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CP/M 2.0 ALTERATION GUIDE

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1. INTRODUCTION

The standard CP/M system assumes operation on an Intel MDS-800 microcomputer development system, but is designed so that the user can alter a specific set of subroutines which define the hardware operating environment. In this way, the user can produce a diskette which operates with any IBM-3741 format compatible drive controller and other peripheral devices.

Although standard CP/M 2.0 is configured for single density floppy disks, field-alteration features allow adaptation to a wide variety of disk subsystems from single drive minidisks through high-capacity "hard disk" systems. In order to simplify the following adaptation process, we assume that CP/M 2.0 will first be configured for single density floppy disks where minimal editing and debugging tools are available. If an earlier version of CP/M is available, the customizing process is eased considerably. In this latter case, you may wish to briefly review the system generation process, and skip to later sections which discuss system alteration for non-standard disk systems.

In order to achieve device independence, CP/M is separated into three distinct modules:

BIOS - basic I/O system which is environment dependent BDOS - basic disk operating system which is not dependent upon the hardware configuration CCP - the console command processor which uses the BDOS

Of these modules, only the BIOS is dependent upon the particular hardware. That is, the user can "patch" the distribution version of CP/M to provide a new BIOS which provides a customized interface between the remaining CP/M modules and the user's own hardware system. The purpose of this document is to provide a step-by-step procedure for patching your new BIOS into CP/M.

If CP/M is being tailored to your computer system for the first time, the new BIOS requires some relatively simple software development and testing. The standard BIOS is listed in Appendix B, and can be used as a model for the customized package. A skeletal version of the BIOS is given in Appendix C which can serve as the basis for a modified BIOS. In addition to the BIOS, the user must write a simple memory loader, called GETSYS, which brings the operating system into memory. In order to patch the new BIOS into CP/M, the user must write the reverse of GETSYS, called PUTSYS, which places an altered version of CP/M back onto the diskette. PUTSYS can be derived from GETSYS by changing the disk read commands into disk write commands. Sample skeletal GETSYS and PUTSYS programs are described in Section 3, and listed in Appendix D. In order to make the CP/M system work automatically, the user must also supply a cold start loader, similar to the one provided with CP/M (listed in Appendices A and B). A skeletal form of a cold start loader is given in Appendix E which can serve as a model for your loader.

2. FIRST LEVEL SYSTEM REGENERATION

The procedure to follow to patch the CP/M system is given below in several steps. Address references in each step are shown with a following "H" which denotes the hexadecimal radix, and are given for a 20K CP/M system. For larger CP/M systems, add a "bias" to each address which is shown with a "+b" following it, where b is equal to the memory size - 20K. Values for b in various standard memory sizes are

24K:b = 24K - 20K = 4K = 1000H32K:b = 32K - 20K = 12K = 3000H40K:b = 40K - 20K = 20K = 5000H48K:b = 48K - 20K = 28K = 7000H56K:b = 56K - 20K = 36K = 9000H62K:b = 62K - 20K = 42K = A800H64K:b = 64K - 20K = 44K = B000H

Note: The standard distribution version of CP/M is set for operation within a 20K memory system. Therefore, you must first bring up the 20K CP/M system, and then configure it for your actual memory size (see Second Level System Generation).

(1) Review Section 4 and write a GETSYS program which reads the first two tracks of a diskette into memory. The data from the diskette must begin at location 3380H. Code GETSYS so that it starts at location 100H (base of the TPA), as shown in the first part of Appendix d.

(2) Test the GETSYS program by reading a blank diskette into memory, and check to see that the data has been read properly, and that the diskette has not been altered in any way by the GETSYS program.

(3) Run the GETSYS program using an initialized CP/M diskette to see if GETSYS loads CP/M starting at 3380H (the operating system actually starts 128 bytes later at 3400H).

(4) Review Section 4 and write the PUTSYS program which writes memory starting at 3380H back onto the first two tracks of the diskette. The PUTSYS program should be located at 200H, as shown in the second part of Appendix D.

(5) Test the PUTSYS program using a blank uninitialized diskette by writing a portion of memory to the first two tracks; clear memory and read it back using GETSYS. Test PUTSYS completely, since this program will be used to alter CP/M on disk.

(6) Study Sections 5, 6, and 7, along with the distribution version of the BIOS given in Appendix B, and write a simple version which performs a similar function for the customized environment. Use the program given in Appendix C as a model. Call this new BIOS by the name CBIOS (customized BIOS). Implement only the primitive disk operations on a single drive, and simple console input/output functions in this phase.

(7) Test CBIOS completely to ensure that it properly performs console character I/O and disk reads and writes. Be especially careful to ensure that no disk write operations occur accidently during read operations, and check that the proper track and sectors are addressed on all reads and writes. Failure to make these checks may cause destruction of the initialized CP/M system after it is patched.

(8) Referring to Figure 1 in Section 5, note that the BIOS is placed between locations 4A00H and 4FFFH. Read the CP/M system using GETSYS and replace the BIOS segment by the new CBIOS developed in step (6) and tested in step (7). This replacement is done in the memory of the machine, and will be placed on the diskette in the next step.

(9) Use PUTSYS to place the patched memory image of CP/M onto the first two tracks of a blank diskette for testing.

(10) Use GETSYS to bring the copied memory image from the test diskette back into memory at 3380H, and check to ensure that it has loaded back properly (clear memory, if possible, before the load). Upon successful load, branch to the cold start code at location 4A00H. The cold start routine will initialize page zero, then jump to the CCP at location 3400H which will call the BDOS, which will call the CBIOS. The CRIOS will be asked by the CCP to read sixteen sectors on track 2, and if successful, CP/M will type "A>", the system prompt.

When you make it this far, you are almost on the air. If you have trouble, use whatever debug facilities you have available to trace and breakpoint your CBIOS.

(11) