGH, the microprogram register is loaded on the next clock cycle with the current Y output word plus one (Y+1 \rightarrow μ PC.) Thus sequential microinstructions can be executed. If this least significant C_n is LOW, the incrementer passes the Y output word unmodified and the microprogram register is loaded with the same Y word on the next clock cycle (Y \rightarrow μ PC). Thus, the same microinstruction can be executed any number of times by using the least significant C_n as the control.

The last source available at the multiplexer input is the 4 x 4 file (stack). The file is used to provide return address linkage

when executing microsubroutines. The file contains a built-in stack pointer (SP) which always points to the last file word written. This allows stack reference operations (looping) to be performed without a push or pop.

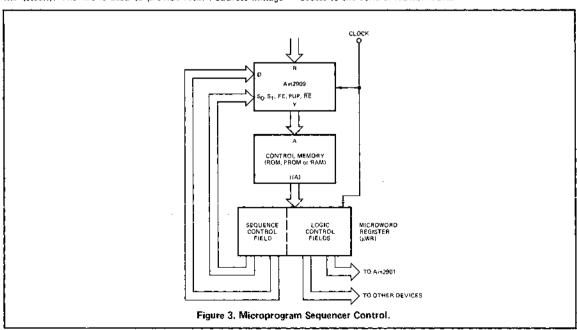
The stack pointer operates as an up/down counter with separate push/pop and file enable inputs. When the file enable input is LOW and the push/pop input is HIGH, the PUSH operation is enabled. This causes the stack pointer to increment and the file to be written with the required return linkage — the next microinstruction address following the subroutine jump which initiated the PUSH.

If the file enable input is LOW and the push/pop control is LOW, a POP operation occurs. This implies the usage of the return linkage during this cycle and thus a return from subroutine. The next LOW-to-HIGH clock transition causes the stack pointer to decrement. If the file enable is HIGH, no action is taken by the stack pointer regardless of any other input.

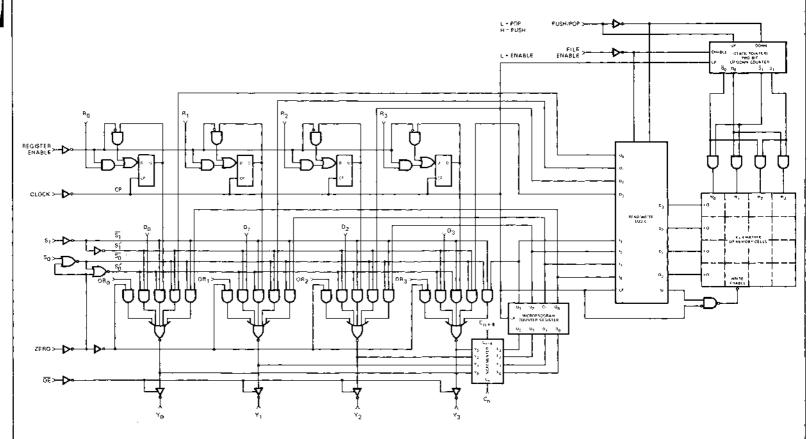
The stack pointer linkage is such that any combination of pushes, pops or stack references can be achieved. One microinstruction subroutines can be performed. Since the stack is 4 words deep, up to four microsubroutines can be nested.

The ZERO input is used to force the four outputs to the binary zero state. When the ZERO input is LOW, all Y outputs are LOW regardless of any other inputs (except \overline{OE}). Each Y output bit also has a separate OR input such that a conditional logic one can be forced at each Y output. This allows jumping to different microinstructions on programmed conditions.

The Am2909/Am2911 feature three-state Y outputs. These can be particularly useful in military designs requiring external Ground Support Equipment (GSE) to provide automatic checkout of the microprocessor. The internal control can be placed in the high-impedance state, and preprogrammed sequences of microinstructions can be executed via external access to the control ROM/PROM.



RAYTHEON



Note: R_i and D_i connected together on Am2911 and OR_i removed.

Figure 4. Microprogram Sequencer Block Diagram.

OPERATION OF THE Am2909/Am2911

H = High L = Low X = Don't Care

Figure 5 lists the select codes for the multiplexer. The two bits applied from the microword register (and additional combinational logic for branching) determine which data source contains the address for the next microinstruction. The contents of the selected source will appear on the Y outputs. Figure 5 also shows the truth table for the output control and

for the control of the push/pop stack. Figure 6 shows in detail the effect of S_0 , S_1 , \overline{FE} and PUP on the Am2909. These four signals define what address appears on the Y outputs and what the state of all the internal registers will be following the clock LOW-to-HIGH edge. In this illustration, the microprogram counter is assumed to contain initially some word J, the address register some word K, and the four words in the push/pop stack contain R_a through R_d .

Address Selection

OCTAL	S ₁	S ₀	SOURCE FOR Y OUTPUTS	SYMBOL
0	L	ЯT	Microprogram Counter Register	μPC REG
2 3	H	L H	Push-Pop stack Direct inputs	STK0 Di

Output Control

ORi	ZERO	ŌĒ	Yi
×	×	н	Z
X	L	L	L
H	H	L	н
L	Н	L	Source selected by \$0 \$1

Z - High Impedance

Synchronous Stack Control

FE	PUP	PUSH-POP STACK CHANGE
н	х	No change
L	н	Increment stack pointer, then push current PC onto STK0
L	L	Pop stack (decrement stack pointer)

Figure 5.

CYCLE	S ₁ , S ₀ , FE, PUP	μРС	REG	sтко	STK1	STK2	STK3	Yout	COMMENT	PRINCIPLE USE
N N+1	0000	J+1	K K	Ra Rb	Rb Rc	Rc Rd	Rd Ra	7	Pop Stack	End Loop
N N+1	0001	J J+1	K	Ra J	Rb Ra	Rc Rb	Rd Rc	–	Push μPC	Set-up Loop
N N+1	0 0 1 X -	J J+1	K K	Ra Ra	Rb Rb	Rc Rc	Rd Rd	_ 1	Continue	Continue
N N+1	0100	ј К+1	K K	Ra Rb	Rb Rc	Rc Rd	Rd Ra	к -	Pop Stack; Use AR for Address	End Loop
N N+1	0101	J K+1	K K	Ra J	Rb Ra	Rc Rb	Rd Rc	K 	Push μPC; Jump to Address in AR	J\$R AR
N N+1	0 1 1 X	J K+1	K	. Ra Ra	Rb Rb	Rc Rc	Rd Rd	K -	Jump to Address in AR	JMP AR
N N+1	1000	J Ra+1	K K	Ra Rb	Rb Rc	Rc Rd	Rd Ra	Ra -	Jump to Address in STK0; Pop Stack	RTS
N N+1	1001	J Ra+î	K K	Ra J	Rb Ra	Rc Rb	Rd Rc	Ra —	Jump to Address in STK0; Push μPC	
N N+1	1 0 1 X	J Ra+1	K	Ra Ra	Rb Rb	Rc Rc	Rd Rd	Ra 	Jump to Address in STK0	Stack Ref (Loop)
N N+1	1 1 0 0	J D+1	K	Ra Rb	Rb Rc	Rc Rd	Rd Ra	D -	Pop Stack; Jump to Address on D	End Loop
N N+1	1101	J D+1	K K	Ra J	Rb Ra	Rc Rb	Rd Rc	D -	Jump to Address on D; Push µPC	JSR D
N N+1	1 1 1 X -	J D+1	K	Ra Ra	Rb Rb	Rc Rc	Rd Rd	D 	Jump to Address on D	JMP D

 $X = Don't care, 0 = LOW, 1 = H(GH, Assume C_n = H(GH, Note; STKO is the location addressed by the stack pointer.$

Figure 6. Output and Internal Next-Cycle Register States for Am2909/Am2911.



Figure 7 illustrates the execution of a subroutine using the Am2909. The configuration of Figure 3 is assumed. The instruction being executed at any given time is the one contained in the microword register (μ WR). The contents of the μ WR also controls (indirectly, perhaps) the four signals S_0 , S_1 , \overline{FE} , and PUP. The starting address of the subroutine is applied to the D inputs of the Am2909 at the appropriate time.

In the columns on the left is the sequence of microinstructions to be executed. At address J+2, the sequence control portion of the microinstruction contains the comand "Jump to sub-

routine at A". At the time T_2 , this instruction is in the μWR , and the Am2909 inputs are set-up to execute the jump and save the return address. The subroutine address A is applied to the D inputs from the μWR and appears on the Y outputs. The first instruction of the subroutine, I(A), is accessed and is at the inputs of the μWR . On the next clock transition, I(A) is loaded into the μWR for execution, and the return address J+3 is pushed onto the stack. The return instruction is executed at T_5 . Figure 8 is a similar timing chart showing one subroutine linking to a second, the latter consisting of only one microinstruction.

CONTROL MEMORY

Éxecute	Micro	program
Cycle	Address	Sequencer Instruction
	J-1	_
T_0)'	_
τ ₁ ,	J+1	_
T ₂	J+2	JSR A
T ₆	J+3	
т ₇	J+4	<u> </u>
.,	_	_
		_
	_	_
	~	_
	-	_
T_3	A	I(A)
T ₄	A+1	_
T ₅	A+2	RTS
•	_	_
	_	_
		_
	-	_
	-	-
	-	_
		-
	I	1

Execute (Cycle .	T _G	T ₁	T ₂	T ₃	T ₄	T 5	T ₆		_ т ₈	Τg
Signa	lock ls										
Am2909 Inputs (from µWR)	S ₁ , S ₀ FE PUP D	0 H X	O H X	3 L H A	0 H X X	о н х х	2 L L X	0 H X X	0 H X X		
Internal Registers	μPC STK0 STK1 STK2 STK3	J+1 - - -	J+2 - - -	J+3 - -	A+1 J+3 — —	A+2 J+3 - -	A+3 J+3 - -	J+4 - - -	J+5		_
Am2909 Output	Y	J+1	J+2	A	A+1	A+2	J+3	J+4	J+5		_
ROM Output	(Y)	1(1)+11	JSR A	I(A)	I(A+1)	RTS	1(J+3)	1(J+4)	!(J+5)		
Contents of µWR (Instruction being executed)	RWμ	l (J)	I(J +1)	JSR A	I(A)	F(A+1)	RTS	1(J+3)	I(J+4)		

Figure 7. Subroutine Execution.

C_n ≈ HIGH

CONTROL MEMORY

	I
Address	Sequencer Instruction
J-1	_
J	_ !
J+1	_
J+2	JSR A
J+3	- 1
	_ (
_	_ 1
_	
_	-
А	- 1
A+1	-
A+2	JSR B
A+3	_
A+4	RTS
_	_
_	-
_	-
	-
₿	RTS
_	-
_	1 -
	J-1 J+1 J+2 J+3 A+1 A+2 A+3 A+4

Execute C	ycle	To	T ₁	T ₂	T ₃	T ₄	T 5	T ₆	T ₇	T ₈	Tg
C Signa	lock Is										
Am2909 Inputs (from μWR)	S ₁ , S ₀ FE PUP D	о н х х	0 H X X	3 L H A	0 H X	0 H X X	3 L H B	2 L L X	0 H X	2 L L	о н х х
Internal Registers	μPC STK0 STK1 STK2 STK3	J+1 - - -	J+2 - - -	J+3 - -	A+1 J+3 - -	A+2 J+3 -	A+3 J+3 	B+1 A+3 J+3 	A+4 J+3 - -	A+5 J+3 - -	J+4 - - - -
Am2909 Output	Y	J+1	J+2	А	A+1	A+2	8	A+3	A+4	J+3	J+4
ROM Output	(Y)	I(J+1)	JSR A	HA	I(A+1)	J\$R B	RTS	I(A+3)	RTS	({J+3)	F(J+4)
Contents of µWR (Instruction being executed)	μWR	1(1)1	I(J+1)	J\$R A	1(A)	I(A+1)	JSR B	RTS	1(E+A)1	RTS	1(J+3)

Figure 8. Two Nested Subroutines, Routine B is Only One Instruction.

Co = HIGH



FUNCTIONAL DESCRIPTION

The Am2915 is a high-performance, low-power Schottky bus transceiver intended for bipolar or MOS microprocessor system applications. The device consists of four D-type edge-triggered flip-flops with a built-in two-input multiplexer on each. The flip-flop outputs are connected to four three-state bus drivers. Each bus driver is internally connected to the input of a receiver. The four receiver outputs drive four D-type latches that feature three-state outputs.

This LSI bus transceiver is fabricated using advanced low-power Schottky processing. All inputs (except the BUS inputs) are one LS unit load. The three-state bus output can sink up to 40mA at 0.5V maximum. The bus enable input (\overline{BE}) is used to force the driver outputs to the high-impedance state. When \overline{BE} is HIGH, the driver is disabled. The VOH and VOL of the bus driver are selected for compatibility with standard and Low-Power Schottky inputs.

The input register consists of four D-type flip-flops with a buffered common clock and a two-input multiplexer at the input of each flip-flop. A common select input (5) controls the four multiplexers. When S is LOW, the A_i data is stored in the register and when S is HIGH, the B_i data is stored. The buffered common clock (DRCP) enters the data into this driver register on the LOW-to-HIGH transition.

Data from the A or B inputs is inverted at the BUS output. Likewise, data at the BUS input is inverted at the receiver output. Thus, data is non-inverted from driver input to receiver output. The four receivers each feature a built-in D-type latch that is controlled from the buffered receiver latch enable (RLE) input. When the RLE input is LOW, the latch is open and the receiver outputs will follow the bus inputs (BUS data

inverted and OE LOW). When the \overline{RLE} input is HIGH, the latch will close and retain the present data regardless of the bus input. The four latches have three-state outputs and are controlled by a buffered common three-state control (\overline{OE}) input. When \overline{OE} is HIGH, the receiver outputs are in the high-impedance state,

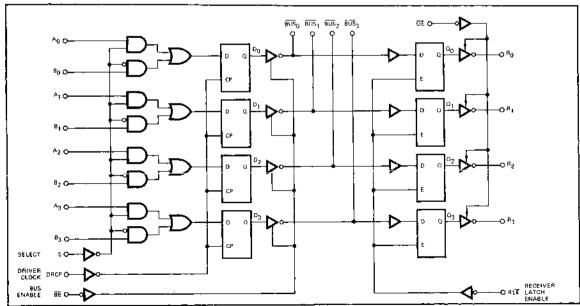
DISTINCTIVE CHARACTERISTICS

- Quad high-speed LSI bus-transceiver
- Three-state bus driver
- Two-port input to D-type register on driver
- Bus driver output can sink 40mA at 0.5V max.
- Receiver has output latch for pipeline operation.
- Three-state receiver outputs sink 12mA
- Advanced low-power Schottky processing
- 100% reliability assurance testing in compliance with MIL-STD-883
- 3.5V minimum output high voltage for direct interface to MOS microprocessors

ORDERING INFORMATION

Package Type	Temperature Range	Order Number
Molded DIP	0°C to +70°C	AM2915PC
Hermetic DIP	0°C to +70°C	AM2915DC
Dice	0°C to +70°C	AM2915XC
Hermetic DIP	55°C to +125°C	AM2915DM
Hermetic Flat Pak	-55°C to +125°C	AM2915FM
Dice	-√55°C to +125°C	AM2915XM

LOGIC DIAGRAM



MAXIMUM RATINGS (Above which the useful life may be impaired)

Storage Temperature	-65°C to +150°C
Temperature (Ambient) Under Bias	-55°C to +125°C
Supply Voltage to Ground Potential	-0.5V to +7V
DC Voltage Applied to Outputs for HIGH Output State	-0.5V to +V _{CC} max.
DC Input Voltage	-0.5V to +5,5V
DC Output Current, Into Outputs (Except Bus)	30mA
DC Output Current, Into Bus	100mA
DC Input Current	-30mA to +5.0mA

ELECTRICAL CHARACTERISTICS

The following conditions apply unless otherwise noted:

BUS INPUT/OUTPUT CHARACTERISTICS OVER OPERATING TEMPERATURE RANGE

Parameters	Description	Test Conditions	Min.	Typ.	Max.	Units	
V	Bus Output LOW Voltage	V _{CC} = MIN.	I _{OL} = 24mA			0.4	Volts
V _{OL}	Bus Gotput LOW Variage	ACC MILIA.	IOL = 40mA			0.5	A COLUZ
VOH	Bus Output HIGH Voltage	V _{CC} = MIN.	I _{OH} = -20mA	2.4			Valts
	Bus Leakage Current		VO = 0.4 V			-200	
i to	(High Impedance)	V _{CC} = MAX. Bus enable = 2.4 V	V _O = 2.4 V			50	μΑ
Í	,	D43 6140016 - 2.4 V	V _O = 4.5 V			100	!
IOFF	Bus Leakage Current	V _O = 4.5 V	•			100	,
"0	(Power OFF)	VCC = 0 V				100	μΑ
V _{IH}	Receiver Input HIGH Threshold	Bus enable = 2.4 V		2.0			Volts
VIL	Receiver Input LOW Threshold	Bus enable = 2.4 V				0.8	Volts
Isc	Bus Output Short Circuit Current	V _{CC} = MAX. V _O = 0 V		-50	 -85	130	mA



ELECTRICAL CHARACTERISTICS

The following conditions apply unless otherwise noted:

Am2915XC(COM'L) $T_A = 0^{\circ}C$ to +70°C $V_{CC}MIN. = 4.75V$ $V_{CC}MAX. = 5.25V$

Am2915XM(MIL) TA = -55°C to +125°C VCCMIN. = 4.50 V VCCMAX. = 5.50 V

rameters	Description	Test Cond	itions (r	lote 1)	Min.	Typ. (Note 2)	Max.	Units
		V _{CC} = MIN,	MIL:	OH = -1.0mA	2.4	3.4		
V _{OH}	Receiver Output HIGH Voltage	V _{1N} = V _{1L} or V _{1H}	COM,	L: I _{OH} = -2.6mA	2.4	3.4		Valts
		VCC = 5.0 V, LOH = -	100μΑ		3.5			L
	0	V _{CC} = MIN.		I _{OL} = 4.0mA		0.27	0.4	
VOL ,	Output LOW Voltage (Except Bus)	VIN = VIL or VIH		I _{QL} = 8.0mA		0.32	0.45	Volts
	1-1000	IOL = 12mA				0.37	0.5	
VIH	Input HIGH Level (Except Bus)	Guaranteed input logical HIGH for all inputs			2.0			Volts
VIL	Input LOW Level (Except Bus)	Guaranteed input logical LOW for all inputs					0.8	Volt
V _I	Input Clamp Voltage (Except Bus)	VCC = MIN., 11N = -1	8mA				-1.2	Volt
1	Input LOW Current	V MAY V (.4.17	BE, RLE		Ţ	-0.72	
Կ <u>L</u>	Input COM Current	V _{CC} = MAX., V _{IN} = 0.4V All other inputs					0.36	mA
Чн .	Input HIGH Current (Except Bus)	V _{CC} = MAX., V _{IN} = 2	.7 V				20	μА
ij.,	Input HIGH Current (Except Bus)	VCC = MAX., VIN = 7	.0 V				100	μА
I _{SC}	Output Short Circuit Current (Except Bus)	V _{CC} = MAX.		-30		-85	mA	
Icc	Power Supply Current	V _{CC} = MAX.				60	90	mΑ
10	Off-State Output Current	$V_{CC} = MAX$. $V_{O} = 2.4 V$ $V_{O} = 0.4 V$		V _O = 2.4 V			20	
.0	(Receiver Outputs)				ļ	-20	μА	

SWITCHING CHARACTERISTICS OVER OPERATING TEMPERATURE RANGE

			,	4m2915X	М	,	4m2915X(:		
Parameters	Description	Test Conditions	Min.	Typ. (Note 2)	Max.	Min.	Typ. (Note 2)	Max.	Units	
tPHL :	Driver Clock (DRCP) to Bus			21	36		21	32		
tPLH .	Driver Clock (DRCP) to Bus	C_L (BUS) = $50pF$		21	36		21	32	ns	
tZH, tZL		R _L (BUS) = 130Ω		13	26		13	23		
tHZ, tLZ	Bus Enable (BE) to Bus	X .		13	26		13	23	ns	
t _s	~~~~	77.	23	ļ		50			,	
th	Data Inputs (A or B)		8.0	1		6.0			ПŚ	
ts	3	SO/2 \ 1	28	ļ		25	T	•		
th	Select Input (\$)	No.Ca	8.0			6.0			ns	
tpW	Driver Clock (DRCP) Pulse Width (HIGH)	R _L (BUS) = 130Ω R _L (BUS) = 130Ω R _L (BUS) = 130Ω R _L (BUS) = 130Ω	20			17			ns	
tpLH	Bus to Receiver Output	JOY AR	V	18	30		18	27		
tPHL .	(Latch Enable)	CL = 15pAO &		18	30		18	27	ns	
tPLH (Latch Enable to Receiver Output	R _L = 2.0kΩ / C	160	21	30		21	27	пs	
tPHL ;	Later Enable to Receiver Output	•	7	21	30		21	27	115	
ts			17	1		14				
th	Bus to Latch Enable (RLE)		6.0	Ţ <u> </u>		4.0	T		ПŠ	
tZH, tZL	00			14	56		14	23		
tHZ, tLZ	Output Control to Receiver Output			14	26		14	23	ns	

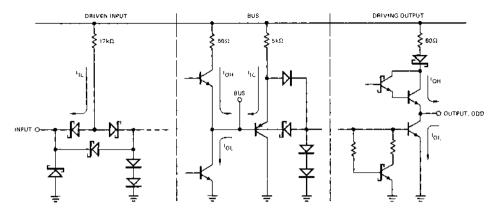
Notes: 1. For conditions shown as MIN, or MAX., use the appropriate value specified under Electrical Characteristics for the applicable device type.

2. Typical limits are at V_{CC} = 5.0 V, 25 °C ambient and maximum loading.

3. Not more than one output should be shorted at a time. Duration of the short circuit test should not exceed one second.

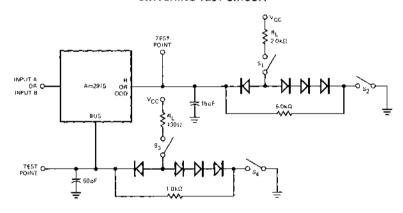


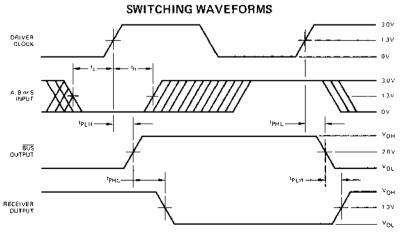
INPUT/OUTPUT CURRENT INTERFACE CONDITIONS



Note: Actual current flow direction shown.

SWITCHING TEST CIRCUIT





Note: Bus to Reciver output delay is measured by clocking data into the driver register and measuring the BUS to R combinatorial delay.



FUNCTIONAL TABLE

			INPUT	s			INTE TO D	NAL	BU\$	оитрит	FUNCTION
s	Αį	Bi	DRCP	BE	RLE	ŌĒ	Dj	α_{i}	BUSi	Ri	
Х	X	Х	Х	Н	Х	Х	X	Х	Z	Х	Driver output disable
Х	Х	Х	Х	×	X	Н	Х	Х	Х	Z	Receiver output disable
Х	Х	x	Х	Н	٦	L	х	Ļ	L	H	Driver output disable and receive data
Х	х	х	x	н	L	L	×	н	Н	L	via Bus input
Х	Х	Х	Х	X	Н	X	Х	NC	Х	Х	Latch received data
L	L	Х	1	×	Х	X	L	X	Х	X	
L	н	х	†	×	×	×	н	x	×	x	Load driver register
н	X	L	↑ ↑	х	×	х	Ł	Х	Х	x	Load di iver register
Н	х	Н	1	х	X	х	Н	×	×	×	
Х	Х	х	L	Х	×	Х	NC	Х	Х	х	No driver clock restrictions
x	х	х	H	х	×	х	NC	x	х	×	140 dilyer clock restrictions
Х	Х	х	Х	L	Х	Х	L	х	Н	х	Data Bara
X	х	×	X	L	×	х	Н	Х	L	×	Drive Bus

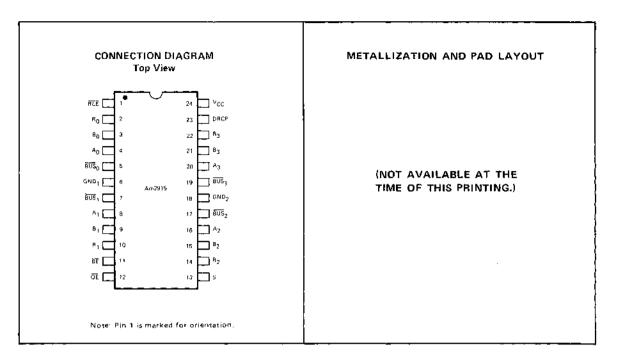
H = HIGH Z = HIGH Impedance X = Don't Care
L = LOW NC = No Change f = LOW-to-HIGH Transition

i = 0, 1, 2, 3

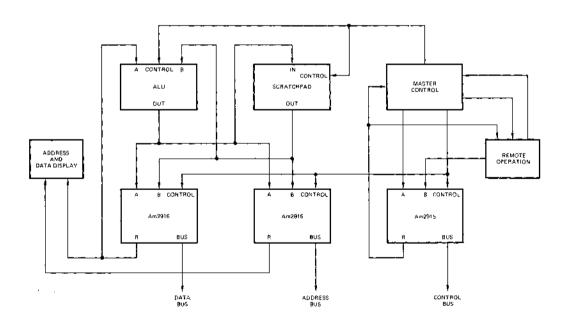
DEFINITIONS

A ₀ , A ₁ , A ₂ , A ₃	The "A" word data input into the two input multiplexer of the driver register.	$\overline{\mathtt{BUS}}_0$, $\overline{\mathtt{BUS}}_1$ $\overline{\mathtt{BUS}}_2$, $\overline{\mathtt{BUS}}_3$	The four driver outputs and receiver inputs (data is inverted).
B ₀ , B ₁ , B ₂ , B ₃	The "B" word data input into the two input multiplexers of the driver register,	R ₀ , R ₁ , R ₂ , R ₃	The four receiver outputs. Data from the bus is inverted while data from the A or B inputs is non-inverted.
S DRCP	Select. When the select input is LOW, the A data word is applied to the driver register. When the select input is HIGH, the B word is applied to the driver register. Driver Clock Pulse, Clock pulse for the	RLÉ	Receiver Latch Enable. When RLE is LOW, data on the BUS inputs is passed through the receiver latches. When RLE is HIGH, the receiver latches are closed and will retain the data independent of
21.00	driver register.		alf other inputs.
BE	Bus Enable. When the Bus Enable is HIGH, the four drivers are in the high impedance state.	ŌĒ	Output Enable. When the \overline{OE} input is HIGH, the four three state receiver outputs are in the high-impedance state,





APPLICATIONS



The Am2915 is a universal Bus Transceiver useful for many system data, address, control and timing input/output interfaces.



FUNCTIONAL DESCRIPTION

The Am2916 is a high-performance, low-power Schottky bus transceiver intended for bipolar or MOS microprocessor system applications. The device consists of four D-type edge-triggered flip-flops with a built-in two-input multiplexer on each. The flip-flop outputs are connected to four three-state bus drivers. Each bus driver is internally connected to the input of a receiver. The four receiver outputs drive four D-type latches. The device also contains a four-bit odd parity checker/generator.

The LSI bus transceiver is fabricated using advanced low-power Schottky processing. All inputs (except the BUS inputs) are one LS unit load. The three-state bus output can sink up to 40mA at 0.5V maximum. The bus enable input (\overline{BE}) is used to force the driver outputs to the high-impedance state. When \overline{BE} is HIGH, the driver is disabled.

The input register consists of four D-type flip-flops with a buffered common clock and a two-input multiplexer at the input of each flip-flop. A common select input (S) controls the four multiplexers. When S is LOW, the A₁ data is stored in the register and when S is HIGH, the B₁ data is stored. The buffered common clock (DRCP) enters the data into this driver register on the LOW-to-HIGH transition.

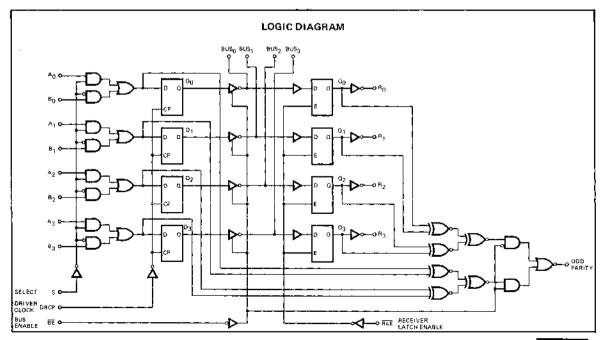
Data from the A or B input is inverted at the BUS output, Likewise, data at the BUS input is inverted at the receiver output. Thus, data is non-inverted from driver input to receiver output. The four receivers each feature a built-in D-type latch that is controlled from the buffered receiver latch enable (RLE) input. When the RLE input is LOW, the latch is open and the receiver outputs will follow the bus

inputs (BUS data inverted). When the RLE input is HIGH, the latch will close and retain the present data regardless of the bus input.

The Am2916 features a built-in four-bit odd parity checker/ generator. The bus enable input (BE) controls whether the parity output is in the generate or check mode. When the bus enable is LOW (driver enabled), odd parity is generated based on the A or B field data input to the driver register. When BE is HIGH, the parity output is determined by the four latch outputs of the receiver. Thus, if the driver is enabled, parity is generated and if the driver is in the high-impedance state, the BUS parity is checked.

DISTINCTIVE CHARACTERISTICS

- Quad high-speed LSI bus-transceiver
- Three-state bus driver
- Two-port input to D-type register on driver
- Bus driver output can sink 40mA at 0.5V max.
- Internal odd 4-bit parity checker/generator
- Receiver has output latch for pipeline operation
- Receiver outputs sink 12mA
- Advanced low-power Schottky processing
- 100% reliability assurance testing in compliance with MIL-STD-883
- 3.5V minimum output high voltage for direct interface to MOS microprocessors





ORDERING INFORMATION

Package	Temperature	Order
Type	Range	Number
Molded DIP	0°C to +70°C	AM2916PC
Hermetic DIP	0°C to +70°C	AM2916DC
Dice	0°C to +70°C	AM2916XC
Hermetic DIP	~55°C to +125°C	AM2916DM
Hermetic Flat Pak	-55°C to +125°C	AM2916FM
Dice	-55°C to +125°C	AM2916XM

MAXIMUM RATINGS (Above which the useful life may be impaired)

Storage Temperature	−65°C to +150°C
Temperature (Ambient) Under Bias	-55°C to +125°C
Supply Voltage to Ground Potential	-0.5V to +7V
DC Voltage Applied to Outputs for HIGH Output State	-0.5V to +V _{CC} max.
DC Input Voltage	-0.5V to +5.5V
DC Output Current, Into Outputs (Except Bus)	30mA
DC Output Current, Into Bus	100mA
DC Input Current	-30mA to +5.0mA

ELECTRICAL CHARACTERISTICS

The following conditions apply unless otherwise noted:

Am2916XC(COM'L) TA = 0°C to +70°C VCCMIN. = 4.75V VCCMAX. = 5.25V

Am2916XM(MIL) $T_A = -55^{\circ}C \text{ to } +125^{\circ}C - V_{CC}MIN. = 4.50 V - V_{CC}MAX. = 5.50 V$

BUS INPUT/OUTPUT CHARACTERISTICS OVER OPERATING TEMPERATURE RANGE

Parameters	Description	Test Conditions	Min.	Typ.	Max.	Units	
V	Bus Output LOW Voltage	14 MINI	IOL = 24 mA			0.4	Volts
VOL	Bus Output LOVY Voltage	V _{CC} = MIN.	IDL = 40mA			0.5	VOITS
v_{OH}	Bus Output HIGH Voltage	VCC = MIN.	I _{OH} = -20mA	2.4			Volts
	First solve 8		VO = 0.4 V			-200	
10	Bus Leakage Current (High Impedance)	V _{CC} = MAX. Bus enable = 2.4 V	V _O = 2.4 V			50	μА
		Dus chable - 2.4 ¥	V _O = 4.5 V			100	
OFF	Bus Leakage Current	V _O = 4.5 V			100		
-077	(Power OFF)	V _{CC} = 0 ∨			100	μA	
VIH	Receiver Input HIGH Threshold	Bus enable = 2.4 V	"	2.0			Volts
VIL	Receiver Input LOW Threshold	Bus enable = 2.4 V				8.0	Volts
Isc	Bus Output Short Circuit Current	V _{CC} ≈ MAX. V _O = 0 V		-50	-85	-130	mA



ELECTRICAL CHARACTERISTICS

The following conditions apply unless otherwise noted:

DC CHARACTERISTICS OVER OPERATING TEMPERATURE RANGE

arameters	Description	SOVER OPERATING TEMPERATURE RANGE		Min.	Typ.	Max.	Units	
·	_	VCC * MIN. MIL		_H =1.0mA	2.4	3.4		
v_{OH}	Receiver Output HIGH Voltage	VIN = VIL or VIH	COM'L	OH = -2.6mA	2.4	3.4		Volts
	output man voltage	V _{CC} = 5.0 V, I _{OH} = -	100 μΑ		3.5			1
v _{oh}	Parity	V _{CC} = MIN., I _{OH} = -	660µA	MIL	2.5	3.4		Volts
VOH	Output HIGH Voltage	V _{1N} = V _{1H} or V _{1L}		COM.F	2.7	3.4		Voits
		1 3/ - MINI		IOL = 4.0mA		0.27	0.4	
v_{OL}	Output LOW Voltage (Except Bus)	VIN = VII OF VIH		lo∟ = 8.0mA		0.32	0,45	Valts
	(Example 5 at)			I _{OL} = 12mA		0.37	0.5	<u> </u>
V _{(H}	Input HIGH Level (Except Bus)	Guaranteed input logical HIGH for all inputs		2.0			Volt	
VIL	Input LOW Level (Except Bus)	Guaranteed input logic for all inputs	al LOW				8.0	Volt
νı	Input Clamp Voltage (Except Bus)	V _{CC} = MIN., I _{IN} = -1	8mA	1		}	1.2	Volt
	Level OM Correct Francis Book	W 144 V W - 5		BE, RLE			-0.72	
lir.	Input LOW Current (Except Bus)	VCC = MAX., VIN = 0.4 V All other input:		All other inputs		1	-0.36	mA
IJН	Input HIGH Current (Except Bus)	V _{CC} = MAX., V _{IN} = 2	2.7 V				20	μА
11	Input HIGH Current (Except Bus)	VCC = MAX., VIN = 7	.0 V			<u> </u>	100	μА
I _{SC}	Output Short Circuit Current (Except Bus)	V _{CC} = MAX.			-30		-85	πА
I _{CC}	Power Supply Current	V _{CC} = MAX., All Inpu	its = GND			75	110	mA

SWITCHIN	G CHARACTERISTICS	,	Am2916X	М	Δ				
Parameters	Description	Test Conditions	Min.	Typ. (Note 2)	Max.	Min.	Typ. (Note 2)	Max.	Units
tenL		Cը (BUS) = 50pF		21	36	l	21	32	
tPLH	Driver Clock (DRCP) to Bus	CL (BUS) = 130 Ω RL (BUS) = 130 Ω RL (BUS) = 130 Ω RL (BUS) = 150 P RL (BUS) = 150 P		21	36		21	32	2.0
tzH, tZL	Bus Enable (BE) to Bus		V 70	13	26		13	23	ns
tHZ. tLZ	Bus chable (BE) to Bus		ζÇ,	13	26	1	13	23] "5
ts	Data Inputs (A or B)		8 4.23	ļ .		20			ns
th	Data inputs (A or B)	112. 62	8.0		—	6.0			1 '''
t _s	Colored In core IO	Per An	28			25			пs
th	Select Inputs (S)	The Che City	8.0	Ì		6.0] "" _
tpW	Clock Pulse Width (HIGH)		20	ì		17			ns ns
tp∟H	Bus to Receiver Output	CIC, MI		18	30		18	27	
tPHL	(Latch Enabled)	cci, age		18	30		18	27	ns
tPLH	Last Factors Section Co.	es, Hu.		21	30		21	27	
text	Latch Enable to Receiver Output	q C		21	30		21	27	ns
tş	Secretary Secret	7	17			14			T
th	Bus to Latch Enable (RLE)	CL = 15pF	6.0	1		4.0	-		ns
tPLH	A or B Data to Odd Parity Output	R _L = 2.0kΩ		21	36		21	32	
tPHL	(Driver Enabled)			21	36		21	32	ns.
tpLH	Bus to Odd Parity Output	1		21	36		21	32	ns
tPHL.	(Driver Inhibited, Latch Enabled)	•		21	36		21	32] ''`
tPLH	Latch Enable (RLE) to	1		21	36	[21	32	0.5
t PHL	Odd Parity Output			21	36		21	32	ns

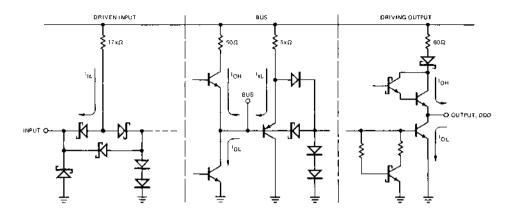
Notes: 1. For conditions shown as MIN, or MAX., use the appropriate value specified under Electrical Characteristics for the applicable device type,

2. Typical limits are at V_{CC} = 5.0 V, 25 C ambient and maximum loading.

3. Not more than one output should be shorted at a time. Duration of the short circuit test shoul not exceed one second.

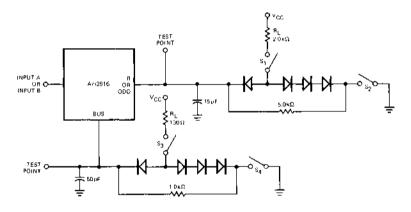


INPUT/OUTPUT CURRENT INTERFACE CONDITIONS

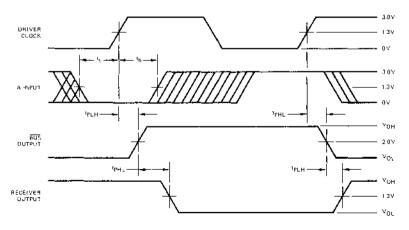


Note: Actual current flow direction shown.

SWITCHING TEST CIRCUIT



SWITCHING WAVEFORMS



Note: Bus to Receiver output delay is measured by clocking data into the driver register and measuring the BUS to R combinatorial delay.



FUNCTION TABLE

			INPUT	s			INTE TO DI	RNAL EVICE	BUS	ООТРИТ	FUNCTION	
\$	Ai	Вį	DRCP	BE	RLE	ŌĔ	Di	αi	<u>BUS</u> i	R;	. 5,15 (101)	
Х	Х	Х	×	н	х	х	X	×	Z	Х	Driver output disable	
х	Х	Х	Х	x	Х	Н	Х	Х	Х	Z	Receiver output disable	
χ	х	х	Х	н	L	Ĺ	Х	L	L	н	Driver output disable and receive data	
X	×	×	×	н	L	L	×	н	н	L L	via 8us input	
Х	Х	х	X	Х	Н	х	_X	NC	х	X	Latch received data	
L	L	х	1	Х	Х	Х	L	Х	X	Х		
L	Н	х	1	×	X	X	н	х	×	×	Load driver register	
Н	х	L	1	_ ×	×	×	L	х	×	×	Load driver register	
Н	х	Н	1	Х	X	×	Н	х	Х	X		
Х	х	х	Ļ	Х	×	х	NC	Х	Х	×	No driver clock restrictions	
Х	х	х	н	×	×	х	NC	×	x	x	IAO GITAGE CLOCK LESTILICIONS	
Х	х	Х	Х	L	х	Х	L	х	Н	Х	Drive Bus	
х	х	х	х	L	×	х	H	×	L	×		

H = HIGH L ≈ LOW

DEFINITIONS

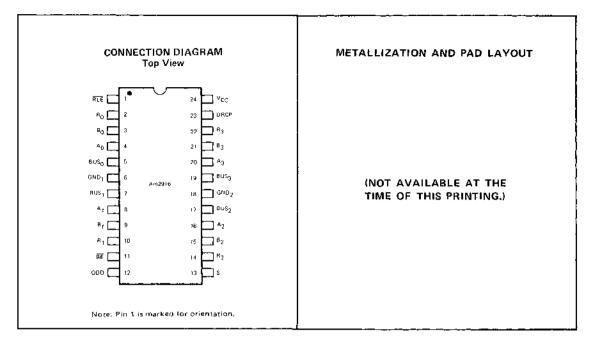
A ₀ , A ₁ , A ₂ , A ₃	The "A" word data input into the two input multiplexer of the driver register.	BUS _O , BUS ₁ BUS ₂ , BUS ₃	The four driver outputs and receiver in- puts (data is inverted).
B ₀ , B ₁ , B ₂ , B ₃	The "B" word data input into the two input multiplexers of the driver register.	R ₀ , R ₁ , R ₂ , R ₃	The four receiver outputs. Data from the bus is inverted while data from the A or B
S	Select. When the select input is LOW, the A data word is applied to the driver register. When the select input is HIGH, the B word is applied to the driver register.	RLE	inputs is non-inverted. Receiver Latch Enable. When RLE is LOW, data on the BUS inputs is passed through the receiver latches. When RLE is HIGH, the receiver latches are closed
DRCP	Driver Clock Pulse. Clock pulse for the driver register.		and will retain the data independent of all other inputs.
BE	Bus Enable. When the Bus Enable is HIGH, the four drivers are in the high impedance state.	O E	Output Enable. When the $\overline{\text{OE}}$ input is HIGH, the four three state receiver outputs are in the high-impedance state.



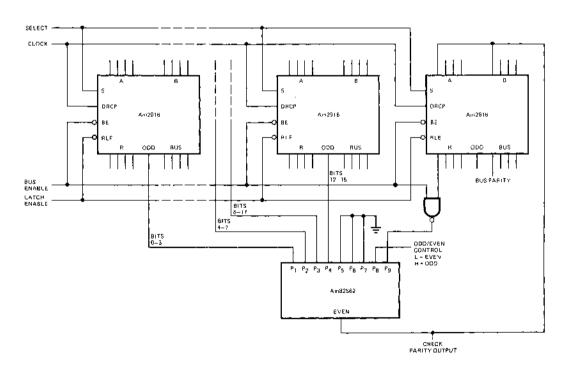
Z = HIGH Impedance NC = No change

X = Don't care
1 = LOW-to-HIGH transition

i = 0, 1, 2, 3



APPLICATIONS



Generating or checking parity for 16 data bits.



FUNCTIONAL DESCRIPTION

The Am2917 is a high-performance, low-power Schottky bus transceiver intended for bipolar or MOS microprocessor system applications. The device consists of four D-type edge-triggered flip-flops. The flip-flop outputs are connected to four three-state bus drivers. Each bus driver is internally connected to the input of a receiver. The four receiver outputs drive four D-type latches, that feature three-state outputs. The device also contains a four-bit odd parity checker/generator.

The LSI bus transceiver is fabricated using advanced low-power Schottky processing. All inputs (except the BUS inputs) are one LS unit load. The three-state bus output can sink up to 40mA at 0.5V maximum. The bus enable input (\overline{BE}) is used to force the driver outputs to the high-impedance state. When \overline{BE} is HIGH, the driver is disabled.

The input register consists of four D-type flip-flops with a buffered common clock. The buffered common clock (DRCP) enters the A_i data into this driver register on the LOW-to-HIGH transition.

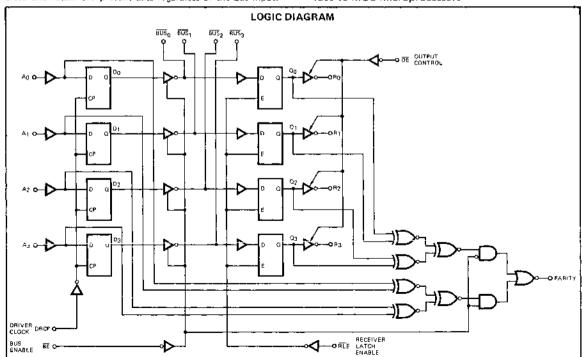
Data from the A input is inverted at the BUS output. Likewise, data at the BUS input is inverted at the receiver output. Thus, data is non-inverted from driver input to receiver output. The four receivers each feature a built-in D-type latch that is controlled from the buffered receiver latch enable (RLE) input. When the \overline{RLE} input is LOW, the latch is open and the receiver outputs will follow the bus inputs (BUS data inverted and \overline{OE} LOW). When the \overline{RLE} input is HIGH, the latch will close and retain the present data regardless of the bus input.

The four latches have three-state outputs and are controlled by a buffered common three-state control (\overline{OE}) input. When \overline{OE} is HIGH, the receiver outputs are in the high-impedance state

The Am2917 features a built-in four-bit odd parity checker/generator. The bus enable input (\overline{BE}) controls whether the parity output is in the generate or check mode. When the bus enable is LOW (driver enabled), odd parity is generated based on the A field data input to the driver register. When \overline{BE} is HIGH, the parity output is determined by the four latch outputs of the receiver. Thus, if the driver is enabled, parity is generated and if the driver is in the high-impedance state, the BUS parity is checked.

DISTINCTIVE CHARACTERISTICS

- Quad high-speed LSI bus-transceiver
- · Three-state bus driver
- D-type register on driver
- Bus driver output can sink 40mA at 0.5V max
- Internal odd 4-bit parity checker/generator
- Receiver has output latch for pipeline operation
- Three-state receiver outputs sink 12mA
- Advanced low-power Schottky processing
- 100% reliability assurance testing in compliance with MIL-STD-883
- 3.5V minimum output high voltage for direct interface to MOS microprocessors





ORDERING INFORMATION

Package Type	Temperature Range	Order Number
Molded DIP	0°C to +70°C	AM2917PC
Hermetic DIP	0°C to +70°C	AM2917DC
Dice	0°C to +70°C	AM2917XC
Hermetic DIP	-55°C to +125°C	AM2917DM
*Hermetic Flat Pak	-55°C to +125°C	AM2917FM
Dice	–55°C to +125°C	AM2917XM

^{*}Available on special order.

MAXIMUM RATINGS (Above which the useful life may be impaired)

Storage Temperature	-65°C to +150°C
Temperature (Ambient) Under Bias	-55°C to +125°C
Supply Voltage to Ground Potential	-0.5 V to +7 V
DC Voltage Applied to Outputs for HIGH Output State	-0.5 V to +Vcc max.
DC Input Voltage	-0.5 V to +5.5 V
DC Output Current, Into Outputs (Except BUS)	30 mA
DC Output Current, Into Bus	100 mA
DC Input Current	-30 mA to +5.0 mA

ELECTRICAL CHARACTERISTICS

The following conditions apply unless otherwise noted:

Am2917XC (COM'L) $T_A = 0^{\circ}C$ to $470^{\circ}C$ $V_{CC}MIN. = 4.75V - V_{CC}MAX. = 5.25V$

Am2917XM (M1L) $T_A = -65^{\circ}C \text{ to} + 125^{\circ}C \text{ V}_{GC} \text{MiN.} = 4.50 \text{ V}_{CC} \text{MAX.} = 5.50 \text{ V}$

BUS INPUT/OUTPUT CHARACTERISTICS OVER OPERATING TEMPERATURE RANGE

Parameters	Description	Test Conditions	(Note 1)	Min.	Тур.	Max.	Units	
V	Bus Output LOW Voltage	VCC = MIN.	IOL = 24 mA			0.4	Volts	
VOL	Bus Output LOW Voltage	ACC - MICO.	IOL = 40mA		7.7	0.5	VOIIS	
V _{OH}	Bus Output HIGH Voltage	VCC = MIN.	I _{OH} = -20mA	2.4			Volts	
			VO = 0.4 V			-200		
10	Bus Leakage Current (High Impedance)	V _{CC} = MAX. Bus enable = 2.4 V	V _O = 2.4 V			50	μА	
		Das enoble - 2.4 v	VO = 4.5 V			100		
IOFF	Bus Leakage Current	V _O = 4.5 V			,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,	100		
.022	(Power OFF)	$V^{CC} = 0 V$				100	μΑ	
V _{IH}	Receiver Input HIGH Threshold	Bus enable = 2.4 V		2.0	,		Volts	
VIL	Receiver Input LOW Threshold	Bus enable = 2,4 V				0.8	Volts	
I _{SC}	Bus Output Short Circuit Current	V _{CC} = MAX. V _O = 0 V		-50	-85	130	mA	



ELECTRICAL CHARACTERISTICS

The following conditions apply unless otherwise noted:

Am2917XC (COM'L) TA = 0°C to +70°C VCCMIN. = 4.75 V VCCMAX. = 5.25 V Am2917XM (MIL) TA = -55°C to +126°C VCCMIN. = 4.50 V VCCMAX. = 5.50 V

DC CHARACTERISTICS OVER OPERATING TEMPERATURE RANGE

arameters	Description	Test Conc			Min.	Typ. (Note 2)	Max.	Units
,		$V_{CC} = MIN.$ $MIL: I_{OH} = -1.0 mA$ $V_{IN} = V_{IL} \text{ or } V_{IH}$ $COM^*L: I_{OH} = -2.6 m$		I _{OH} = -1,0mA	2.4	3.4		
V OH	Receiver Output HIGH Voltage			'L: I _{OH} = -2.6mA	2.4	3.4		Volts
	output more variage	V _{CC} = 5.0 V, I _{OH} = -	100µA		3.5	1		1
VoH	Parity	VCC = MIN., IOH = -	Αμ066	MIL	2.5	3.4		Volts
- OH	Output HIGH Voltage	AIM = AIH or AIT		COM1L	2.7	3.4		Voits
		34 AMIN I		1 _{OL} = 4.0mA		0.27	0.4	
VOL	Output LOW Voltage (Except Bus)	V _{CC} = MIN. V _{IN} = V _{IL} or V _{IH}		I _{OL} = 8.0mA		0.32	0.45	Volts
	(CAGO)	IOL = 12mA			0.37	0.5	1	
V _{IH}	Input HIGH Level (Except Bus)	Guaranteed input logical HIGH for all inputs			2.0			Volts
VIL	Input LOW Level (Except Bus)	Guaranteed input logic for all inputs	Guaranteed input logical LOW for all inputs				8.0	Volt
V _I	Input Clamp Voltage (Except Bus)	VCC = MIN., IIN = -	8mA				-1.2	Volt
,	Input LOW Current (Except Bus)	N		BE, RLE			-0.72	
III.	INDUT LOW Current (Except Bus)	VCC = MAX., VIN = (J,4 V	All other inputs			-0.36	mΑ
ΉΗ	Input HIGH Current (Except Bus)	V _{CC} = MAX., V _{IN} = 2	?.7 V				20	μΑ
Τį	Input HIGH Current (Except Bus)	V _{CC} = MAX., V _{IN} = 7.0 V				100	μА	
I _{SC}	Output Short Circuit Current (Except Bus)	V _{CC} = MAX.		-30		-85	mA	
1 _{CC}	Power Supply Current	V _{CC} = MAX.		~		63	95	mA
Io	Off-State Output Current	V _{CC} = MAX.		V _O = 2.4 V			20	
·U	(Receiver Outputs)	, , CC = MICAY;	Vo = 0.4 V		··†	-20	μΑ	

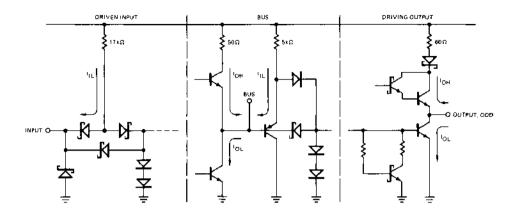
	IG CHARACTERISTICS ERATING TEMPERATURE RA	ANGE	Am2917XM Typ.				Am2917XC Typ.			
arameters	Description	Test Conditions	Min,	(Note 2)	Max.	Min.	(Note 2)	Max.	Unit	
^t PHL		C _L (BUS) = 50pF		21	36	1	21	32		
tPLH	Driver Clock (DRCP) to Bus	R _L BUS) = 130Ω		21	36		21	32	d ns	
tzH, tZL	Bus Enable (BE) to Bus		1	13	26		13	23		
tHZ, tLZ	Bus Enable (BE) to Bus			13	26]	13	23	nş	
ts) _	23			20			ns	
th	A Data Inputs		8.0	1		6.0			1 "	
tpw	Clock Pulse Width (HIGH)		20			17			nş	
tPLH .	Bus to Receiver Output	e \ \ / / .		18	30	ì	18	27	1	
PHL	(Latch Enabled) C/	FIG		18	30		18	27	ns	
tPLH	Latch Enable to Receiver Output	GE ATIO	11.	21	30		21	27		
tPHL .	Laten Enable to Receiver Output	WITH	YO	21	30	T	21	27	- ns	
tş		ECIFICATIONS NGEWITHOUS ARES CL = 15 PF NOTION	17			14				
th	Bus to Latch Enable (RLE)	CL = 15pt-100	8.00			4.0			l ns	
tpLH	A Data to Odd Parity Out	NL - 2.0K32 1/	£ , C>	21	36	[21	32	T	
tPHL .	(Driver Enabled)			21	36		21	32	ns ns	
^t PLH	Bus to Odd Parity Out			21	36		21	32	ns	
tPHL	(Driver Inhibit)			21	36		21	32] "`	
tPLH .	Latch Enable (RLE) to Odd]		21	36		21	32		
tPHL	Parity Output			21	36		21	32] "	
tZH, tZL	Output Control to Output]		14	26		14	23		
t _{HZ} , t _{LZ}	Output Control to Output			14	26		14	23	nş	

Notes: 1. For conditions shown as MIN, or MAX., use the appropriate value specified under Electrical Characteristics for the applicable device type.

Typical limits are at V_{CC} = 5.0 V, 25°C ambient and maximum loading.
 Not more than one output should be shorted at a time. Duration of the short circuit test should not exceed one second.

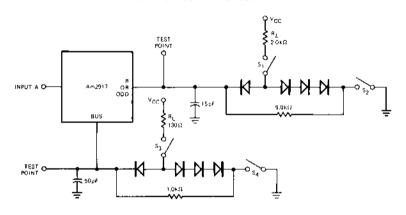
RAYTHEON

INPUT/OUTPUT CURRENT INTERFACE CONDITIONS



Note: Actual current flow direction shown,

SWITCHING TEST CIRCUIT



Note: Bus to Receiver output delay is measured by clocking data into the driver register and measuring the BUS to R combinatorial delay.



FUNCTION TABLE

	INPUTS					NAL	BUS	ОШТРИТ	FUNCTION	
Ai	DRCP	BE	RLE	ŌĒ	Di	Qi	BUSi	Ri	. 5.1.2.7.2.7	
Х	X	Н	×	×	Х	х	z	Х	Driver output disable	
Х	Х	Х	×	Н	Х	Х	Х	Z	Receiver output disable	
x	х	н	L	L	Х	L	L	н	Driver output disable and receive data	
х	х	Н	L	L	x	Н	Н	L	via Bus input	
х	X	Х	н	Х	X	NC	Х	Х	Latch received data	
L	1	Х	Х	Х	L	Х	Х	Х	Load driver register	
Н	t	Х	x	×	н	Х	х	X	Load driver register	
х	ı.	Х	X	Х	NC	×	Х	×	No driver clock restrictions	
х	H	Х	х	х	NC	×	х	×	No driver clock restrictions	
х	х	L	х	Х	L	X	Н	×	Deiro Bur	
х	x	L	x	х	н	×	L	x	Drive Bus	

H = HIGH

Z = High Impedance X = Don't Care

1 : 0, 1, 2, 3

L = LOW

NC = No Change

1 = LOW-to-HIGH Transition

PARITY OUTPUT FUNCTION TABLE

BE	ODD PARITY OUTPUT				
L	ODD = A0 @ A1 @ A2 @ A3				
н	ODD = $Q_0 \oplus Q_1 \oplus Q_2 \oplus Q_3$				

DEFINITIONS

DRCP Driver Clock Pulse. Clock pulse for the driver register.

BE Bus Enable. When the Bus Enable is LOW, the four drivers are in the high impedance state.

BUS₀, BUS₁, BUS₂, BUS₃ The four driver outputs and receiver inputs (data is inverted).

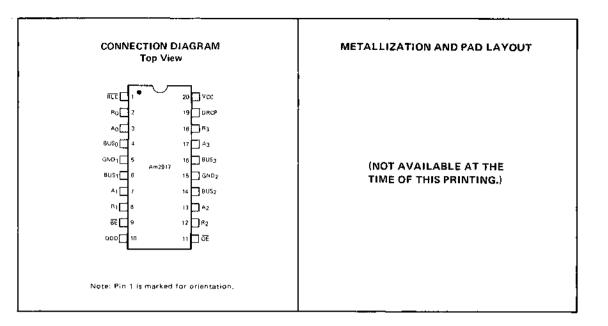
Ro, R1, R2, R3. The four receiver outputs. Data from the bus is inverted while data from the A or B inputs is noninverted.

RLE Receiver Latch Enable, When RLE is LOW, data on the BUS inputs is passed through the receiver latches. When RLE is HIGH, the receiver latches are closed and will retain the data independent of all other inputs.

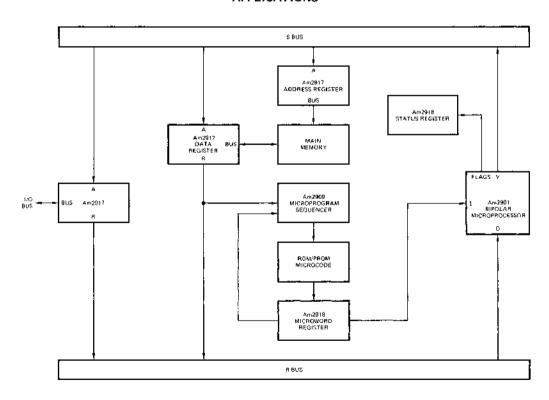
ODD Odd parity output. Generates parity with the driver enabled, checks parity with the driver in the high-impedance state.

OE Output Enable. When the OE input is HIGH, the four three-state receiver outputs are in the high-impedance state.





APPLICATIONS



The Am2917 can be used as an I/O Bus Transceiver and Main Memory I/O Transceiver in high-speed Microprocessor Systems.



FUNCTIONAL DESCRIPTION

New Schottky circuits such as the Am2918 register provide the design engineer with additional flexibility in system configuration.— especially with regard to bus structure, organization and speed. The Am2918 is a quadruple D-type register with four standard totem pole outputs and four three-state bus-type outputs. The 16-pin device also features a buffered common clock (CP) and a buffered common output control (OE) for the Y outputs. Information meeting the set-up and hold requirements on the D inputs is transferred to the Q outputs on the LOW-to-HIGH transition of the clock.

The same data as on the Q outputs is enabled at the three-state Y outputs when the "output control" $\{\overline{OE}\}$ input is LOW. When the \overline{OE} input is HIGH, the Y outputs are in the high-impedance state.

The Am2918 register can be used in bipolar microprocessor designs as an address register, status register, instruction register or for various data or microword register applications. Because of the unique design of the three-state output, the device features very short propagation delay from the clock to the Q or Y outputs. Thus, system performance and architectural design can be improved by using the Am2918 register. Other applications of Am2918 register can be found in microprogrammed display systems, communication systems and most general or special purpose digital signal processing equipment.

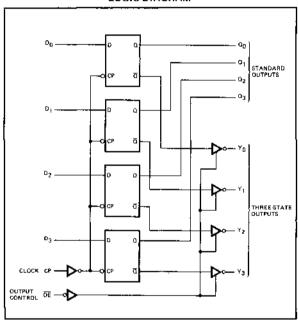
DISTINCTIVE CHARACTERISTICS

- Advanced Schottky technology
- Four D-type flip-flops
- Four standard totem-pole outputs
- Four three-state outputs
- 75 MHz clock frequency
- 100% reliability assurance testing in compliance with MIL-STD-883

ORDERING INFORMATION

	ckage ype	Temper Rang		Order Number
Mole	fed DIP	0°C to −	+70°C	AM2918PC
Herm	etic DIP	0°C to −		AM2918DC
(Dice	0°C to −		AM2918XC
Herm	etic DIP	–55°C to		AM2918DM
Hermeti	c Flat Pack	_55°C to		AM2918FM
į.	Dice	55°C to	+125°C	AM2918XM

LOGIC DIAGRAM





	0
Storage Temperature	_65°C to +150°C
Temperature (Ambient) Under Bias	-55°C to +126°C
Supply Voltage to Ground Potential (Pin 16 to Pin 8) Continuous	-0,5V to +7V
DC Voltage Applied to Outputs for HIGH Output State	-0.5V to +V _{CC} max.
DC Input Voltage	-0.5V to +5.5V
DC Output Current, Into Outputs	30mA
DC Input Current	-30mA to +5.0mA

ELECTRICAL CHARACTERISTICS OVER OPERATING TEMPERATURE RANGE (Unless Otherwise Noted)

Am2918XC Am2918XM	$T_A = 0^{\circ}C \text{ to } + 70^{\circ}C$ $T_A = -55^{\circ}C \text{ to } +125^{\circ}C$	V _{CC} = 5.0V ± 5% (COM'L) V _{CC} = 5.0V + 10% (MIL)		1N = 4 75V IN. = 4.5V		X = 5.25V X. = 5.5V	Typ.		
Parameters	Description	Test Conditions (Note 1)			Min.	(Note 2)	Max.	Units	
v oн	Output HIGH Voltage	V _{CC} = MIN.,	$\begin{array}{c c} \Omega & I_{OH} = -1 \text{mA} \end{array}$	MIL	2.5	3,4			
)H = -1mA	COMIL	2.7	3.4		Volts
		VIN - VIH or VIL		-2mA	2.4	3,4		01/103	
			[`] :	XC, I _{ОН} = -	-6.5mA	2.4	3.4		1
v ol	Output LOW Voltage (Note 6)	V _{CC} = MIN., I _{OL} = 20mA V _{IN} = V _{IH} or V _{IL}					0.5	Volts	
V _{1H}	Input HIGH Level	Guaranteed input logical HIGH voltage for all inputs			2.0			Volts	
VIL	Input LOW Level	Guaranteed input logical LOW voltage for all inputs					9.0	Valts	
v ₁	Input Clamp Voltage	V _{CC} = M(N ₋ , t _{1N} = -18mA					-1.2	Volts	
I _{IL} (Note 3)	Input LOW Current	V _{CC} = MAX., V _{IN} = 0.5V					-2,0	mA	
I _{IH} (Note 3)	Input HIGH Current	V _{CC} = MAX., V _{IN} = 2.7V					50	μА	
i ₁	Input HIGH Current	V _{CC} = MAX., V _{IN} = 5.5V					1.0	mA	
' o	Y Оцірит Off-State Leakage Current	$V_{CC} \approx MAX$. $V_{O} = 2.4V$ $V_{O} = 0.4V$		V _Q = 2	.4V			50	μA
						50			
¹sc	Output Short Circuit Current (Note 4)	V _{CC} = MAX.		-40		-100	mA		
¹ CC	Power Supply Current	V _{CC} = MAX. (Note 5	;}		1		80	120	mA

For conditions shown as MIN, or MAX., use the appropriate value specified under Electrical Characteristics for the applicable device type.
 Typical limits are at V_{CC} = 5.0V, T_{A'} = 25°C ambient and maximum loading.
 Actual input currents = Unit Load Current x Input Load Factor (see Loading Rules).
 Not more than one output should be shorted at a time, Duration of the short circuit test shoud not exceed one second.
 I_{CC} is measured with all inputs at 4.5V and all outputs open.
 Measured on Q outputs with Y outputs open, Measured on Y outputs with Q outputs open.

Switching Characteristics ($T_A = +25^{\circ}C$, $V_{CC} = 5.0V$, $R_L = 280\Omega$)

arameters	Description	Test Conditions	Min.	Тур.	Max.	Units
t₽LH	Charles O. O. da			6.0	9.0	
¹₽HL	Clock to Q Output	C _L = 15pF		8.5	13	п\$
tpw .	Clock Pulse Width		7.0			nş
t _š	Dața		5.0			пѕ
th	Data		3.0			ns
tPLH	Clock to Y Output (OE LOW)			6.0	9.0	
tPHL				8.5	13	ns
tzH	_	CL = 5pF		12.5	19	
^t ZL	Output Control to Output	оц орг		12	18]
tHZ		C _L = 50 pF		4.0	6.0	nş
^t LZ		□ □ 50 pr		7.0	10.5]
fmax	Maximum Clock Frequency	CL = 15pF	75	100	İ	MHz



DEFINITIONS

D; The four data inputs to the register.

Q_i The four data outputs of the register with standard totern-pole active pull-up outputs. Data is passed non-inverted.

 Y_i The four three-state data outputs of the register. When the three-state outputs are enabled, data is passed non-inverted. A HIGH on the "output control" input forces the Y_i outputs to the high-impedance state.

CP Clock. The buffered common clock for the register. Enters data on the LOW-to-HIGH transition.

OE Output Control. When the OE input is HIGH, the Yi outputs are in the high-impedance state. When the OE input is LOW, the TRUE register data is present at the Yi outputs.

TRUTH TABLE

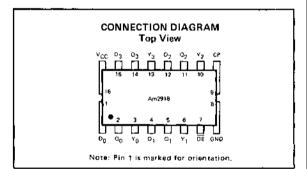
	PUTS	ООТ	INPUTS		
NOTES	Υ	a	D	CLOCK CP	ŌĒ
	Z	NC	×		н
_	2	NC	х	н	н
_	z	L	L	[t	н
-	z	н	н	F	н
l –	Ļ	L	L	[†	L
_	н	н :	H	1 1	L
1	L	L	_	_	L
1 1	н	н	_	1 -	L

L = LOW H = HIGH X = Don't care NC = No change

1 = LOW to HIGH transition

Z = High impedance

Note: 1. When \overline{OE} is LOW, the Y output will be in the same logic state as the Q output.

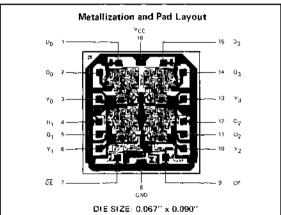


LOADING RULES (In Unit Loads)

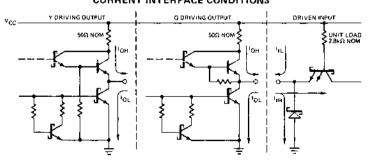
			Fan-out		
Input/Output	Pin No.'s	Input Unit Load	Output HIGH	Output LOW	
D ₀	1	1			
a ₀	2	-	20	10*	
Y ₀	3	-	40/130	10*	
D ₁	4	1		_	
Q ₁	5		20	10*	
Y 1	6	-	40/130	10*	
ŌĒ	7	1	-	-	
GND	8	-	_		
CP	9	1			
Y2,	10		40/130	10*	
Q 2	11		20	10*	
D ₂	12	1			
Y ₃	13		40/130	10*	
	14	-	20	10*	
D ₃	15	1	·		
V _{CC}	16	-	_		

A Schottky TTL Unit Load is defined as 50µA measured at 2.7V HIGH and -2.0mA measured at 0.5V LOW.

*Fan-out on each Ω_1 and Y_4 output pair should not exceed 15 unit loads (30mA) for $i=0,\,1,\,2,\,3,$



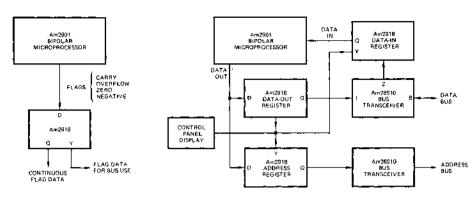
SCHOTTKY INPUT/OUTPUT CURRENT INTERFACE CONDITIONS



Note: Actual current flow direction shown,

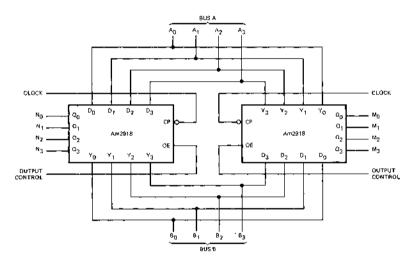


APPLICATIONS

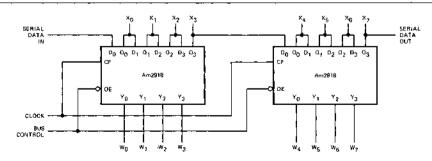


The Am2918 as a 4-Bit status register

The Am2918 used as data-in, data-out and address registers.



The Am2918 can be connected for bi-directional interface between two buses. The device on the left stores data from the A-bus and drives the A-bus. The device on the right stores data from the B-bus and drives the A-bus. The output control is used to place either or both drivers in the high-impedance state. The contents of each register are available for continuous usage at the N and M ports of the device.

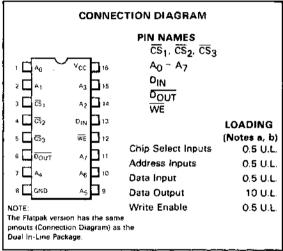


8-Bit serial to parallel converter with three-state output (W) and direct access to the register word (X).



DESCRIPTION

The 93411 and 93411A are high speed 256-bit TTL random access memories with full decoding on chip. They are organized 256 words by one bit and are designed for scratchpad, buffer and distributed main memory applications. The devices have three chip select lines to simplify their use in larger memory systems. Address input pin locations are specifically chosen to permit maximum packaging density and for ease of PC board layout. An uncommitted collector output is provided to permit "OR-ties" for ease of memory expansion.



DISTINCTIVE CHARACTERISTICS

- Replacement for 54/74S206 and equivalent devices
- Organization 256 words by one bit
- Three high speed chip select inputs
- Typical access time

93411A	Commercial	40nsec
93411	Commercial	45nsec
93411	Military	45nsec

- On chip decoding
- Power dissipation 1.2mW/bit
- Power dissipation decreases with temperature

**Available in both ceramic DIP package options.

Inverted data output

ORDERING INFORMATION

ORDER CODE	PACKAGE	TEMPERATURE*
45 ns Typical Access Time		
93411DC	Ceramic DIP**	Commercial
93411DM	Ceramic DIP**	Military
93411PC	Plastic DIP	Commercial
93411FM	Fiatpak	Mittary
40 ns Typical Access Time		
93411ADC	Ceramic DIP**	Commercial
*Commercial = 0°c to +75°C Military = ~55°C to +125°C		

NOTES

- a. 1 Unit Load (U.L.) = 40 µA HIGH / 1.6 mA ŁOW
- b. 10 U L is the output LOW drive factor. An external pull-up resistor is needed to provide HIGH level drive capability. This output will sink a maximum of 16 mA at $V_{OUT} = 0.45$ V.

LOGIC SYMBOL LOGIC DIAGRAM ADDRESS DECODER ① An -(f) AT .55, (1) ARPAY (€) A₀ 42 Аз 93411/93411A Α4 SENSE AFTE Α5 WILLIAM VER 10 DOUT JECODER Y AUDRESS V_{CC} = Pin 16 GND = Pin 8 V_{CC} = Pin 16 A_C വി.എി.എി = Pin Numbers GND = Pin 8



FUNCTIONAL DESCRIPTION

The 93411/93411A are fully decoded 256-bit random access memories organized 256 words by one bit. Word selection is achieved by means of an eight-bit address, Ag thru A7.

Three chip select inputs are provided for logic flexibility. For larger memories, the fast chip select access time permits the decoding of Chip Select, CS, from the address without increasing address access time.

The read and write operations are controlled by the state of the active LOW Write Enable (\overline{WE} , pin 12). With \overline{WE} held LOW and the chip selected, the data at D_{1N} is written into the addressed location. To read, \overline{WE} is held HIGH and the chip selected. Data in the specified location is presented at $\overline{D_{OUT}}$.

Uncommitted collector outputs are provided to allow maximum flexibility in output connection. In many applications, such as memory expansion, the outputs of several 93411's or 93411A's can be tied together. In other appli-

cations the wired-OR is not used. In either case an external pull-up resistor of value R_L must be used to provide a HIGH at the output when it is off. Any value of R_L within the range specified below may be used.

$$\frac{V_{CC} \text{ (MAX)}}{16 - \text{ F.O. (1.6)}} \le R_L \le \frac{V_{CC} \text{ (MIN)} - V_{OH}}{n \text{ (ICEX)} + \text{ F.O. (0.04)}}$$

R_I is in $k\Omega$

n = number of wired-OR outputs tied together F.O. = number of TTL Unit Loads (U.L.) driven ICEX = Memory Output Leakage Current in mA

VOH = Required Output HIGH level at Output Node

The minimum value of R_L is limited by output current sinking ability. The maximum value of R_L is determined by the output and input leakage current which must be supplied to hold the output at V_{OH}.

TABLE 1 TRUTH TABLE

	INPUTS				OUTPUT]
PIN 3	ČŠ ₂ PÍŇ 4	PIN 5	WE	D _{IN}	D _{OUT}	MODE
Н	х	×	×	x	H	Not Selected
х	н	x	×	l x	н	Not Selected
x	×	н	×	×	H	Not Selected
L	L	L	l ι	L	H	Write "0"
L	L	L	Ĺ	н	H	Write "1"
L	Ĺ	L	н	×	Tuod	Read inverted data from addressed location

H = HIGH Voltage Level

ABSOLUTE MAXIMUM RATINGS (above which the useful life may be impaired)

Storage Temperature

Temperature (Ambient) Under Bias

V_{CC} Pin Potential to Ground Pin

*Input Voltage (dc)

*Input Current (dc)

**Voltage Applied to Outputs (output HIGH)

Output Current (dc) (output LOW)

-65°C to +150°C -55°C to +125°C

-0.5 V to +7.0 V -0.5 V to +5.5 V

-0.5 V to +5.5 V -12 mA to +5.0 mA

~0.5 V to +5.50 V

+20 mA

^{**}Output Current Limit Required.

GUARANTEED	OPERATING	RANGES

PART NUMBER	SI	JPPLY VOLTAGE (VC	C)	AMBIENT TEMPERATURE
PART NOMBER	MIN	TYP	MAX	Note 4
93411AXC, 93411XC	4.75 V	5.0 V	5.25 V	0°C to +75°C
93411XM	4.50 V	5.0 V	5.50 V	-55°C to +125°C

X = package type, F for Flatpak, D for Ceramic Dip, P for Plastic Dip, See Packaging Information Section for packages available on this product.

L = LOW Voltage Level

X - Don't Care (HIGH or LOW)

^{*}Either Input Voltage limit or Input Current limit is sufficient to protect the inputs.

DC CHARACTERISTICS: Over Operating Temperature Ranges, Notes 1, 2 and 4

SYMBOL	PARAMETER			LIMITS		UNITS	CONDIT	IONE	
3 I WIBQL			MIN	TYP (Note 3)	MAX	UNITS	CONDI	ions	
VOL	Output LOW Vo		0.3	0.45	٧	V _{CC} = MiN, I _O	_L = 16 mA		
V _{IH}	Input HIGH Voli	age	2.0	1.6		v	Guaranteed Input Logical HIGH Voltage for all Inputs		
۷۱۲	Input LOW Volt	age		1.5	0.85	v	Guaranteed Input Logical LOW Voltage for all Inputs		
<u>η</u> μ	Input LOW Curr	rent		-530	-800	μΑ	V _{CC} = MAX, V _{IN} = 0 V		
l _H	input HIGH Cur	rent		1.0	20	μΑ	V _{CC} = MAX, V	/ _{IN} = 4.5 V	
CEX	Output Leakage	Current		1.0	50	μΑ	V _{CC} = MAX, V	OUT = 4.5 V	
V _{CD}	Input Clamp Die	ode Voltage		~1.0	- 1.5	V	V _{CC} = MAX. I	N = -10 mA	
lcc	Power Supply	93411XC 93411AXC		90 100	124 135	mA	$T_{A} = +75^{\circ}C$ $T_{A} = 0^{\circ}C$	V _{CC} = MAX, WE Grounded, all other inputs @ 4.5 V	
-ÇL	Current 93411XM			90 100	117 T _A		$T_A = +125$ °C $T_A = -55$ °C		

AC CHARACTERISTICS: Over Guaranteed Operating Ranges. Notes 1, 2, 4, 5, 6

	·	9	3411A	XC		934112	кc	9	34112	M	j	1
SYMBOL	CHARACTERISTIC	MIN	TYP (Note 3)	MAX	MIN	TYP (Note 3)	MAX	MIN	TYP (Note 3)	MAX	UNITS	CONDITIONS
READ MODE TACS TRCS TAA	DELAY TIMES Chip Select Time Chip Select Recovery Time Address Access Time		25 25 40	30 25 45		25 25 45	30 25 55		25 25 45	40 35 65	ns	See Test Circui and Waveform Note 5
WRITE MODE [†] WS [†] WR	DELAY TIMES Write Disable Time Write Recovery Time INPUT TIMING	10	20 25	35 40	10	20 25	35 40	10	20 25	4 5 50	ns	
t _{VV}	REQUIREMENTS Write Pulse Width (to guarantee write)	40	25		40	25		50	25			See Test Circu and Waveform
twsp	Data Set-Up Time Prior to Write	0	0		0	0	,	0	0			Note 6
I _{WHD}	Data Hold Time After Write	5	0		5	0		5	O		1	
^t WSA	Address Set-Up Time	0	0		0	0		0	0		ns	
^t WHA	Address Hold Time	5	0		5	0		5	0		1	Į.
twscs twhcs	Chip Select Set-Up Time Chip Select Hold Time	0 5	0		5	0		0 5	0			
C _l	Input Lead Capacitance Output Lead Capacitance		4 7	5 8		4 7	5 8		4 7	5 8	ρF	Measured with pulse technique

NOTES:

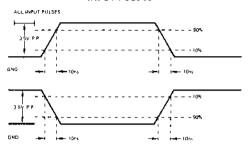
- 1. Conditions for testing, not shown in the Table, are chosen to guarantee operation under "worst case" conditions.
- 2. The specified LIMITS represents the "worst case" value for the parameters. Since these "worst case" values normally occur at the temperature and supply voltage extremes, additional noise immunity and guard banding can be achieved by decreasing the allowable system operating ranges
- 3. Typical limits are at V_{CC} = 5.0 V, T_A = +25°C, and MAX loading
- 4 The Temperature Ranges are guaranteed with transverse air flow exceeding 400 linear feet per minute and a two minute warm-up. Temperature range of operation refers to case temperature for Flatpaks and ambient temperature for all other packages. Typical thermal resistance values of the package at maximum temperature are:
 - θ_{JA} (Junction to Ambient) (at 400 fpm air flow) = 50°C/Watt, Ceramic DIP; 65°C/Watt, Plastic DIP; NA, Flatpak,
 - $\theta_{\rm JA}$ (Junction to Ambient) (still air) = 90°C /Waii, Ceramic DIP, 110°C /Waii, Plastic DIP, NA, Flatpak, $\theta_{\rm JC}$ (Junction to Case) = 25°C / Waii, Ceramic DIP; 25°C / Waii, Plastic DIP; 10°C / Wait, Flatpak.
- 5. The MAX address access time is guaranteed to be the "worst case" bit in the memory using a pseudo random testing pattern.
- 6. t_{W} measured at t_{WSA} = MIN, t_{WSA} measured at t_{W} = MIN.



AC TEST LOAD AND WAVEFORM

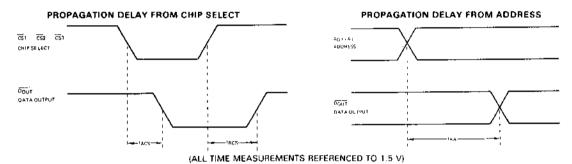
B3411/93411A 256-BHT READWAITE MEMORY DOUT DOUT DOUT DOUT JOOF JOAPACITANCE INCLUDING SCOPE AND JIG)

INPUT PULSES

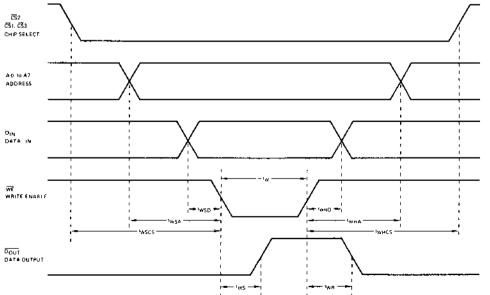


AC WAVEFORMS

READ MODE



WRITE MODE



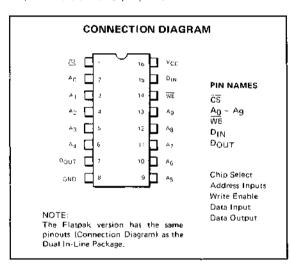
(ALL TIME MEASUREMENTS REFERENCED TO 1.5 V)



DESCRIPTION

The 93415 and 93415A are 1024-bit read/write random access memories organized 1024 words by one bit. They are designed for buffer control storage and high performance main memory applications. The devices have typical access times of 40nsec for the 93415 and 30nsec for the 93415A.

The 93415 and 93415A include full decoding on chip, separate Data Input and Data Output lines and an active LOW Chip Select. They are fully compatible with standard DTL and TTL logic families and have an uncommitted collector output for ease of memory expansion.



DISTINCTIVE CHARACTERISTICS

- Uncommitted collector output
- · TTL inputs and output
- Non-inverting data output
- Organized 1024 words by one bit
- Typical read access time

93415A	Commercial	30nsec
93415	Commercial	40nsec
93415	Military	40nsec

- Chip select access time 15nsec typical
- Power dissipation 0.5mW/bit typical
- Power dissipation decreases with increasing temperature

ORDERING INFORMATION

ORDER CODE	PACKAGE	TEMPERATURE
40 ns Typical Access Time		
93415DC	Ceramic DIP	Commercial
93415DM	Ceramic DtP	Military
93415PC	Plastic DIP	Commercial
93415FM	Flatpak	Military
30 ns Typical Access Time		
93415ADC	Caramic DIP	Commercial
93415APC	Plastic DIP	Commercial

LOGIC SYMBOL

GND = Pin 8

LOGIC DIAGRAM

CS DIN ACRO DOTY (125 32 x 32 p8645 Αı A٥ Aз 44 93415/93415A Aş Αö WAITE DRIVER Α7 A₈ A4 DECOUR V_{CC} = Pin 16

45 —

F 62 -

A3 -

V_{CC} ≈ Pin 16 GND ≈ Pin 8

O = Pin Numbers

DECCCEA

^{*}Commercial = 0°C to +75°C Military = +55°C to +125°C

FUNCTIONAL DESCRIPTION

The 93415/93415A are fully decoded 1024-bit random access memories organized 1024 words by one bit. Bit selection is achieved by means of a 10-bit address, An to Ag.

The Chip Select input provides for memory array expansion. For large memories, the fast chip select access time permits the decoding of chip select (CS) from the address without affecting system performance.

The read and write operations are controlled by the state of the active LOW Write Enable (WE, Pin 14). With WE held LOW and the chip selected, the data at DIN is written into the addressed location. To read, WE is held HIGH and the chip selected. Data in the specified location is presented at DOUT and in non-inverted.

Uncommitted collector outputs are provided to allow maximum flexibility in output connection. In many applications such as memory expansion, the outputs of many 93415's or 93415A's can be tied together. In other applications the wired-OR is not used. In either case an external pull-up resistor of RL value must be used to provide a HIGH at the output when it is off. Any Rt value within the range specified below may be used.

$$\frac{V_{CC}(MIN)}{I_{OL} - FO(1.6)} \le R_L \le \frac{V_{CC}(MIN) - V_{OH}}{n(I_{CEX}) + FO(0.04)}$$

R_L is in kΩ

n = number of wired-OR outputs tied together FO = number of TTL Unit Loads (UL) driven

ICEX = Memory Output Leakage Current

VOH = Required Output HIGH Level at Output Node Int = Output LOW Current

The minimum Rt value is limited by output current sinking ability. The maximum RL value is determined by the output and input leakage current which must be supplied to hold the output at VOH. One unit load = 40µA HIGH/ 1.6mA LOW.

TABLE I - TRUTH TABLE

	INPUT	S	QUTPUT	MODE
cs	WE	DIN	Open Collector	•
н	х	х	н	NOT SELECTED
Ł	L.	L	Н	WRITE "0"
L	L	н	Н	WRITE "1"
L	н	х	Dout	READ

H = HIGH Voltage Level

L = LOW Voltage Level

X = Don't Care (HIGH or LOW)

ABSOLUTE MAXIMUM RATINGS (above which the useful life may be impaired)

Storage Temperature

Temperature (Ambient) Under Bias V_{CC} Pin Potential to Ground Pin

"Input Voltage (dc)

*Input Current (dc)

Voltage Applied to Outputs (Output HIGH)

Output Current (dc) (Output LOW)

-65°C to +150°C

-55°C to +125°C

-0.5 V to +7.0 V

-0.5 V to +5.5 V

-12 mA to +5.0 mA

-0.5 V to +5.5 V

+20 mA

*Either input voltage or input current	limit is sufficient to protect the input.
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PART NUMBER	\$L	IPPLY VOLTAGE (VC	c)	AMBIENT TEMPERATURE (TA
PART NUMBER	MIN	TYP	MAX	(Note 4)
93415XC, 93415AXC	4.75 ∨	5.0 ∨	5.25 ∨	0°C to +75°C
93415XM	4.50 V	5.0 V	5.50 V	-55° C to +125° C

X = package type; F for Fistpak, D for Ceramic DIP, P for Plastic DIP. See Packaging Information Section for packages available on this product,



DC CHARACTERISTICS: Over Operating Temperature Ranges (Notes 1, 2, 4)

			LIMITS			CONDITIONS			
SYMBOL 	CHARACTERISTIC	MIN	TYP (Note 3)	MAX	UNITS				
VOL.	Output LOW Voltage		0.3	0.45	٧	V _{CC} = MIN, I _{OL} = 16 mA			
V _{IH}	Input HIGH Voltage	2.1	1.6		V	Guaranteed Input HIGH Voltage for all Inputs			
VIL	Input LOW Voltage		1.5	0.8	V	Guaranteed Input LOW Voltage for all Inputs			
I _{IL}	Input LOW Current		-250	-400	μΑ	V _{CC} = MAX, V _{IN} = 0.4 V			
	Lucus (UCH Courses		1.0	40	μА	V _{CC} = MAX, V _{IN} = 4.5 V			
IIН	Input HIGH Current			1.0	mΑ	V _{CC} = MAX, V _{IN} = 5.25 V			
ICEX	Output Leakage Current		1,0	100	μΑ	V _{CC} = MAX, V _{OUT} = 4.5 V			
V _{CD}	Input Diode Clamp Voltage		-1,0	-1.5	V	V _{CC} = MAX, I _{IN} = ~10 mA			
			95	130	mΑ	TA ≥ 75°C V MAY			
Icc	Power Supply Current			155	mΑ	T _A = 0°C V _{CC} = MAX, All Inputs Grounded			
•	ļ			170	mA	T _A = -55°C			

AC CHARACTERISTICS: Over Guaranteed Operating Ranges (Notes 1, 2, 4, 5, 6)

		9	3415A	xc	,	3415X	:c	9	3415X	M] !	
SYMBOL	CHARACTERISTIC	MIN	TYP	MAX	MIN	TYP	MAX	MIN	TYP	MAX	UNITS	CONDITIONS
			(Note 3	3)		(Note 3	3)		(Note 3	3)		
READ MODE	DELAY TIMES											
[†] ACS	Chip Select Time		15	30		16	40		15	45		See Test Circui
^L BC\$	Chip Select Recovery Time		16	30		20	40		20	50	ris.	and Waveforms
^t AA	Address Access Time		30	45		40	70		40	75		
WRITE MODE	DELAY TIMES											
tws	Write Disable Time		20	30	 !	20	40		20	45	ns	
™R	Write Recovery Time		20	40		25	50		45	55	115	
	INPUT TIMING				1							
	REQUIREMENTS				Ì					ļ	,	
1 _{VV}	Write Pulse Width	35	25		50	25		55	25		i	See Test Circui
	(to guarantee write)											and Waveforms
\WSD	Data Set-Up Time Prior to Write	5	0		5	0		5	0			and waveloning
™HD	Data Hold Time After Write	5	0		5	0	ı	5	0		ns	
^t WSA	Address Set-Up Time	5	0		15	0		15	0		115	
WHA .	Address Hold Time	5	0		5	0	I	5	0			
*WSCS	Chip Select Set-Up Time	5	0		5	0	ı	5	0			
*WHCS	Chip Select Hold Time	5	0		5	0		5	(1			
C ₁	Input Pin Capacitance		4	5		4	5		4	5	рF	
co	Output Pin Capacitance		7	8		7	8,		7	8	P .	

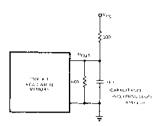
NOTES.

- 1. Conditions for testing, not shown in the Table, are chosen to guarantee operation under "worst case" conditions.
- The specified LIMITS represents the "worst case" value for the parameters. Since these "worst case" values normally occur at the temperature and supply voltage extremes, additronal noise immunity and guard banding can be achieved by decreasing the allowable system operating ranges.
- 3. Typical values at V_{CC} = 5.0 V, T_A = 125°C, and MAX loading.
- Temperature Ranges are guaranteed with transverse air flow exceeding 400 linear feet per minute and a two minute warm-up. Typical thermal resistance values of the package at maximum temperature are:
 - θ JA (Junction to Ambient) lat 400 fpm air flow) = 50° C/Watt, Ceramic DIP; 65° C/Watt, Plastic DIP; NA, Flatpak.
 - $\theta_{
 m JA}$ [Junction to Ambient) [still air] = 90° C/Watt, Ceramic DIP; 110° C/Watt, Plastic DIP; NA, Flatpak.
 - H_{JC} (Junction to Case) = 25° C/Watt, Ceramic DIP, 25° C/Watt, Plastic DIP; 10° C/Watt, Flatpak.
- 5. The MAX address access time is guaranteed to be the "worst case" bit in the memory using a pseudo random testing pattern.
- 6 t_W measured at t_{WSA} = MIN, t_{WSA} measured at t_W MIN.

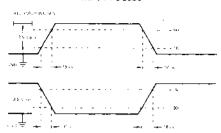


AC TEST LOAD AND WAVEFORM

LOADING CONDITION

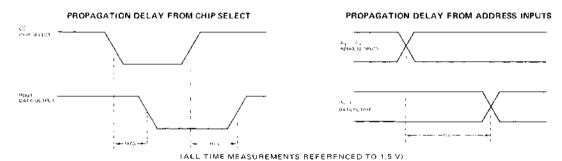


INPUT PULSES

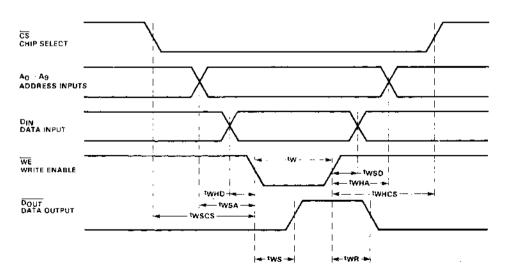


AC WAVEFORMS

READ MODE



WRITE MODE

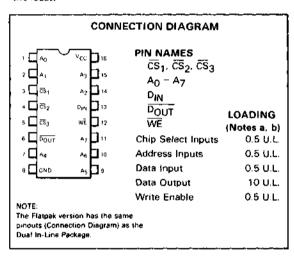


(ALL TIME MEASUREMENTS REFERENCED TO 1.5 V)



DESCRIPTION

The 93421 and 93421A are high speed 256-bit random access memories with full decoding on chip. They are organized 256 words by one bit and are designed for scratchpad, buffer and distributed main memory applications. The devices have three chip select lines to simplify their use in larger memory systems. Address input pin locations are specifically chosen to permit maximum packaging density and for ease of PC board layout. A 3-state output is provided to drive bus organized systems and/or highly capacitive loads.



DISTINCTIVE CHARACTERISTICS

- 3-state output
- Replacement for 54/74S200 and equivalent devices
- Organization 256 words by one bit
- Three high speed chip select inputs
- Typical read access time

93421A	Commercial	30nsec
93421	Commercial	35nsec
93421	Military	35nsec

- On chip decoding
- Power dissipation 1.2mW/bit
- Power dissipation decreases with temperature
- Inverted data output

ORDERING INFORMATION

ORDER CODE	PACKAGE	TEMPERATURE*
35 ns Typical Access Time		
93421DC	Ceramic DIP**	Commercial
93421DM	Ceramic DIP**	Military
93421PC	Plastic DIP	Commercial
93421FM	Flatpak	Military
30 ns Typical Access Time		
93421ADC	Ceramic DIP**	Commercial

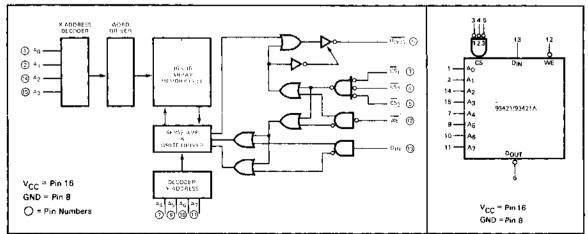
^{*}Commercial = 0°C to +75°C

NOTES

- a. 1 Unit Load (U.L.) = 40 μ A HIGH / 1.6 mA LOW
- b. 10 U.L. is the output LOW drive factor. This output will sink a maximum of 16 mA at V_{QUT} = 0.45 V, and will source a minimum of 10 mA at 2.4 V.

LOGIC DIAGRAM

LOGIC SYMBOL





Military = -55°C to +125°C

^{**}Available in both Ceramic DIP package options

FUNCTIONAL DESCRIPTION

The 93421/93421A are fully decoded 256-bit random access memories organized 256 words by one bit. Word selection is achieved by means of an eight-bit address. An through A7.

Three Chip Select inputs are provided for logic flexibility. For larger memories, the fast chip select access time permits the decoding of Chip Select, \overline{CS} , from the address without increasing address access time.

The read and write operations are controlled by the state of the active LOW Write Enable (\overline{WE} , pin 12). With \overline{WE} held

LOW and the chip selected, the data at D_{IN} is written into the addressed location. To read, \overline{WE} is held HIGH and the chip selected. Data in the specified location is presented at $\overline{D_{OUT}}$.

The 3-state output provides drive capability for higher speeds with high capacitive load systems. The third state (high impedance) allows bus organized systems where multiple outputs are connected to a common bus.

During writing, the output is held in the high impedance state.

TABLE I -	TRUTH	TABLE
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INPUTS				OUTPUT			
ĈŜ₁	cs ₂	CS ₃	CS ₃ WE D _{IN}		TUO ^Q	MODE	
Н	×	×	×	×	HIGH Z	Not Selected	
Х	н	х	x	х	HIGH Z	Not Selected	
х	x	н	x	x	HIGH Z	Not Selected	
L	L	L	L	L	HIGH Z	Write "0"	
L	L	L	L	Н	HIGH Z	Write "1"	
L	L	L	н	×	DOUT	Read inverted data from addressed location	

H = HIGH Voltage Level

TABLE 2 FUNCTION TABLE

5 INSTITUTE	IN	INPUTS				
FUNCTION	CHIP SELECT WRITE ENABLE		OUTPUT			
Write	L	L	HIGH Z			
Read	L	Н	Stored Data			
Not Selected	н	× ×	HIGH Z			

ABSOLUTE MAXIMUM RATINGS (above which the useful life may be impaired)

Storage Temperature

Temperature (Ambient) Under Bias

V_{CC} Pin Potential to Ground Pin

*input Voltage (dc)

*Input Current (dc)

**Voltage Applied to Outputs (output HIGH)
Output Current (dc) (output LOW)

Opipar Carretti (ac) (opipar 20

~65°C to +150°C -55°C to +125°C

-0.5 V to +7.0 V -0.5 V to +5.5 V

-12 mA to +5.0 mA

-0.5 V to +5.50 V

+20 mA

GUARANTEED OPERATING RANGES

PART NUMBER	5	AMBIENT TEMPERATURE		
	MIN	TYP	MAX	Note 4
93421AXC. 93421XC	4.75 V	5.0 V	5.25 V	0°C to +75°C
93421XM	4.50 V	5.0 V	5.50 V	~55°C to +125°C

X = package type: F for Flatpak, D for Ceramic Dip, P for Plastic Dip, See Packaging Information Section for packages available on this product



L = LOW Voltage Level

X - Don't Care (HIGH or LOW)

HIGH Z = HIGH Impedance

^{*}Either Input Voltage limit or Input Current limit is sufficient to protect the inputs.

^{**}Output Current Limit Required.

SYMBOL	DADIMETE	PARAMETER		LIMITS			COMPLETIONS						
STIVIBOL PARAMETE				MMETEM		PARAMETER		PARAMETER		TYP (Note 3)	MAX		CONDITIONS
VOL	Output LOW V		0.3	0.45	٧	V _{CC} = MIN, I _{OL} = 16 mA							
V _{IH}	Input HIGH Vo	ltage	2.0	1.6		٧	Guaranteed Input Logical HIGH Voltage for all Inputs						
V _{IL}	Input LOW Vo	Itage		1.5	0.85	v	Guaranteed Input Logical LOW Voltage for all Inputs						
յլ Կե	Input LOW Cu	rrent		-530	800	μΑ	V _{CC} = MAX, V _{IN} = 0 V						
¹н	Input HIGH Cu	rrent		1.0	20	μΑ	V _{CC} = MAX, V _{IN} = 4.5 V						
OFF	Output Curren	t (HIGH Z)			50 -50	μА	V _{CC} = MAX, V _{OUT} = 2.4 V V _{CC} = MAX, V _{OUT} = 0.5 V						
V _{CD}	Input Clamp D	iode Voltage		-1.0	-1.5	٧	V _{CC} = MAX, I _{IN} = -10 mA						
lcc	Power Supply	93421XC 93421AXC		90 100	124 135	mA	T _A = +75°C V _{CC} = MAX, WE Grounded all other inputs @ 4.5 V						
	Current	93421XM		90 100	117 143		T _A = +125°C T _A = -55°C						
Voн	Output HIGH	93421XC,AXC	2.4	TT		V	I _{OH} = -10.3 mA						
· OH	Voltage	93421XM	2.4		·	V	I _{OH} 5.2 mA						
ios	Output Curren Short Circuit t				-100	mA	V _{CC} = MAX, Note 7						

		93421AXC		93421XC		XC.	93421XM		M			
SYMBOL	CHARACTERISTIC		TYP (Note 3)	MAX	MIN	TYP (Note 3)	MAX	MIN	TYP (Note 3)	MAX	UNITS	CONDITIONS
READ MODE IACS IZRCS IACS Chip Select Access Time Chip Select to HIGH Z Address Access Time			20 20 30	30 30 40		20 20 35	30 30 50		25 20 35	40 40 60	ns	See Test Circuit and Waveforms Note 5
WRITE MODE [†] ZWS [‡] WR	DELAY TIMES Write Disable to HIGH Z Write Recovery Time INPUT TIMING REQUIREMENTS	10	20 25	35 40	10	20 25	35 40	10	20 25	45 50	ns	
tw twsd twsa twsa twscs twecs	Minimum Write Pulse Width Data Set-Up Time Prior to Write Data Hold Time After Write Address Set-Up Time Address Hold Time Chip Select Set-Up Time Chip Select Hold Time	30 0 5 0 5 0 5	0000000		30 0 5 0 5 0 5	000000		40 0 5 0 5 0 5	10 0 0 0 0 0		ns	See Test Circuit and Waveforms Note 6
c _i	Input Capacitance Output Capacitance		2.5 5	3.5 7		2.5 5	3.5 7		2.5 5	3.5 7	pF	Measured with pulse technique

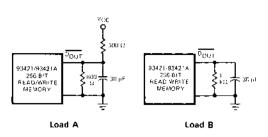
NOTES:

- 1. Conditions for testing, not shown in the Table, are chosen to guarantee operation under "worst case" conditions
- 2. The specified LIMITS represents the "worst case" value for the parameters. Since these "worst case" values normally occur at the temperature and supply voltage extremes, additional noise immunity and guard banding can be achieved by decreasing the allowable system operating ranges
- Typical limits are at V_{CC} = 5.0 V. T_A = +25°C, and MAX loading
 The Temperature Ranges are guaranteed with transverse air flow exceeding 400 linear feet per minute and a two minute warm-up. Temperature range of operation refers to case temperature for Flatpaks and ambient temperature for all other packages. Typical thermal resistance values of the package at maximum temperature are
 - $\theta_{\rm JA}$ (Junction to Ambient) (at 400 fpm air flow) = 50°C / Watt, Ceramic DIP, 65°C / Watt, Plastic DIP, NA, Flatpak, $\theta_{\rm JA}$ (Junction to Ambient) (still air) = 90°C / Watt, Ceramic DIP; 110°C / Watt, Plastic DIP; NA, Flatpak, $\theta_{\rm JC}$ (Junction to Case) = 25°C / Watt, Ceramic DIP; 25°C / Watt, Plastic DIP; 10°C / Watt, Flatpak.
- The MAX address access time is guaranteed to be the "worst case" bit in the memory using a pseudo random testing pattern. τ_W measured at $\tau_W = M$ measured at $\tau_W = M$. Duration of short circuit should not exceed one second.

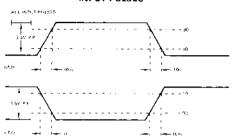


AC TEST LOAD AND WAVEFORM

LOADING CONDITION



INPUT PULSES

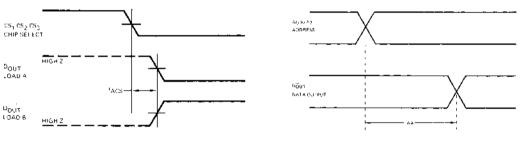


AC WAVEFORMS

READ MODE

PROPAGATION DELAY FROM CHIP SELECT

PROPAGATION DELAY FROM ADDRESS



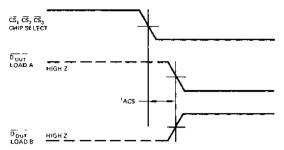
(ALL TIME MEASUREMENTS REFERENCED TO 1.5 V)

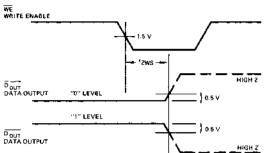
(ALL TIME MEASUREMENTS REFERENCED TO 1.5 V)

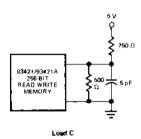


PROPAGATION DELAY FROM CHIP SELECT TO HIGH Z

WRITE ENABLE TO HIGH Z DELAY







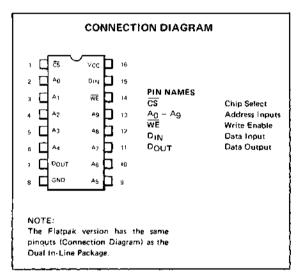
IAII τ_{ZXXX} parameters are measured at a delta of 0.5 V from the logic level and using Load C.)



DESCRIPTION

The 93425 and 93425A are 1024-bit read/write random access memories organized 1024 words by one bit. They are designed for buffer control storage and high performance main memory applications. The devices have typical address access times of 40nsec for the 93425 and 30nsec for the 93425A.

The 93425 and 93425A include full decoding on chip, separate Data Input and Data Output lines and an active LOW Chip Select and Write Enable. They are fully compatible with standard DTL and TTL logic families. A 3-state output is provided to drive bus organized systems and/or highly capacitive loads.



DISTINCTIVE CHARACTERISTICS

- 3-state output
- Organized 1024 words by one bit
- TTL inputs and output 16mA drive capability
- Typical read access time

93425A	Commercial	30nsec
93425	Commercial	40nsec
93425	Military	40nsec

- Chip select access time 15nsec typical
- Non-inverting data output
- Power dissipation 0.5mW/bit typical
- Power dissipation decreases with increasing temperature

ORDERING INFORMATION

ORDER CODE	PACKAGE	TEMPERATURE		
93425DC	Ceramic DIP	Commercial		
93425DM	Ceramic DIP	Military		
93425PC	Plastic DIP	Commercial		
93425FM	Fiatpak	Military		
93425ADC	Ceramic DIP	Commercial		
93425APC	Plastic DIP	Commercial		

*Commercial = 0°C to +75°C Military = -55°C to +125°C

LOGIC DIAGRAM LOGIC SYMBOL DIN ENSE AMP WRITE ORIVERS 93425/93425A (14) Ag ABORESS DECOGER An DOUT V_{CC} = Pin 16 23466 90000 GND = Pin 8 V_{CC} = Pin 16 GND = Pin 8 O = Pin Numbers

FUNCTIONAL DESCRIPTION

The 93425/93425A are fully decoded 1024-bit random access memories organized 1024 words by one bit. Word selection is achieved by means of a 10-bit address, An to Ag.

The chip select (CS) input provides for memory array expansion. For large memories, the fast chip select access time permits the decoding of chip select from the address without increasing address access time.

The read and write operations are controlled by the state of the active LOW write enable (WE, Pin 14), With WE and CS held LOW, the data at DIN is written into the addressed

location. To read, WE is held HIGH and CS held LOW, Data in the specified location is presented at DOLLT and is noninverted.

The 3-state output provides drive capability for higher speeds with high capacitive load systems. The third state (high impedance) allows bus organized systems where multiple outputs are connected to a common bus.

During writing, the output is held in the high impedance

TABLE 1 - TRUTH TABLE

<u> </u>	INPUTS		OUTPUT	
CS	WE	DIN	DOUT	MODE
Н	х	х	HIGH Z	NOT SELECTED
L	L	L	HIGH Z	WRITE "0"
L	L	н	HIGH Z	WRITE "1"
Ļ	н	×	Dout	R€AD

H = H1GH Voltage Level L = LOW Voltage Level X = Don'r Care (HIGH or LOW)

ABSOLUTE MAXIMUM RATINGS (above which the useful life may be impaired.)

Storage Temperature

Temperature (Ambient) Under Bias V_{CC} Pin Potential to Ground Pin

- *Input Voltage (dc)
- *Input Current (dc)
- **Voltage Applied to Outputs (Output HIGH) Output Current (dc) (Output LOW)
- **Output Current Limit Required

-65°C to +150°C -55°C to +125°C -0.5 V to +7.0 V -0.5 V to +5.5 V

+20 mA

-12 mA to +5.0 mA -0.5 V to +5.5 V

- *Either input voltage or input current limit is sufficient to protect the input.

GUARA	NTEED	OPERA	TING	RANGES

PART NUMBER	su	JPPLY VOLTAGE (VC	AMBIENT TEMPERATURE (TAI	
PART NOWBER	MIN	TYP	MAX	(Note 4)
93425XC, 93425AXC	4.75 V	5.0 V	5.25 V	0° C 10 +75 C
93425×M	4 50 V	5.0 V	5.50 ∨	-55 C to +125 C

package type, Fifor Flarpak, Difor Ceramic DIP, Pifor Plastic DIP. See Packaging Information Section for packages available on this product.



DC CHARACTERISTICS: Over Operating Temperature Ranges (Notes 1, 2, 4).

CVIMOO:	i CHARACTERISTIC			LIMITS		LUNITO	CONDITIONS
SYMBOL	CHARACIERIS	110	MIM	TYP (Note 3)	MAX	UNITS	CONDITIONS
VOL	Output LOW Voltage			0.3	0.45	v	VCC - MIN, LOL = 16 mA
VIH	Input HIGH Voltage		2.1	1.6	i	V	Guaranteed Input HIGH Voltage for all Inputs
V _{IL}	Input LOW Voltage			1.5	8.0	V	Guaranteed Input LOW Voltage for all Inputs
TIL	Input LOW Current			-250	-400	μА	V _{CC} = MAX, V _{IN} = 0.4 V *
1	I Input HIGH Current			1.0	40	μА	V _{CC} = MAX, V _{IN} = 4.5 V
li li	Input migh Current			T	1.0	mA	V _{CC} = MAX, V _{IN} = 5.25 V
	Output Current IHIGH	7)			50	μA	V _{CC} = MAX, V _{OUT} = 2.4 V
OFF	L	۷,			-50	7,4	VCC = MAX, VOUT - 0.5 V
los	Output Current Short (Circuit			~-100	mA	V _{CC} - MAX, Note 7
	0	93425XC	2.4			V	I _{OH} = -10.3 mA, V _{CC} 5.0 V · 5%
۷он	Output HIGH Voltage	93425XM	2.4			٧	I _{OH} = -5.2 mA
VCD	Input Diode Clamp Vol	tage	1	-1.0	-1.5	_^_ :-	V _{CC} - MAX, I _{IN} = +10 mA
			T	1	130	mΑ	TA > 75°C
1CC	. Power Supply Current			95	155	mΑ	TA = 0°C All Inputs Grounded
	l				170	mA	TA = -55°C

AC CHARACTERISTICS: Over Guaranteed Operating Ranges (Notes 1, 2, 4, 5, 6)

	Į Į	9:	3425A	XC.	93425XC		9	3425×	M.			
SYMBOL /	CHARACTERISTIC	MIN	TYP	MAX	MIN	TYP	MAX	MIN	TYP	MAX	UNIT\$	CONDITIONS
		! '	Note.	31		(Note :	31	1	Note 3	3)		
READ MODE	DELAY TIMES	!										
¹ ACS	, Chip Select Time		15	30		15	40		15	45	ns	See Test Circuit
†ZRCS	Chip Select to HIGH Z	l	15	30		20	40		20	50	""	and Waveforms
†AA	Address Access Time	l	30	45		40	70		40	75		
WRITE MODE	DELAY TIMES	Ţ~~ ~										
^t ZWS	Write Disable to HIGH Z	į	20	30		20	40		20	45	ns	
twa .	, Write Recovery Time	!	20	40		25	50		45	55	l ns	
	INPUT TIMING REQUIREMENTS											
tw	Write Pulse Width (to guarantee write)	30	25		50	25		55	25			 See Test Circuit
twso	Data Set-Up Time Prior to Write	5	0		5	0		5	0			
(WHD	Data Hold Time After Write	5	0		5	0		5	0			į anu wavelornis į
twsa	Address Set-Up Time	10	0		15	0		15	0		ns	
WHA	Address Hold Time	5	0		5	0		5	0		•	
twscs	Chip Select Set-Up Time	5	0		5	0		5	0			
^t WHCS	Chip Select Hold Time	5	0		5	0		5	0			
c ₁	Input Pin Capacitance		4	5		4	5		4	5	ρF	Measure with
co	Output Pin Capacitance		7	8		7	8		7	8 .	Pr	Pulse Technique

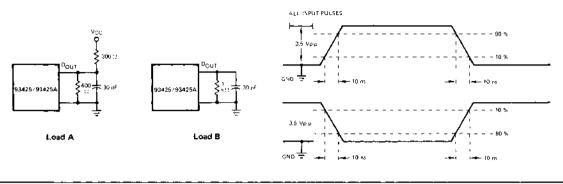
NOTES:

- 1. Conditions for testing not shown in the Table, are chosen to guarantee operation under: "worst case" conditions
- The specified LIMITS represents the "worst case" value for the parameters. Since these "worst case" values normally occur at the temperature and supply voltage extremes, additional noise immunity and guard banding can be achieved by decreasing the allowable system operating ranges.
- Typical values are at V_{CC} = 5.0 V, T_{Δ} = 25 $^{\circ}$ C, and MAX loading.
- 4 The Temperature Ranges are guaranteed with transverse air flow exceeding 400 linear feet per milliote and a two minute warm up. Typical thermal resistance values of the package at maximum temperature are.
 - HyA (Junction to Ambient) (at 400 fpm air flow) = 50, C/Watt, Ceramic DIP: 65, C/Watt, Plastic OIP, NA, Flatpak.
 - $\theta_{
 m JA}$ (Junction to Ambient) (still air) = 90 C/Watt, Ceramic DIP 110 C/Watt, Plastic DIP, NA, Flatpak, $\theta_{
 m JC}$ (Junction to Case) = 25 C/Watt, Ceramic DIP: 25 C/Watt Plastic DIP, 10 C/Watt, Flatpak.
- The MAIX address access time is guaranteed to be the "worst case" bit in the memory using a pseudo random testing pattern.
 tw measured at two = MIN, two measured at tw = MIN.
 Duration of short circuit should not exceed one second.

AC TEST LOAD AND WAVEFORMS

LOADING CONDITIONS

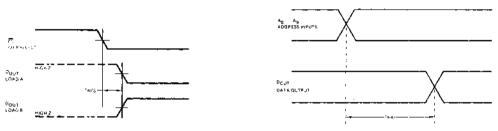
INPUT PULSES



READ MODE

PROPAGATION DELAY FROM CHIP SELECT

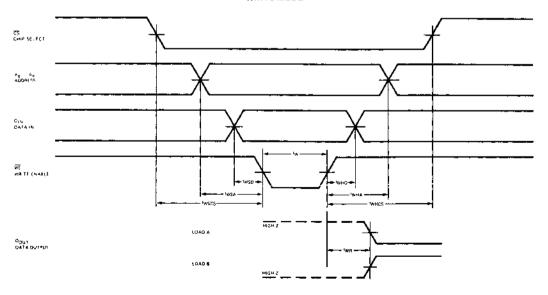
PROPAGATION DELAY FROM ADDRESS INPUTS



(All time measurements referenced to 1.5 V)

AC WAVEFORMS (Cont'd)

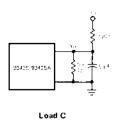
WHITE MODE

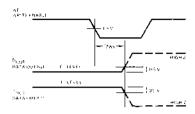


(All above measurements referenced to 1.5 V)

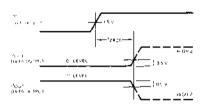


WRITE ENABLE TO HIGH Z DELAY





PROPAGATION DELAY FROM CHIP SELECT TO HIGH Z



(All τ_{ZXXX} parameters are measured at a delta of 0.5 V from the logic level and using Load C,I



FUNCTIONAL DESCRIPTION

The R29600 and the R29601 are electrically programmable Low-Power Schottky TTL read only memories. Both devices are organized as 256 words of 8 bits each; the R29600 has open collector outputs and the R29601 has three-state outputs. The devices are shipped with all bits HIGH and each bit in the memory can be programmed to a LOW by applying appropriate voltages to the circuit. At each bit location on the circuit there is a narrow link of fuse material which is conductive, but which can be opened by passing a short high-current pulse through it.

The programming voltage is applied at the output pin for the bit to be programmed, while the word to be programmed is selected by normal TTL levels on the address lines. The passage of current through the link is controlled by a programming pulse on the chip select input.

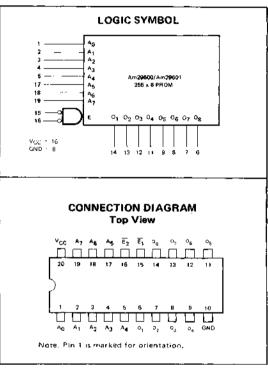
After programming, the device can be used for microprogram storage or random logic function generation, like any read-only memory. If either chip select input is held HIGH, the outputs will all turn off, so the outputs of several memories can be tied together for expansion.

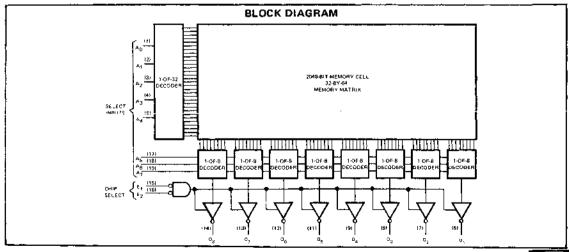
ORDERING INFORMATION

٠.		
Package	Temperature	Order
Type	Range	Number
	Open Collectors	
Hermetic DIP	0°C to +75°C	R29600DC
Hermetic DIP	-55°C to +125°C	R29600DM
Hermetic Flat Pak	-55°C to +125°C	R29600FM
	Three-State Outputs	
Hermetic DIP	0°C to +75°C	R29601DC
Hermetic DIP	-55°C to +125°C	R29601DM
Hermetic Flat Pak	-55°C to +125°C	R29601FM

DISTINCTIVE CHARACTERISTICS

- Field programmable read only memory
- Highly reliable nichrome fuses
- Pin compatible with other popular 256 by 8 PROMS.
- Typical fusing time of 95µs/bit
- 70ns access time at 25°C
- Three-state and open-collector versions





MAXIMUM RATINGS (Above which the useful life may be impaired)

Storage Temperature	-65°C to +150°C
Temperature (Ambient) Under Bias	−55°C to +125°C
Supply Voltage to Ground Potential (Pin 20 to Pin 10) Continuous	−0.5V to +7V
DC Voltage Applied to Outputs (Except During Programming)	$-0.5V$ to $+V_{CC}$ max.
DC Input Voltage (Address Inputs)	-0.5V to +5.5V
DC Voltage Applied to Outputs During Programming	26V
Output Current into Outputs During Programming	125 mA
DC Input Voltage (Chip Select Input - Pin)	-0.5V to +33V
DC Input Current	-30mA to +5mA

OPERATING RANGE

R29600XC, R29601XC	T _A = 0°C to +75°C	V _{CC} = 5.0 V ± 5%	COM'L
R29600XM, R29601XM	$T_A = -55^{\circ}C \text{ to } +125^{\circ}C$	V _{CC} = 5.0 V ±10%	MIL

ELECTRICAL CHARACTERISTICS OVER OPERATING RANGE (Unless Otherwise Noted)

Parameters	Description	Test Canditions	Min.	Τγ ρ. {Note 1}	Max.	Units
V _{OH} (R29601 Onl	Output HIGH Voltage	V _{CC} = MIN., I _{OH} = -2.0mA V _{IN} = V _{IH} or V _{IL}	2.4			Volts
v ol	Output LOW Voltage	V _{CC} = MIN. I _{OL} = 8mA V _{IN} = V _{IH} or V _{IL} I _{OL} = 16mA			0.4 0.45	Volts
v _{IH}	Input HIGH Level	Guaranteed input logical HIGH voltage for all inputs	2.0			Volts
v _{IL}	Input LOW Level	Guaranteed input logical LOW voltage for all inputs			0.8	Volts
1 ₁ L	Input LOW Current	V _{CC} = MAX., V _{IN} = 0.4V		60	-250	μА
14	Input HIGH Current	V _{CC} = MAX., V _{IN} = 2.7V			10	μΑ
1,	Input HIGH Current	V _{CC} = MAX., V _{IN} = 5.5V			1.0	mA
1 ₍ (Pin 15 Only)	input HIGH Current	V _{CC} = MAX., V _{IN} = 4.5V			1.0	mA
I _{SC} (R29601 Only	Output Short Circuit	V _{CC} = MAX., V _{OUT} = 0.0V	-12	-3 5	-85	mΑ
¹cc	Power Supply Current	All inputs = GND V _{CC} = MAX.	f	90	130	mA
v _I	Input Clamp Voltage	V _{CC} = MIN., I _N = 18mA			-1.5	V
CEX	Output Leakage Current	$V_{CC} = MAX$ $V_{CS} = 2.4V$ $V_{O} = 4.5V$ $V_{O} = 2.4V$ $V_{O} = 0.4V$			100 40 40	μα

Note 1. Typical limits are at $V_{CC} = 5.0V$ and $T_{A} = 25^{\circ}C$

^{2.} Not more than one output should be shorted at a time, Duration of the short circuit should not be more than one second.

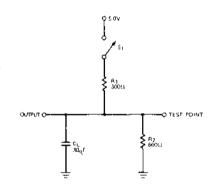


SWITCHING CHARACTERISTICS OVER OPERATING RANGE

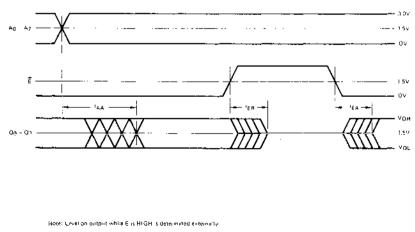
			Тур.				
Parameter	Description	Test Conditions	5V 25°C	25°C	Çom'l	Mil	Units
†AA	Address Access Time	C _L = 30 pF R _L = 300 Ω to V _{CC}	50				ns
tEA	Enable Access Time	and 600 Ω to GND	20				ns
tER	Enable Recovery Time	(16 mA Load) Note 1	20			_	ns.

Note 1, 300 Ω resistor opened for t_{EA} and t_{EA} measurements between HIGH and OFF states.

AC TEST CIRCUIT



SWITCHING WAVEFORMS



KEY TO TIMING DIAGRAM

WAVEFORM	INPUTS	OUTPUTS
	MUST B€ STEADY	WILL BE STEADY
	MAY CHANGE	WIEL BE CHANGING FROM H TO L
	MAY CHANGE FROM L TO H	WILL BE CHANGING FROM L TO M
XXX	DON'T CARE. ANY CHANGE PERMITTED	CHANGING. STATE UNKNOWN
}}	OQES NOT APPLY	CENTER LINE IS HIGH IMPEDANCE "OFF" STATE



DEVICE DESCRIPTION

ne device is manufactured with all outputs high in all storage locations. To make an output low at a particular word, a nifrome fusible link must be changed from a low resistance to a high resistance. This procedure is called programming. There be 1024 fusible links on the chip. Programming equipment can be obtained from Data I/O, Inc., and other manufacturers of FROM programming equipment.

PROGRAMMING DESCRIPTION

a select a particular fusible link for programming, the word address is presented with TTL levels on A_0 through A_7 , a V_{CC} 5.50 V is applied or left applied, and the program pin (Enable E_1) and the output to be programmed are taken to an evated voltage to supply the required current to program the fuse. The outputs must be programmed one output at a time, noe internal decoding circuitry is capable of sinking only one unit of programming current at a time.

ENABLE E2

nable E_2 (pin 16) is a logic enable and is not used during programming. It may be high, low or open during programming, then checking that an output is programmed (which is called verification) enables E_1 and E_2 must be low to activate the exice. Since E_2 must be low during verification and the state is irrelevant during programming, the simplest procedure is to round E_2 during programming and verification.

TIMING

he programming procedure involves the use of the program pin (an enable) and the output pin. In order to guarantee that is output transistor is off before increasing the voltage on the output pin, the program pin's voltage pulse must come before the output pin's programming pulse. 100 his delay is adequate. The programming pulse applied to the output pin and program pin must have a 50 to 70 microseconds rise time. See Figure 4.

VERIFICATION

fter programming a device, it can be checked for a low output by taking both enables low. Since we must guarantee operator at minimum and maximum V_{CC} , current and temperature, the device must be required to sink 12 mA at 4.20 V V_{CC} and 0.2 mA at 6.0 V V_{CC} at room temperature.

BOARD PROGRAMMING

nits may be programmed at the board level by bringing the program pin of each package to the card connector. To program particular package "A", the program pin of package A and one output of package A, which may or may not be "OR" tied other packages, are taken to the required programming voltage. An alternate procedure is to tie the enable and outputs gether as required by the system function and only apply V_{CC} to the device to be programmed. The number of units aldered on a board should be consistant with expected programming yields to avoid rework.

UNPROGRAMMABLE UNITS

isual inspection at 200X prior to encapsulation, test fuses and decoding circuitry tests are used to guarantee a high proamming yield of the device in the field. However, because of random defects, it is impossible to guarantee that a link will ben without actually programming it. UNITS RETURNED TO RAYTHEON AS UNPROGRAMMABLE MUST BE AC-DMPANIED BY A COMPLETE DEVICE TRUTH TABLE WITH THE LOCATION WHICH COULDN'T BE PRO-RAMMED, OR WHICH FALSELY PROGRAMMED, CLEARLY INDICATED.

ROGRAMMING SPEED

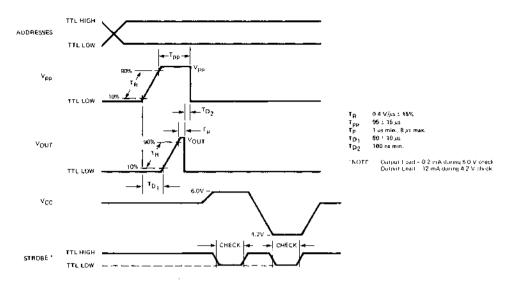
pically fuses will blow on the rise time of the pulse.

automated programmers which must copy devices in a short time because of production requirements, the following pulse d voltage sequences have been found to maximize reliability, programming yield, and thruput. The device should be veried after each programming attempt and is advanced to the next bit if the device has programmed.

PULSE NUMBER	PROGRAM PIN VOLTAGE	OUTPUT VOLTAGE
1 to 3	27 V	20 V
4 to 6	30 V	23 V
7 to 9	33 V	26 V



PROGRAMMING TIMING



REQUIRED INFORMATION FOR RAYTHEON TO PROGRAM TO YOUR TRUTH TABLE

TRUTH TABLES

Raytheon can program devices at our facility from Raytheon truth table forms (available on request). For customers desiring to make their own forms, an example is shown below:

		OUT	PUTS					
WORD	PIN> 14	13	12	11	9	8	7	6
NUMBER	o ₈	07	06	05	$\mathbf{o_4}$	o^3	02	01
a	н	н	н	L	н	L	L	н
1	L	н	L.	н	H	Н	H	Ł
•	•	•		•		•	•	•
•	•	•	•	•	•	•	-	•
255	L	Н	н	н	н	H	Н	н

Note: A high voltage on the data out lines is signified by an "H". A low voltage on the data out lines is signified by an "L". The word number assumes positive logic on the address pins, so for example, word 255 = LHHHHHHH.

PAPER TAPE FORMAT

Truth tables can also be sent to Raytheon in an ASCII tape format. Information can be sent to us by air mail or TWX 910-379-6481. The tape reading equipment at Raytheon only recognizes ASCII characters 8, P, N, and F and interprets them respectively as Start, Begin a word, High data, Low data, Finish a word, and End of tape. All other characters such as carriage returns, line feeds, etc. are ignored so that comments and spaces may be sent in the data field to improve readability. Comments, however, should not use the characters 8, H, L, F. Word addresses must begin with zero and count sequentially to word 255.

In order to assist the machine operator in determining where the heading information stops and the data field begins, 25 bell characters or rubout characters should precede the start of the truth table. Any type of 8 level paper tape (mylar, fanfold, etc.) is acceptable. Channel 1 is the most significant bit and channel 8 (parity) is ignored. Sprocket holes are located between channels 3 and 4. Note that the order of the outputs between characters B and F is 08, 07, 06, 05, 04, 03, 02, 01, not 01, 02, 03, 04, 05, 06, 07, 08.

A typical list of characters and their machine interpretations are shown in the next page.



SUGGESTED IMPLEMENTATION OF THE VERIFICATION CIRCUITRY

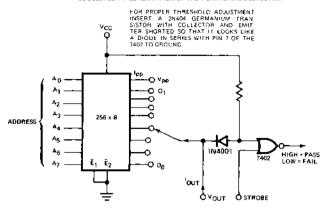
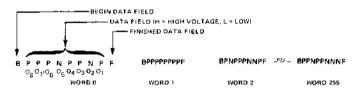


FIGURE 5. THE 1N4001 DIODE PROTECTS THE INPUT OF THE 7402 FROM THE HIGH PROGRAMMING VOLTAGES



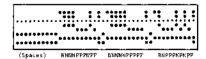
The required beading information at the biginning of the tape is as follows:

CUSTOMERS NAME AND PHONE. THUTH TABLE NUMBER -CUSTOMERS TWX NUMBER .. NUMBER OF TRUTH TABLES -PURCHASE ORDSA NUMBER ___ TOTAL NUMBER OF PARTS ... BAYTHEON FART NUMBER ______ NUMBER OF PARTS OF EACH TRUTH TABLE CUSTOMER SYMBOUIZED PART NUMBER _ 25 SELL OR RUDOUT CHARACTERS _

An example is shown below:

| BLARNEY ELECTRONICS | 408-735-8140 | TWX | 911-336-9225 | F0142 | BYNNPPPNPE BS | 2960! | BYNNNNNNP BN

B LEVEL TWX.



Interpretation:



PROGRAMMING PARAMETERS - Do Not Test These Limits or You May Program the Device

ľ			}			
				1		
SYMBOL	PARAMETERS	TEST CONDITION See Figure 4	MIN.	OR OPTIMUM	MAX.	UNITS
lbb	Current Program Pin During Programming, Before and After Fuse Has Blown	V _{CC} = 5.50V V _{out} = 5.0V to 25V V _{pp} = 4.50V		0		mA
		V _{pp} = 27V		77		mA
J _{out}	Current into Output During Programming Before the Fuse Has Programmed	V _{pp} = 27V, V _{CC} = 5,50V V _{out} = 9.0V		0.1		mA
		V _{out} = 20V	<u> </u>	16		mA
lout	Current into Output During Programming After the Fuse Has Programmed	V _{pp} = 27V V _{out} =20V V _{CC} = 5.50V		0.1		mA
TR	Rise Time of Program Pulse Applied to the Data Out or Program Pin		0.34	0.4	0.46	V/μs
VCCP	V _{CC} Required During Programmin	g	5.40	5.50	5.60	V
'oLV ₁	Output Current Required During Verification	Both Chip Enables Low T _A = 25°C, V _{CC} = 4.2V	11	12	13	mA
OLV2	Output Current Required During Verification	Both Chip Enables Low T _A = 25°C, V _{CC} = 6.0V	0.19	0,2	0.21	mA
MDC	Maximum Duty Cycle During Automatic Programming of Program Pin and Output Pin	T _P ⊤C			25	%
v _{pp}	Required Programming Voltage on Program Pin		27	27	33	V
V _{out}	Required Programming Voltage on the Output Pin		20	20	26	v
ı _L	Required Current Limit of the Power Supply Feeding the Program Pin and the Output During Programming	V _{pp} = 33V V _{out} = 26V V _{CC} = 5.50V	150			mA
T _{pp}	Required Coincidence Among the Program Pin, Output, Address and V _{CC} for Programming		80	95	110	įis
T _{D1}	Required Time Delay Between Disabling the Memory Output and Application of the Output Program ming Pulse	Measure at 10% Levels	70	во	90	μs
т _{О2}	Required Time Delay Between Removal of Programming Pulse and Enabling the Memory Output	Measure at 10% Levels	100			ns



FUNCTIONAL DESCRIPTION

The R29660 and R29661 are electrically programmable Low-Power Schottky TTL read only memories. Both devices are organized as 256 words of 4 bits each: the R29660 has open collector outputs and the R29661 has three-state outputs. The devices are shipped with all bits HIGH and each bit in the memory can be programmed to a LOW by applying appropriate voltages to the circuit. At each bit location on the circuit there is a narrow link of fuse material which is conductive, but which can be opened by passing a short, high-current pulse through it.

The programming voltage is applied at the output pin for the bit to be programmed, while the word to be programmed is selected by normal TTL levels on the address lines. The passage of current through the link is controlled by a programming pulse on the chip select input.

After programming, the device can be used for microprogram storage or random logic function generation, like any read-only memory. If either chip select input is held HIGH, the outputs will all turn off, so the outputs of several memories can be tied together for expansion.

DISTINCTIVE CHARACTERISTICS

- Field programmable read only memory
- Highly reliable nichrome fuses
- Pin compatible with other popular 256 by 4 PROMS
- Typical fusing time of 95μs/bit
- 60ns access time at 25°C

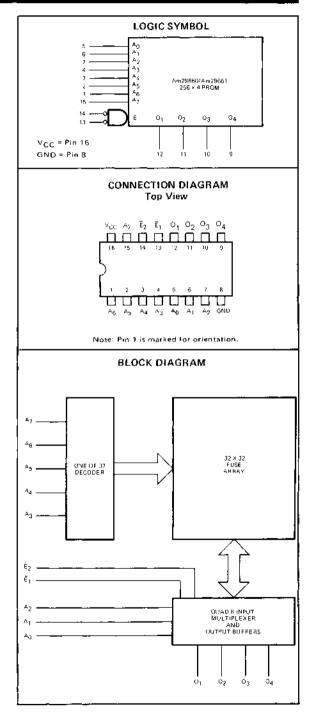
Three-state and open-collector versions

ORDERING INFORMATION

Package Type	Temperature Range	Order Number
	Open Collectors	
Hermetic DIP	0°C to +75°C	R29660DC
Hermetic DIP	-55°C to +125°C	R29660DM
Hermetic Flat Pak	-55°C to +125°C	R29660FM

Three-State Outputs

Hermetic DIP	0°C to +75°C	R29661DC
Hermetic DIP	-55°C to +125°C	R29661DM
Hermetic Flat Pak	-55°C to +125°C	R29661FM





MAXIMUM RATINGS (Above which the useful life may be impaired)

Storage Temperature	65°C to +150°C
Temperature (Ambient) Under Bias	55°C to +125°C
Supply Voltage to Ground Potential (Pin 16 to Pin 8) Continuous	-0.5V to +7V
DC Voltage Applied to Outputs (Except During Programming)	-0.5V to +V _{CC} max.
DC Input Voltage (Address Inputs)	-0,5V to +5,5V
DC Voltage Applied to Outputs During Programming	26V
Output Current into Outputs During Programming	125 mA
DC Input Voltage (Chip Select Input – Pin 13)	-0.5V to +33V
DC Input Current	-30mA to +5mA

OPERATING RANGE

R29660XC, R29661XC	T _A = 0°C to +75°C	V _{CC} = 5.0 V ±5%	COM'L
R29660XM, R29661XM	T _A = -55°C to +125°C	V _{CC} = 5.0 V ±10%	MIL

ELECTRICAL CHARACTERISTICS OVER OPERATING RANGE (Unless Otherwise Noted)

Parameters	Description	Test Conditions	Min.	Typ. (Note 1)	Max.	Units
V _{OH} (Am29661 Onl	Output HIGH Voltage y)	V _{CC} = MIN., I _{OH} = -2.0mA V _{IN} = V _{IH} or V _{IL}	2.4			Volts
v _{OL}	Output LOW Voltage	V _{CC} = MIN. OL = 8mA V _{IN} = V _{IH} or V _{IL} I _{OL} = 16mA			0.4 0.45	Volts
V _{IH}	Input HIGH Level	Guaranteed input logical HIGH voltage for all inputs	2.0			Volts
v _{iL}	Input LOW Level	Guaranteed input logical LOW voltage for all inputs			8.0	Volts
1 _{IL}	Input LOW Current	V _{CC} = MAX., V _{IN} = 0.4V		-60	-250	μΑ
1,	Input HIGH Current	V _{CC} = MAX., V _{IN} = 2.7V			10	μΑ
1	Input HIGH Current	V _{CC} * MAX., V _{IN} = 5.5∨			1.0	mA
l ₁ (Pin 13 Only)	Input HIGH Current	V _{CC} = MAX., V _{IN} = 4.5V			1.0	mA
^I SC (Am29 6 61)	Output Short Circuit Current	V _{CC} = MAX., V _{OUT} = 0.0V	-12	-35	-85	mA
Icc	Power Supply Current	All inputs = GND V _{CC} = MAX.		90	130	mA,
v _i	Input Clamp Voltage	V _{CC} = MIN., I _{IN} = 18mA			-1.5	٧
CEX	Output Leakage Current	$V_{CC} = MAX$ $V_{O} = 4.5V$ $V_{O} = 2.4V$ $V_{O} = 0.4V$			100 40 -40	μΑ

Note 1. Typical limits are at V_{CC} = 5.0V and T_A = 25 $^{\circ}C$

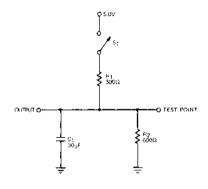
^{2.} Not more than one output should be shorted at a time. Duration of the short circuit should not be more than one second.

SWITCHING CHARACTERISTICS OVER OPERATING RANGE

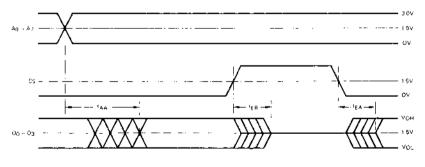
			Тур.		Max.		
Parameter	Description	Test Conditions	5∨ 25°C	25°C	Comʻl	Mil	Units
tAA	Address Access Time	C _L = 30 pF R ₁ = 300 Ω to V _{CC}	45	60	70	80	nş
t _{EA}	Enable Access Time	end 600 Ω to GND	20	30	35	40	nş
tER	Enable Recovery Time	(16 mA Load) Note 1	20	30	35	40	ns

Note 1, 300 $\!\Omega$ resistor opened for $t_{\mbox{EA}}$ and $t_{\mbox{EB}}$ measurements between HIGH and OFF states.

AC TEST CIRCUIT



SWITCHING WAVEFORMS



Note: Level on output while E is HIGH is determined externally.

KEY TO TIMING DIAGRAM

WAVEFORM	MPUTS	(Ic)TPUTS
	MUST BE STEADY	WILL BE STEADY
	MAY CHANGE FROM HITO L	WILL BE CHANGING FROM HITO C
	MAY CHANGE FROM J MORA	WILL BE CHANGING FROM I, TO H
XXXX	DON'T CARE, ANY CHANGE PERMITTED	CHANGING STATE UNKNOWN
₩	OGES NOT APPLY	CENTER LINE IS HIGH IMPEDANCE IDEFT STATE
,		



PROGRAMMING INSTRUCTIONS

1) DEVICE DESCRIPTION

The device is manufactured with all outputs high in all storage locations. To make an output low at a particular word, a nichrome fusible link must be changed from a low resistance to a high resistance. This procedure is called programming. There are 1024 fusible links on the chip. Programming equipment can be obtained from Data I/O, Inc., and other manufacturers of P/ROM programming equipment.

2) PROGRAMMING DESCRIPTION

To select a particular fusible link for programming, the word address is presented with TTL levels on A_0 through A_7 , a V_{CC} of 5.50 V is applied or left applied, and the program pin (Enable E_1) and the output to be programmed are taken to an elevated voltage to supply the required current to program the fuse. The outputs must be programmed one output at a time, since internal decoding circuitry is capable of sinking only one unit of programming current at a time.

3) ENABLE E2

Enable E₂ (pin 14) is a logic enable and is not used during programming. It may be high, low or open during programming. When checking that an output is programmed (which is called verification) enables E₁ and E₂ must be low to activate the device. Since E₂ must be low during verification and the state is irrelevant during programming, the simplest procedure is to ground E₂ during programming and verification.

4) TIMING

The programming procedure involves the use of the program pin (an enable) and the output pin. In order to guarantee that the output transistor is off before increasing the voltage on the output pin, the program pin's voltage pulse must come before the output pin's programming pulse and leave after the output pin's programming pulse. 100 ns delay is adequate. The programming pulse applied to the output pin and program pin must have a 50 to 70 microseconds rise time. See Flaure 4.

VERIFICATION

After programming a device, it can be checked for a low output by taking both enables low. Since we must guarantee operation at minimum and maximum V_{CC} , current and temperature, the device must be required to sink 12 mA at 4.20 V V_{CC} and 0.2 mA at 6.0 V V_{CC} at room temperature.

6) BOARD PROGRAMMING

Units may be programmed at the board level by bringing the program pin of each package to the card connector. To program a particular package "A", the program pin of package A and one output of package A, which may or may not be "OR" tied to other packages, are taken to the required programming voltage. An alternate procedure is to tie the enable and outputs together as required by the system function and only apply V_{CC} to the device to be programmed. The number of units soldered on a board should be consistent with expected programming yields to avoid rework.

7) UNPROGRAMMABLE UNITS

Visual inspection at 200X prior to encapsulation, test fuses and decoding circuitry tests are used to guarantee a high programming yield of the device in the field. However, because of random defects, it is impossible to guarantee that a link will open without actually programming it. UNITS RETURNED TO RAYTHEON AS UNPROGRAMMABLE MUST BE ACCOMPANIED BY A COMPLETE DEVICE TRUTH TABLE WITH THE LOCATION WHICH COULDN'T BE PROGRAMMED, OR WHICH FALSELY PROGRAMMED, CLEARLY INDICATED.

PROGRAMMING SPEED

Typically fuses will blow on the rise time of the pulse.

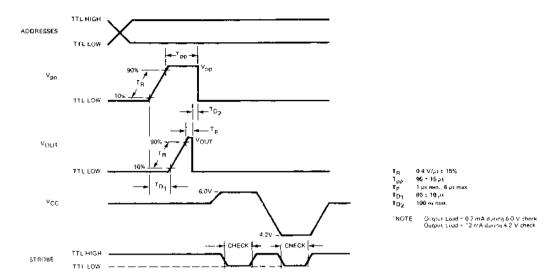
P

In automated programmers which must copy devices in a short time because of production requirements, the following pulse and voltage sequences have been found to maximize reliability, programming yield, and thruput. The device should be verified after each programming attempt and is advanced to the next bit if the device has programmed.

PULSE NUMBER	PROGRAM PIN VOLTAGE	OUTPUT VOLTAGE
1 to 3	27 V	20 V
4 to 6	30 V	23 V
7 to 9	33 V	26 V



PROGRAM TIMING



TRUTH TABLES

Raytheon can program devices at our facility from Raytheon truth table forms (available on request). For customers desiring to make their own forms, an example is shown below:

				DUTPUTS	
WORD	PIN 💳	9	10	11	12
NUMBER		04	03	02	Q ₁
0		·H	Н	Н	L,
1		L	н	L	Н
•		-			-
•					-
255		L	Н	Н	н

Note: A high voltage on the data out lines is signified by an "H". A low voltage on the data out lines is signified by an "L". The word number assumes positive logic on the address pins, so for example, word 255 = HHMHHHHH.

PAPER TAPE FORMAT

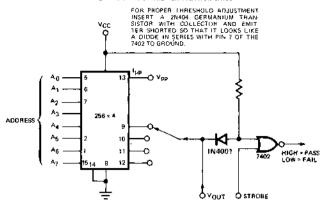
Truth tables can also be sent to Raytheon in an ASCII tape format. Information can be sent to us by air mail or TWX 910-379-6481. The tape reading equipment at Raytheon only recognizes ASCII characters S, B H, L, F and E and interprets them respectively as Start, Begin a word, High data, Low data, Finish a word, and End of tape. All other characters such as carriage returns, line feeds, etc. are ignored so that comments and spaces may be sent in the data field to improve readability. Comments, however, should not use the characters S, B, H, L, F, E. Word addresses must begin with zero and count sequentially to word 255.

In order to assist the machine operator in determining where the heading information stops and the data field begins, 25 bell characters or rubout characters should precede the start of the truth table. Any type of 8 level paper tape (mylar, fanfold, etc.) is acceptable. Channel 1 is the most significant bit and channel 8 (parity) is ignored. Sprocket holes are located between channels 3 and 4. Note that the order of the outputs between characters B and F is 04, 03, 02, 01, not 01, 02, 03, 04.

A typical list of characters and their machine interpretations are shown on the next page.

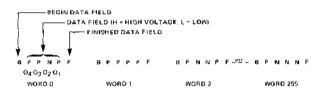


SUGGESTED IMPLEMENTATION OF THE VERIFICATION CIRCUITRY



THE 1N4001 DIODE PROTECTS THE INPUT OF THE 7402 FROM THE HIGH PROGRAMMING VOLTAGES

FIGURE 5.
THE 1N4001 DIDDE PROTECTS THE INPUT OF THE 7402 FROM THE HIGH PROGRAMMING VOL.) AGES



The required heading information at the beginning of the tape is as follows

CUSYOMERS NAME AND PHONE	TRUTH TABLE NUMBER
CUSTOMERS TWX NUMBER	NUMBER OF TRUTH TABLES
PUHCHASE ORDER NUMBER	701AL NUMBER OF PARTS
RAYTHEON PART NUMBER	NUMBER OF PARTS OF EACH TRUTH TABLE
CHSTOMATS Sykadyl, 1251) SwBT by Backet	OC CELL ON BURGHAY CHARACTERS

An example is shown below:

BLARNEY ELECTRONICS 408 735 H140

TWX 911-338-9225

 PO142
 BX:NNPF
 BNNNNF
 BNPPF
 12 8 level 1 TWX

3





PROGRAMMING PARAMETERS — Do Not Test These Limits or You May Program the Device

	PARAMETERS		<u> </u>	LIMITS		
SYMBOL		TEST CONDITION See Figure 4	MIN.	TYPICAL OR OPTIMUM	MAX.	UNITS
ℓ _{pp}	Current Program Pin During Programming, Before and After Fuse Has Blown	V _{CC} = 5.50V V _{out} = 5.0V to 25V V _{pp} = 4.50V	:	0		mA
		V _{pp} = 27∨		77		mA
lout	Current into Output During Programming Before the Fuse	V _{pp} = 27V, V _{CC} = 5.50V V _{out} = 9.0V		0,1		mA
	Has Programmed	V _{out} = 20V		16		mA
lout	Current into Output During Programming After the Fuse Has Programmed	V _{pp} = 27V V _{out} =20V V _{CC} = 5.50V		0.1		mA
т _R	Rise Time of Program Pulse Applied to the Data Out or Program Pin		0.34	0.4	0.46	V/µs
v _{CCP}	V _{CC} Required During Programming		5.40	5.50	5.60	v
I _{OLV1}	Output Current Required During Verification	Both Chip Enables Low T _A = 25°C, V _{CC} = 4.2V	11	12	13	mA
I _{OLV2}	Output Current Required During Verification	Both Chip Enables Low T _A = 25°C, V _{CC} = 6.0V	0.19	0.2	0.21	mA
MDC	Maximum Duty Cycle During Automatic Programming of Program Pin and Output Pin	^Т Р — Тс	İ		25	96
V _{pp}	Required Programming Voltage on Program Pin		27	27	33	V
Vout	Required Programming Voltage on the Output Pin		20	20	26	V
ı	Required Current Limit of the Power Supply Feeding the Program Pin and the Output During Programming	Vpp = 33V Vout = 26V VCC = 5.50V	150			mA
т _{рр}	Required Coincidence Among the Program Pin, Output, Address and V _{CC} for Programming		80	95	110	μs
т _{D1}	Required Time Delay Between Disabling the Memory Output and Application of the Output Programming Pulse	Measure at 10% Levels	70	80	90	μs
т _{D2}	Required Time Delay Between Removal of Programming Pulse and Enabling the Memory Output	Measure at 10% Levels	100			ns



FUNCTIONAL DESCRIPTION

The R29690 and R29691 are field programmable random logic arrays. The R29691 has three-state outputs and the R29690 has open collector outputs. The devices have 16 data inputs and contain 48 intermediate product terms. Random logic functions are implemented by using the device like 8 large AND-OR-INVENT gates. Up to 48 AND functions may be generated in the device, where each AND function is the product of any or all of the 16 inputs (I_0-I_{15}) or their complements. Any of the AND functions which have been generated may then be OR'ed to form an output function on one of the F outputs. The F output may also be programmed to invert, i.e., form an AND-OR-INVERT function, rather than AND-OR-

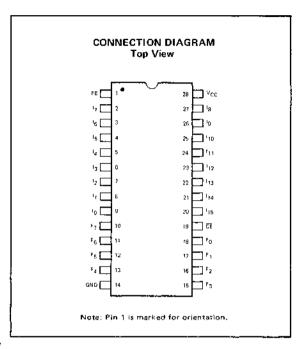
Units are shipped with all fuses intact, and all outputs LOW. Each of the 48 AND functions contains both true and complement of each of the 16 input variables. (Therefore all AND functions are initially false). A variable (or its complement) is inserted into an AND gate by blowing the fuse for the undesired state. A variable is removed from an AND gate by blowing both the true and complement fuses. Each OR function contains the outputs of all 48 AND functions. AND functions are removed from the OR gates by blowing the proper fuse. If the output is to be negated, then a fuse is blown on an EXCLUSIVE-OR gate at the output. That output then becomes an AND-OR-INVERT function of the inputs. The chip enable input (CE) forces all outputs OFF when HIGH. When LOW, the outputs will be OFF or LOW for the R29690, or HIGH or LOW for the R29691

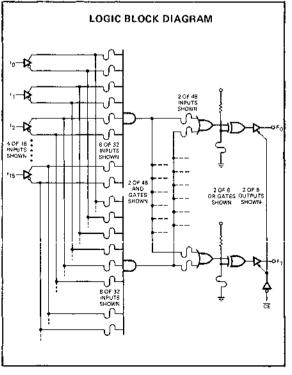
DISTINCTIVE CHARACTERISTICS

- LSI replacement for random logic and inefficiently used ROMs and PROMs
- 16 inputs 8 outputs 48 product terms
- Logic equation for each output field programmed by fusing nichrome links
- Each output can be programmed to be active HIGH or active LOW
- 100% processing in accordance with MIL-STD-883
- Three-state or open collector outputs controlled by active LOW chip enable.

ORDERING INFORMATION

This data is based on design goals and is subject to change following complete characterization.







OPERATING RANGE

Part No.	Ambient Temperature	Vc c
R29690DC, R29691DC	$T_A = 0^{\circ} C \text{ to } +75^{\circ} C$	5.0∨ ±5%
R29690DM, R29691DM	T _A = 55°C to +125°C	5.0V ±10%



ELECTRICAL CHARACTERISTICS OVER OPERATING RANGE (Unless Otherwise Noted)

Parameters	Description	Test Co	Test Conditions		Typ. (Note 1)	Máx.	Units
V _{OH} (RA29691 Only)	Output HIGH Voltage	V _{CC} = MIN., I _{OH} = −2.0mA V _{IN} = V _{IH} or V _{IL}		2.4			Volts
VoL	Output LOW Voltage	V _{CC} = MIN., I _{OL} = 12mA V _{IN} = V _{IH} or V _{IL}			0.3	0,45	Volts
VIH	Input HIGH Level	Guaranteed input logical HIGH voltage for all inputs		2.0			Volts
VIL	Input LOW Level	Guaranteed input logical LOW voltage for all inputs				0.8	Volts
liL	Input LOW Current	V _{CC} = MAX., V _{IN} = 0.45V			-0.03	-0,25	mΑ
hн	Input HIGH Current	V _{CC} = MAX., V _{IN} = 2.4V			<1	25	μА
Isc (Am29791 Only)	Output Short Circuit Current	V _{CC} = MAX., V _{QUT} = 0.0V		-12	-35	-90	mΑ
too	Power Supply Current	CE = GND, all other inputs = 4.5V VCC = MAX.			120		mA
v _c	Input Clamp Voltage	V _{CC} = MIN., I _{IN} = -18mA				-1,2	Volts
	Outro A Landana Outro	U = - 2 AU	ΛO = ΛCC	}		40	
CEX, OLK	Output Leakage Current	∨ <u>CE</u> = 2.4∨	V _Q = 0.45V	-40			μΑ

Note 1. Typical limits are at V_{CC} = 5.0V and T_A = 25°C.

SWITCHING CHARACTERISTICS OVER OPERATING RANGE

RA29690DC

RA29690DM RA29691DM

Parameters	Description	Test Conditions	Min.	Тур,	Max.	Mia,	Тур.	Max.	Units
tpZL(CE)	Delay Chip Select to Output Active HIGH or LOW	C _L = 30pF R _L = 300Ω to V _{CC} and 600Ω to GND (16mA Load) Note 2		25			40	·	ns
tPLZ(CE)	Delay Chip Select H1GH to Output OFF			25			40		ns
tpLH(f)	Delay Input to Output HIGH Delay Input to Output LOW			30		"	50		nş :

Note 2. 300 Ω resistor opened for tops and ten measurements between HIGH and OFF states,

PROGRAMMING PROCEDURE

Programming is facilitated by two internal 1-of-48 decoders used to select one of the 48 product terms. One of these decoders is activated by applying high-voltage logic to inputs l_0 - l_6 , and the other by applying high voltage logic to outputs F_0 - F_6 .

Input variables which are not desired in a particular AND gate are fused out by selecting the AND gate with the decoder on output F_0 - F_5 , applying a HIGH or LOW TTL level to the

input variable to be fused, applying high voltage levels to all other input variables, and pulsing the chip enable input.

Product terms which are to be removed from an OR gate are fused out by selecting the appropriate AND gate using the decoder on inputs I_0 – I_5 , applying a high voltage level to the output for the desired OR gate, and pulsing the chip enable input.

The polarity of the output may be inverted by deselecting the chip, applying a high voltage to the desired output, and pulsing the FE input.



BASIC SYSTEM ARCHITECTURE

The Am2901 is designed to be used in microprogrammed systems. Figure 15 illustrates such an architecture. The nine instruction lines, the A and B addresses, and the D data inputs normally will all come from registers clocked at the same time as the Am2901. The register inputs come from a ROM or PROM – the "microprogram store". This memory contains sequences of microinstructions, typically 28 to 40 bits wide, which apply the proper control signals to the Am2901's and other circuits to execute the desired operation.

The address lines of the microprogram store are driven from the Am2909 microprogram sequencer. This device has facilities for storing an address, incrementing an address, jumping to any address, and linking subroutines. The Am2909 is controlled by some of the bits coming from the microprogram store. Essentially these bits are the "next instruction" control.

Note that with the microprogram register in-between the microprogram memory store and the Am2901's, an instruction accessed on one cycle is executed on the next cycle. As one instruction is executed, the next instruction is being read from microprogram memory. In this configuration, system speed is improved because the execution time in the Am2901's occurs in parallel with the access time of the microprogram store. Without the "pipeline register", these two functions must occur serially.

EXPANSION OF THE Am2901

The Am2901 is a four-bit CPU slice. Any number of Am2901's can be interconnected to form CPU's of 12, 16, 24, 36 or more bits, in four-bit increments. Figure 16 illustrates the interconnection of three Am2901's to form a 12-bit CPU, using ripple carry. Figure 17 illustrates a 16-bit CPU using carry lookahead, and Figure 18 is the general carry lookahead scheme for long words.

With the exception of the carry interconnection, all expansion schemes are the same. Refer to Figure 14. The Q_3 and RAM3 pins are bidirectional left/right shift lines at the MSB of the device. For all devices except the most significant, these lines are connected to the Q_0 and RAM9 pins of the adjacent more

significant device. These connections allow the Q-registers of all Am2901's to be shifted left or right as a contiguous n-bit register, and also allow the ALU output data to be shifted left or right as a contiguous n-bit word prior to storage in the RAM. At the LSB and MSB of the CPU, the shift pins should be connected to three-state multiplexers which can be controlled by the microcode to select the appropriate input signals to the shift inputs. (See Figure 19)

The open collector F=0 outputs of all the Am2901's are connected together and to a pull-up resistor. This line will go HIGH if and only if the output of the ALU contains all zeroes, Most systems will use this line as the Z (zero) bit of the processor status word.

The overflow and F₃ pins are generally used only at the most significant end of the array, and are meaningful only when two's complement signed arithmetic is used. The overflow pin is the Exclusive-OR of the carry-in and carry-out of the sign bit (MSB). It will go HIGH when the result of an arithmetic

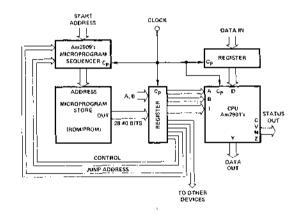


Figure 15. Microprogrammed Architecture Around Am 2901's.

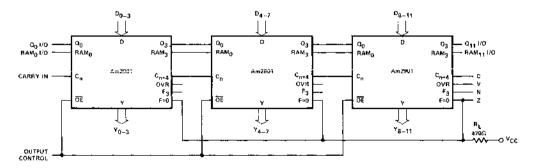


Figure 16. Three Am2901's used to Construct 12-Bit CPU with Ripple Carry. Corresponding A, B, and I Pins on all Devices are Connected Together.



operation is a number requiring more bits than are available, causing the sign bit to be erroneous. This is the overflow (V) bit of the processor status word. The F_3 pin is the MSB of the ALU output. It is the sign of the result in two's complement notation, and should be used as the Negative (N) bit of the processor status word.

The carry-out from the most significant Am2901 (C_{n+4} pin) is the carry-out from the array, and is used as the carry (C) bit of the processor status word.

Carry interconnections between devices may use either ripple carry or carry lookahead. For ripple carry, the carry-out (C_{n+4}) of each device is connected to the carry-in (C_n) of the next more significant device. Carry lookahead uses the Am2902 lookahead carry generator. The scheme is identical to that used with the 74181/74182. Users unfamiliar with this technique should refer to AMD's application note on Arithmetic Logic Units. Figures 17 and 18 illustrate single and multiple level lookahead.

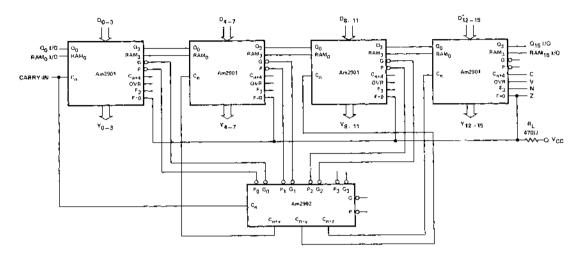


Figure 17, Four Am2901's in a 16-Bit CPU using the Am2902 for Carry Lookahead.

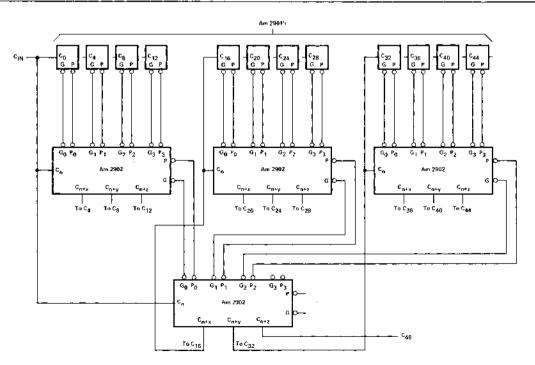


Figure 18. Carry Lookahead Scheme for 48-Bit CPU using 12 Am2901's. The Carry-Out Flag (C48) Should be Taken from the Lower Am2902 Rather than the Right-Most Am2901 for Higher Speed.



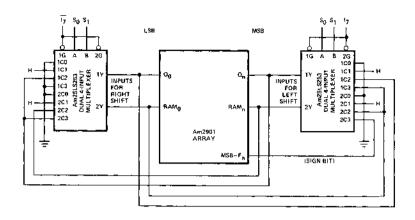


Figure 19. Three-State Multiplexers Used on Shift I/O Lines.

SHIFT I/O LINES AT THE END OF THE ARRAY

The Q-register and RAM up/down shift data transfers occur between devices over bidirectional lines. At the ends of the array, three-state multiplexers are used to select what the new inputs to the registers should be during shifting. Figure 19 shows two Am25LS253 dual four-input multiplexers connected to provide four shift modes. Instruction bit I_7 (from the Am2901) is used to select whether the left-shift multiplexer or the right-shift multiplexer is active. The four shift modes in this example are:

Zero A LOW is shifted into the MSB of the RAM on a down shift. If the Q-register is also shifted, then a LOW is deposited in the Q-register MSB. If the RAM or both registers are shifted up, LOWs are placed in the LSBs.

One Same as zero, but a HIGH level is deposited in the LSB or MSB.

Rotate A single precision rotate. The RAM MSB shifts into the LSB on a right shift and the LSB shifts into the MSB on a left shift. The Q-register, if shifted, will rotate in the same manner.

Arithmetic A double-length Arithmetic Shift if Q is also shifted. On an up shift a zero is loaded into the Q-register LSB and the Q-register MSB is loaded into the RAM LSB. On a down shift, the RAM LSB is loaded into the Q-register MSB and the ALU output MSB (Fn, the sign bit) is loaded into the RAM MSB. (This same bit will also be in the next, less significant RAM bit.)

Code		Source of New Data				01:6	7		
17	S ₁	S ₀	Ω_0	Q _n	RAM ₀	RAMn	Shift	Туре	
Н	L	L	0	Q _{n-1}	o	F _{n-1}	Up	Zero	
Н	L	Н	1	Q _{n-1}	1	F _{n-1}	(Right)	One	
Н	Н	L	o _n	Q_{n-1}	Fn	F _{n-1}		Rotate	
Н	Н	Н	o	Q _{n-1}	Ωn	F _{n-1}		Arithmetic	
L	L	L	Ω_1	0	F ₁	0	Down	Zero	
L	L	Н	Ω_1	1	F ₁	1	(Left)	One	
L	Н	i.	Ω1	σ^0	F ₁	Fo		Rotate	
L	H	Н	Q ₁	Fo	F ₁	$RAM_n = RAM_{n-1} = F_n$		Arithmetic	

HARDWARE MULTIPLICATION

Figure 20 illustrates the interconnections for a hardware multiplication using the Am2901. The system shown uses two devices for 8 x 8 multiplication, but the expansion to more bits is simple — the significant connections are at the LSB and MSB only.

The basic technique used is the "add and shift" algorithm, One clock cycle is required for each bit of the multiplier, On each cycle, the LSB of the multiplier is examined; if it is a "1", then

the multiplicand is added to the partial product to generate a new partial product. The partial product is then shifted one place toward the LSB, and the multiplier is also shifted one place toward the LSB. The old LSB of the multiplier is discarded. The cycle is then repeated on the new LSB of the multiplier available at $Q_{\rm D}$.

The multiplier is in the Am2901 Q-register. The multiplicand is in one of the registers in the register stack, R_a. The product will be developed in another of the registers in the stack, R_b.



The A address inputs are used to address the multiplicand in R_a , and the B address inputs are used to address the partial product in R_b . On each cycle, R_a is conditionally added to R_b , depending on the LSB of Q as read from the Q_0 output and both Q and ALU output are shifted down one place. The instruction lines to the Am2901 on every cycle will be:

I₈₇₆ = 4 (shift register stack input and Q register left)

 $I_{643} = 0 \qquad (Add)$

I₂₁₀ = 1 or 3 (select A, B or 0, B as ALU sources)

Figure 20 shows the connections for multiplication. The circled numbers refer to the paragraphs below.

 The adjacent pins of the Q-register and RAM shifters are connected together so that the Q-registers of both (or all) Am2901's shift up or down as a unit. Similarly, the entire eight-bit (or more) ALU output can be shifted as a unit prior to storage in the register stack.

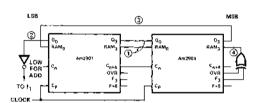


Figure 20. Interconnection for Dedicated Multiplication (8 by 8 Bit) (Corresponding A, B and I Connected Together).

- 2. The shift output at the LSB of the O-register determines whether the ALU source operands will be A and B (add multiplicand to partial product) or 0 and B (add nothing to partial product. Instruction bit I₁ can select between A, B or 0, B as the source operands; it can be driven directly from the complement of the LSB of the multiplier.
- 3. As the new partial product appears at the input to the register stack, it is shifted left by the RAM shifter. The new LSB of the partial product, which is complete and will not be affected by future operations, is available on the RAMo pin. This signal is returned to the MSB of the Q-register. On each cycle then, the just-completed LSB of the product is deposited in the MSB of the Q-register; the Q-register fills with the least significant half of the product.
- 4. As the ALU output is shifted down on each cycle, the sign bit of the new partial product should be inserted in the RAM MSB shift input. The F3 flag will be the correct sign of the partial product unless overflow has occurred. If overflow occurs during an addition or subtraction, the OVR flag will go HIGH and F3 is not the sign of the result. The sign of the result must then be the complement of F3. The correct sign bit to shift into the MSB of the partial product is therefore F3 * OVR; that is, F3 if overflow has not occurred and F3 if overflow has occurred. On the last cycle, when the MSB of the multiplier is examined, a conditional subtraction rather than addition should be performed, because the sign bit of the multiplier carries negative rather than positive arithmetic weight

$$(Y = -Y_i 2^i + Y_{i-1} 2^{i-1} + \dots + Y_0 2^0).$$

This scheme will produce a correct two's complement product for all multiplicands and multipliers in two's complement notation.

Figure 21 is a table showing the input states of the Am2901's for each step of a signed, two's complement multiplication.

Initial Register States R 0 Multiplier 1 Multiplicand 2 X 3 X				Am2901 Microcode Program 2's Comp. Multiply Date 8/5/75 By J. S							Final Register States R 0 Multiplier 1 Multiplicand 2 LSH Product 3 MSH Product				
				Pin States (Octal)									Jump		
s, F ——	Ď	Description	Repeat	А	В	1 ₈₇₆	1543	1210	Cn	Q ₀	O3	RAMo	RAM ₃	То	lf
OVA	α	Move Multiplier to Q	-	0	x	0	3	4	х	х	х	х	×		
ОЛВ	В	Clear R ₃	-	×	3.	2	4	3	х	х	×	×	×		-
(O+B)/2 (A+B)/2	В	Cond. Add & Shift	n-1	1	3	4	0	$\frac{1 \text{ or } 3}{i_1 = Q_0 LQ}$	0		RAMo	_	F ₃ ₩OVR		
(B = O)/2 (B = A)/2	В	Cond. Subt. & Shift	-	1	3	4	1	1 or 3 I ₁ = Q ₀ LO	1		RAM ₀		F ₃ ₩OVR		
0 / 0	В	Move LSH Prod. to R ₂	-	х	2	2	3	2	×	х	х	×	×		

X = Don't Care S = Source F = Function D = Destination



EXAMPLES OF SOME OTHER OPERATIONS

1. Byte Swapping

Occasionally the two halves of a 16-bit word must be swapped. D $_{0...7}$ is interchanged with D $_{8-15}$. The quickest way to perform this operation is to rotate the word in RAM, shifting two bits at a time. Only four shift cycles are required. The same register is selected on both the A and B ports; the two are added together with carry-in connected to carry-out, producing an up shift of one place; then the ALU is shifted up one more place prior to storage.

Byte Swap of Ro

A = B = 0 I = 701 RAM₀ = RAM₁₅ $C_{IN} = C_{OUT}$ Repeat 4 times

2. Instruction Fetch Cycle

Execution of a macroinstruction generally begins with an instruction fetch cycle. The current contents of the PC (in one of the registers) is the address of the macroinstruction to be fetched, and must be read out to the

memory address register. Then the PC is incremented to point to the next macroinstruction. The macroinstruction obtained from memory is then loaded into the Am2909 microprogram sequencer to cause a jump to the microcode for executing the instruction.

The PC can be read out and incremented in one cycle by using the Am2901 destination code 2, and addressing the PC with both the A and B addresses. The current value of PC will appear on the Y outputs, and PC+1 will be returned to the register. If the PC is in register 15, then:

$$A = B = 15$$
, $I = 203$, $Carry-in = 1$

The PC will be on the Y outputs via the RAM A-port. On the clock LOW-to-HIGH transition, the program counter is incremented and the value on the Y outputs is loaded into the memory address register. During the following cycle, the memory is read and, on the next clock LOW-to-HIGH transition the instruction from the memory is dropped into the Am2909 instruction register. The fetch operation requires only two microcycles.



INTRODUCTION:

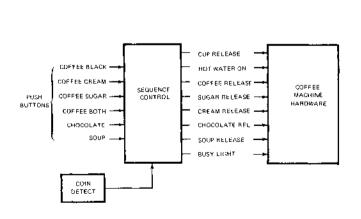
One of the main functions of digital logic systems is the generation of timing and control sequences. In the past, sequence control logic has been implemented by unique designs using flip-flops and gates. Current designs often use microprogram techniques based on combinations of counters and P/ROMs, due to the wide spread availability of low cost, high speed P/ROMs. These combinations of counters and P/ROMs tend to be significantly more flexible than their flip-flop logic, or hard wired, complements. Counter and P/ROM designs also tend to be significantly more powerful in the sense that they can be used to solve problems of much higher complexity and with a much lower chip count than is practical with a hard wired design.

The Am2909 Microprogram Sequencer is, in a sense, a sophisticated counter, which is used in conjunction with high speed P/ROMs in microprogram sequence control designs. Designs using the 2909 cover a broad range. Two of the better known examples include: Computer Control Units, which provide the sequence control for instruction execution in computers, and the I/O controllers for sophisticated computer Peripheral devices, such as high speed magnetic disc files and high speed magnetic tape units. Other examples include control units for high speed communication processing, and general purpose machine control at rates from DC to above 5 MHz. In order to understand the applicability of the Am2909 with sequence control, we will first examine the sequence control problem in general and the methods employed to solve it.

SEQUENCE CONTROL — SIMPLE PROBLEMS AND SOLUTIONS:

An example of a simple sequence control problem is the sequencer for an automatic coffee vending machine. In the automatic coffee vending machine, the user selects a machine operating sequence by pressing the appropriate button on the coffee machine's front panel. There may be six such buttons labeled: coffee black, coffee cream, coffee sugar, coffee cream & sugar, hot chocolate, soup. Each of these buttons must cause a specific sequence of actions inside the coffee machine in order to deliver a filled cup of coffee to the user. The sequence control unit for the coffee machine must generate signals such as: hot water on, coffee release, chocolate release, cream release, etc. In addition, it must generate these signals according to which of the six buttons was pressed. Furthermore, it must ignore the buttons until the coin has been inserted, during sequence operation, and until another coin has been inserted after the selected sequence is complete.

A block diagram of a hypothetical coffee machine and its sequence control are shown in Figure 1. The sequence list corresponding to each of the push buttons is shown in Figure 2. A coffee machine sequence controller can be designed using combinations of flip-flops and gates. An example of this type of design is shown in Figure 3. This approach is called the hard wired logic or random logic approach. The design in Figure 3 uses various gates, flip-flops and one-shots. This approach is typical of solid state sequence controller designs used in the past. The advantage



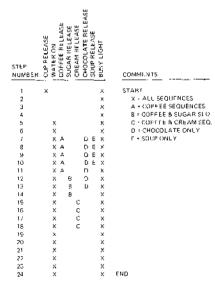


Figure 1. Coffee Machine Block Diagram.

Figure 2. Coffee Machine Combined Sequence List.



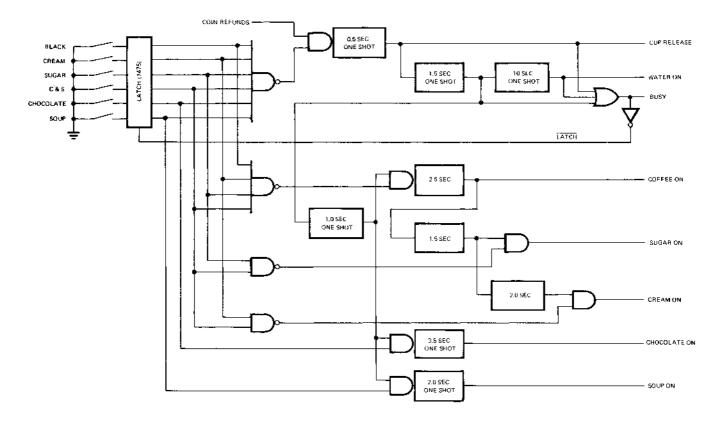


Figure 3. Coffee Machine Sequencer: Random Logic.



of this approach is that it results in a design with a reasonable part count and may closely parallel that of a previous design based on electro-mechanical components, such as relays and timers. There are several disadvantages with this design approach, however. The first disadvantage is that it is a unique design. The sequence controller for a soft drink machine, for instance, will be similar but different. A second, related disadvantage is that a redesign of the control logic may be required if the control sequence must be changed to accommodate different coffee machine hardware, such as using a different type of hot water valve. It involves, in any case, an evaluation of the effect of the sequence change upon the control logic to determine whether the design will have to be modified or redone. These two problem areas expand with the complexity of the sequence control problem. In the random logic design approach, the totality of the design is contained in the connection pattern of the integrated circuits, and the number of connections possible is approximately proportional to the square of the number of integrated circuits. This would indicate that the design complexity, probability of error, and probable number of design changes should also increase at about the same rate.

Microprogram techniques using counters and P/ROMs attack the disadvantages of the random logic approach as described above. The combination of a counter and a P/ROM can generate a single sequence, such as the black coffee sequence for a coffee machine, as shown in Figure 4. The contents of the P/ROM, which is also the sequence chart, is shown in Figure 5. The counter shown in Figure 4 is a binary counter and the P/ROM serves as a programmable decoder. The P/ROM decodes each count

value and activates the corresponding output lines as determined by the P/ROM contents or program. The counter is started at zero and steps to a maximum count of 31 in this case. The counter is incremented by pulses from a 2Hz Oscillator. This allows sequences of up to 16 seconds with a resolution of 1/2 second. This counter and P/ROM combination is directly analogous to an electro-mechanical cam-and-timer type sequence generator with a revolution time of 16 seconds and a programmable cam resolution of 1/32 of a circle or approximately 11 degrees. Notice, however, that the choice of a 5 bit binary counter allowing 32 possible counts, and the corresponding 32 word P/ROM, is arbitrary. It is equally practical to use an 8 bit counter and a 256 word P/ROM. This would allow generation of sequences up to 128 seconds long with a half second resolution. The sequencer shown in Figure 4 creates output signals according to the pattern shown in Figure 5. Figure 5 shows the contents of each of the 32 locations of the P/ROM. The 5 bit binary counter begins with a count value of zero and increments sequentially through a count value of 31. Each count value is decoded by the P/ROM and the corresponding output signals, as indicated in Figure 5 are generated.

The sequence generator shown in Figure 6 has the capability of eight individually selectable sequences. This is done by using a 256 word P/ROM and using the upper 3 bits of the P/ROM address to select one of eight possible thirty-two step sequences. If these 3 bits are 000, the 5 bit counter will step through the address sequence zero through 31. If these 3 bits are 001, the counter will step through 32 through 47, etc. We can apply this sequence

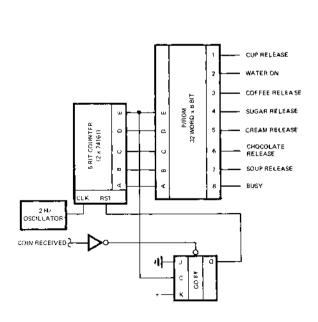


Figure 4. Simple Sequencer: Counter/ROM

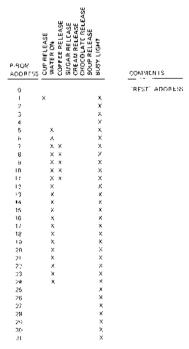


Figure 5. Counter/ROM Sequence Chart (Black Coffee).



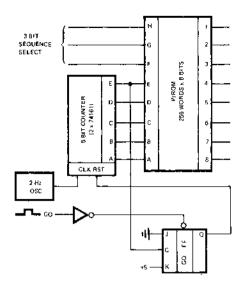


Figure 6. 8 Sequence Controller.

generator to our coffee machine problem by adding some circuitry as shown in Figure 7. Figure 7 shows the six push buttons, a register for holding the button status between the time it was pressed and the end of a sequence, and an encoder. The encoder converts the signal from each push button into a 3 bit binary number which is applied to the most significant 3 bits of the P/ROM. Push buttons 1 through 6 will generate binary codes 0 through 5, respectively. The 74148 encoder used performs this conversion as well as indicating whether or not any button is active. This activity output is used in conjunction with the output of a flip-flop which indicates a coin has been deposited and the button has been pushed. Once these two events have occurred, the sequence counter will step through the appropriate sequence locations in the P/ROM as shown in Figure 8. When the selected sequence is ended, the BUSY P/ROM output bits are used to clear the coin flip-flop, restoring the system to the initial conditions for the next cup..

We have now built a sequence controller, using 9 chips, which is capable of 8 selectable sequences of up to 32 steps each. Also, any and all sequences may be changed by modifying the contents of the P/ROM, a single chip, without any rewiring or other modifications to the hardware.

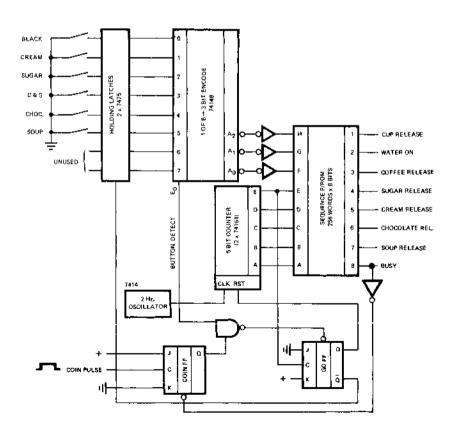


Figure 7. Coffee Machine ROM/Counter Sequencer.



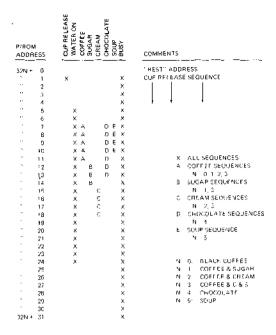


Figure 8. Combined ROM/Counter Sequence Chart.

Although in our coffee machine example the sequencer is stepped along at a rate of 2 hertz, the circuit is practical at rates to several megahertz. At rates above one megahertz, the decoding time of the P/ROM becomes significant. This can be avoided by using a modification of the original counter and P/ROM technique. The modified technique has various names, such as pipelining and look-ahead. This method puts a register at the P/RQM outputs to hold the output values, while the P/ROM itself is decoding the next set of outputs. Figure 9 shows a block diagram of such a system. In this system, a counter is created using an incrementer and a register. This incrementer and register combination is characteristic of the internal design of synchronous counter chips, such as the 74163. In Figure 9 these two functions have been separated to grant access to the path between the incrementer and the register. The incrementer output creates the next count value and applies it to the count register input. The incrementer

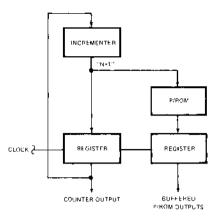


Figure 9. Pipelined Counter /ROM System.

output is also applied to the address input of the sequence P/ROMs. The sequence P/ROMs decode this next count value and apply their decoded outputs to another register. When a clock pulse is applied, the counter will increment since it will be loaded with the current count value plus one. Likewise, the P/ROM buffer register will be loaded immediately with the decoded P/ROM data for that state. The advantage of this approach is that the decoded P/ROM data is available almost immediately after the clock pulse. as opposed to 60 to 80 nanoseconds later, as determined by propagation delay to the P/ROMs. This is shown in the timing diagram of Figure 10. An additional feature of the pipelined approach is that the outputs remain stable from one clock period to the next. In the simpler counter and P/ROM system, the P/ROM outputs are undefined for a period of typically 60 to 80 nanoseconds after the clock pulse. Not only are they undefined during this interval, but they may change rapidly and create positive or negative pulses. This is due to changes occurring in the internal decode logic of the P/ROM.

The Am2909 is designed to be used in pipelined systems. The 2909 contains an incrementer and logic for generating the next P/ROM state. Figure 11 shows a 2909 based pipelined sequencer. Some of the features of the 2909, which will be discussed below, are not shown for the sake of clarity. Note that the sequencer in Figure 11 differs

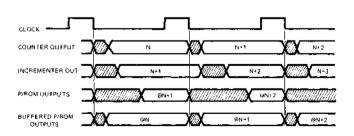


Figure 10. Pipelined Counter /ROM Timing Diagram.



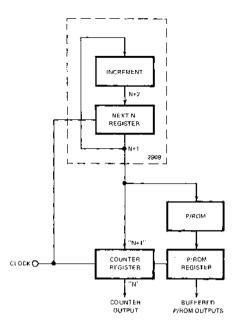


Figure 11. 2909 Pipelined System.

slightly from the sequencer shown in Figure 9. In the 2909, a register is placed after rather than before the incrementer. This register holds the value of the next P/ROM address, the decoded contents of which will be loaded into the P/ROM buffer register. The incrementer takes the 2909 outputs and increments them to form the next address plus one, to be loaded into the 2909 microprogram counter register. This arrangement provides high speed since the next address is available from this next address microprogram counter register almost immediately after the rise of the clock pulse. The inherently slower generation of the incremented value thus occurs between clock pulses. This is, in a sense, further pipelining of the system. Because of this arrangement, microprogram step times as low as 100 nanoseconds are possible.

COMPLEX PROBLEMS AND SOLUTIONS

The simple sequencers described above are useful and practical for small and medium sized sequence control systems. As the system's size expands, in terms of number of selectable sequences, and/or number of steps per sequence, the P/ROM memory required expands proportionally. As the P/ROM memory size, and therefore cost, starts becoming significant in our design, we naturally look for methods of reducing the size of memory required. There is one fundamental technique with several implementation methods for reducing memory size, and this is to identify and eliminate redundant sequences.

There are two general methods of identifying and eliminating redundant sequences. One method is repetition, where a single step or group of steps is repeated a number of times. In our sequence list of Figure 8, we can see several areas in each sequence where the outputs do not

change from one P/ROM address to the next. These redundant states can be reduced to a single state with a repeat counter to cause those outputs to remain unchanged until they have been repeated the appropriate number of clock steps. This same technique can be applied in cases where a number of steps are repeated several times. The second general method is to identify common subsequences where a group of identical steps are used in several different sequences or even several times within the same sequence. A sub-sequence, called a subroutine, typically performs some commonly used function. An example is the cup release sequence shown in the Chart in Figure 8.

A wide variety of ways exist for implementing the repeat function. A simple method of implementing the repeat function for single P/ROM states or micro-instructions is shown in Figure 12. In this figure we have added a synchronous counter, a 74163 in this case, which is loaded from the pipeline P/ROMs at the same time the P/ROM. buffer, or micro-instruction register, is loaded. This counter has an output which detects when the 4 bits of the counter. are all ones. This output is used to gate the clock pulse to the sequencer control system, including the 2909 microprogram counter register and the micro-instruction register. This signal is also used to enable the load enable input to the counter itself. If the counter is loaded with all ones for every clock pulse, the sequencer will operate as before, with the 2909 micro-program counter register and the micro-instruction register being loaded on every clock pulse. If the counter is loaded with other than all ones, for instance with all zeros, the clock for the rest of the system will be disabled until the counter reaches all ones, in this case after 15 counts. This means that the micro-instruction register outputs will remain unchanged for a total of 16 counts, 15 counts corresponding to the clock being disabled plus one count during which the clock was enabled, in order to execute a micro-instruction in N times, one simply loads the repeat counter with a value of 16 minus N.

It may be desirable to repeat a group of several instructions. In this case, we will require somewhat more complicated logic. In order to perform this repeat, we must go back to the first instruction in the repeated group. The logic to perform this is shown in Figure 13. A sequence chart example is shown in Figure 14. In order to repeat a group of instructions, we must be able to conditionally go from the last step of the group sequence to the first step of the group sequence. The hardware in Figure 13 does this. An added P/ROM, buffer register, and a multiplexer in the 2909 provides a method of selecting a next address different from the current address plus one. This next address will typically be the starting address of our repetitive sequence, the repeat address. Our repeat counter is now used to control this multiplexer. When we reach the last step in the repetitive sequence, the repeat counter all ones detect output is used to control 2909 multiplexer selection. If the repeat counter detect output is zero, then the multiplexer selects the repeat address from the P/ROM buffer register. If the repeat counter detect output is one, then the current address plus one is selected, which is in the 2909 microprogram counter register.

Note that the repeat counter is incremented only when it is used to control the multiplexer. In this way it will be incremented once per repetition of the group, or once per loop, rather than once per clock pulse. This circuit can be



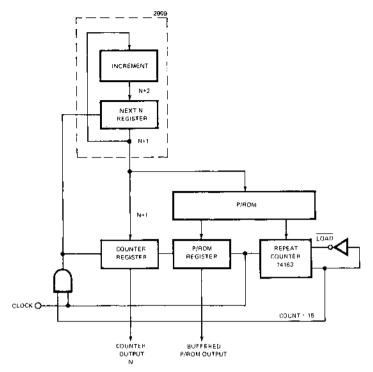


Figure 12. Single Instruction Repeat.

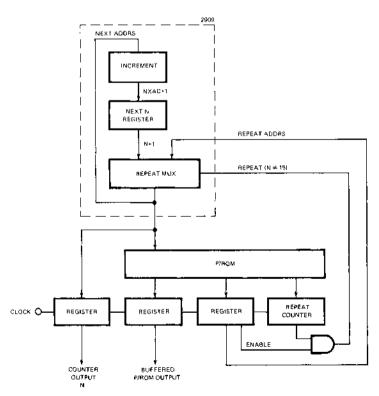


Figure 13. Group Repeat Logic.



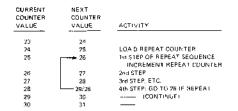


Figure 14. Repeat Sequence Example.

used to implement a single micro-instruction repeat function by simply causing the address P/ROM output to be equal to the address of the current micro-instruction. Note also that we no longer use the repeat counter every cycle. It can therefore be loaded from another source within the overall system at an arbitrary time. In a computer control unit context, our repeat counter might be loaded with a shift count from the instruction register at some time before the shifting operation was actually performed.

The second method of eliminating redundant coding in the P/ROM memory is by using sub-sequences or subroutines. In this case, a sub-sequence exists which is used in a variety of different sequences in the P/ROM memory. For example, the cup release sub-sequence in our coffee vending machine example is used in all sequences. Instead of having six different sets of codes doing the same function, one set of code is used plus a method to get to the code and return to the original sequence when the sub-sequence is complete. With the addition of the address generation P/ROM shown in Figure 13, we now have a method of going from a current location to some entirely different location within the memory. To implement sub-sequences, we need a method of returning to the next location in our original sequence. To be specific, we wish the next P/ROM address to be the first address in our commonly used sub-sequence. or subroutine, and at the end of that subroutine we wish to return to our current address plus one. In order to do this we must be able to (1) Select a desired subroutine address as the next P/ROM value, and (2) Save the current P/ROM address value plus one as the next P/ROM value to be used upon completion of our subroutine. This can be done by using the logic of Figure 13, which allows us to select an arbitrary next P/ROM address value, plus the addition of a register to save our normal next address value, which is the current P/ROM address value plus one. The design of this subroutine logic shown in Figure 15 and its operation is shown in the sequence chart of Figure 16.

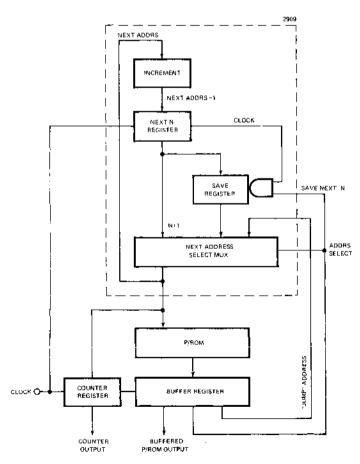


Figure 15. Next Address Save Logic.



CURRENT COUNTER VALUE	NEXT COUNTER VALUE	SAVE REGISTER CONTENTS	ACTIVITY
23	24	_	
24	37		GO YO 37 AND SAVE RETURN
37	38	25	1st STEP OF SUBADUTINE
38	39	25	2nd STEP OF SUBROUTINE
39	40	25	3rd STEP OF SUBROUTINE
40	25	25	4th STEP OF SUBROUTINE AND
			RETURN (SELECT SAVE REG
			AS NEXT COUNTER VALUE;
25	26	25	CONTINUE
26	27	25	
27	28	25	

Figure 16. Subroutine Sequence Example.

The 2909 contains a save register. In fact, it is implemented as a Last In, First Out stack of 4 registers. This may be considered as a set of bi-directional shift registers which allow data to be shifted in or out according to the 2909 controls. These effective shift registers are 4 bits long. Having more than one save register allows one subroutine to

use, or call, another subroutine. This process of calling a subroutine which in turn calls another subroutine, can be repeated up to 4 times in the 2909. At the completion of each subroutine, the last return value which was put into the save shift registers is used to set the next P/ROM address value. This Last-In-First-Out function is called a stack because it is conceptionally similar to a stack of plates. The current output of the 2909 LIFO stack corresponds to the last value shifted into the stack in a similar manner to a stack of plates, where the top-most plate on the stack corresponds to the last one placed there.

There exists another method of sub-sequence selection. In this case, the sequence selection register is decoded more than once. This is shown in Figure 17. Here we are assuming that the original command selection is stored in a register, and that this register has a P/ROM at its output which decodes its contents to produce a P/ROM address value corresponding to the selected sequence. In the above discussions, we have assumed that this command register (the push button register) is decoded once and the selected

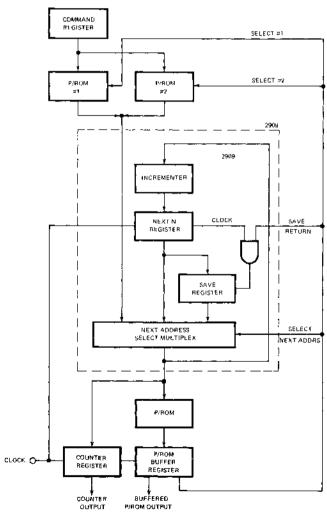


Figure 17. Multiple Decode Subsequence Select.



sequence runs through to completion. Another possibility is to use the same command register to select several subsequences rather than one main sequence. This means that our command register is decoded several times. The first time it is decoded, it selects the first sub-sequence. At the end of this sub-sequence, a second decode P/ROM is enabled which decodes the command register a second time. This selects a second sub-sequence, etc. In this manner, as many sub-sequences as desired may be selected by sequentially decoding this command register. This method has the advantage of not requiring a special set of P/ROM outputs for address generation, in some cases, it may save a micro-instruction step which would be required to generate a subroutine call.

A general purpose sequence control system using 2909's is shown in Figure 18. This is a semi-pipelined system. This means that part of the system is pipelined for high-speed functions which must have signals available at the beginning of the cycle and part of the system is not pipelined, meaning that the next P/ROM address is simply

loaded into a register which is decoded by P/ROMs, the outputs of which are not needed until the middle or the end of the micro-instruction cycle. The timing diagram for this system is shown in Figure 19. Note that in this 2909 system we have included a repeat counter and a condition testing multiplexer which can select the repeat counter output or any of several other external signals to control the source of the next P/ROM address. These external signals are synchronized with the buffer register to insure that their outputs will remain stable during the micro-cycle. This allows the look-ahead P/ROMs sufficient time for their outputs to stabilize and drive the inputs to the micro-instruction register.

Now that our general purpose sequence controller has been designed, all that remains is to determine the number of high-speed, pipelined P/ROM outputs and lower-speed, non-pipelined P/ROM outputs required by our particular application. Also, we must determine the number and source of our condition control inputs. Last of all, we must create the pattern to be placed into our control

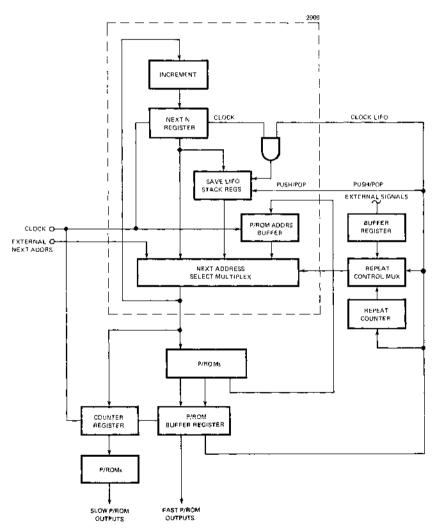


Figure 18. General Purpose 2909 Sequencer.



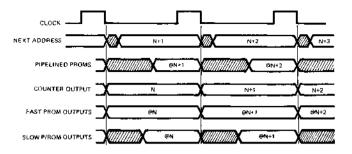


Figure 19. 2909 Sequencer Timing Diagram.

P/ROMs. All three of these tasks are part of the system design for a particular sequence control system. The last, however, can be supported by the RAYASM micro-assembler which is designed to simplify the generation of P/ROM patterns and the corresponding P/ROM programming tapes.

These two methods of sub-sequence generation, subroutine call and multiple command decode, may both be used within a given system. The method of subroutine call is the most general and can be used to create very long effective sequences. It is useful in medium and large systems. The method of multiple command decode has the advantage of eliminating the subroutine call micro-instructions and their corresponding execution time. It is useful in systems which

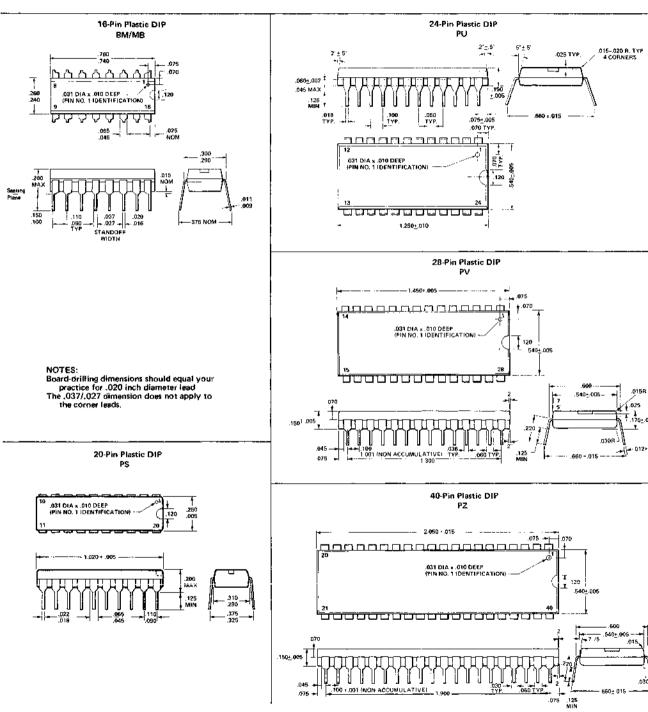
have a large number of short sequences where execution time is important. The subsequences are typically few and short in computer control units since this provides not only fast instruction execution but also minimizes interrupt and DMA latency times. In many computer control units, interrupt and DMA activities are allowed to occur only between instructions. This means that the waiting period, or latency time, between an interrupt or DMA request in the beginning of the micro-program sequence which services that request can be long as the execution time of the longest instruction in the instruction set. For this reason, the multiple decode method is often used in CCU's, since it is desirable and sometimes necessary in these types of systems to keep the instruction execution sequences for all instructions below some maximum figure.



PACKAGING INFORMATION

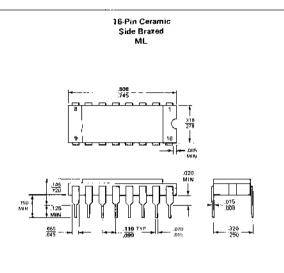
These drawings are preliminary and subject to change when the final product is complete.

DUAL IN-LINE PLASTIC





DUAL IN-LINE CERAMIC SIDE BRAZED

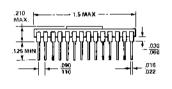


> 28-Pin Ceramic Side Brazed MV

20-Pin Ceramic

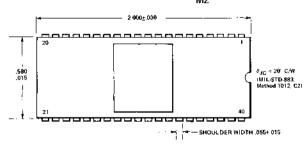
NOTES: All dimensions in inches Leads are intended for insertion in hole rows on .300" centers Board-drilling dimensions should equal your practice for .020 inch diameter lead Leads are gold-plated kovar

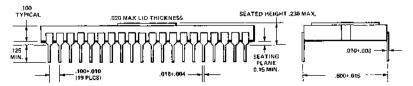






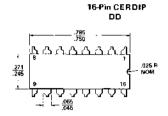
40-Pin Ceramic Side Brazed MZ

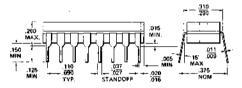






DUAL IN-LINE CERAMIC CERDIP

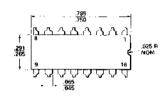


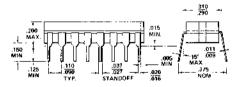


NOTES:
All dimensions in inches
Leads are intended for insertion in hole
rows on .300" centers
They are purposely shipped with "positive"
misalignment to facilitate insertion
Board-drilfing dimensions should equal your
practice for .020" diameter lead
Leads are tin-plated alloy 42
"The .037/.027 dimension does not apply to
the corner leads

40-Pin CERDIP D2 2.080 2.100 2.100 3.518 40 0.258 NOM 165 165 100 17 100 TVP 100 TV

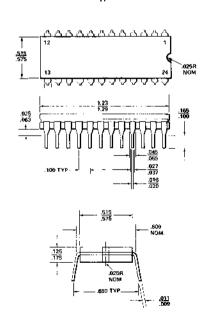
16-Pin CERDIP DM





NOTES:
All dimensions in inches
Leads are intended for insertion in hole
rows on .300" centers.
They are purposely shipped with "positive"
misalignment to facilitate insertion
Board-drilling dimensions should equal your
practice for .020" diameter lead
Leads are tin-plated alloy 42
*The .037/.027 dimension does not apply to
the corner leads

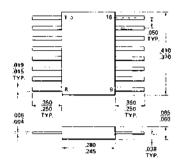
24-Pin CEROIP





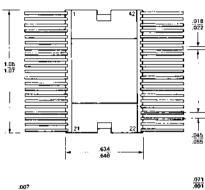
FLAT PACKAGES

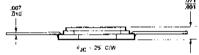




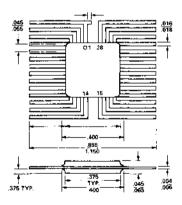
NOTES: All dimensions in inches Leads are tin-plated alloy 42

42-Pin Metal F2





28-Pin Metal FV









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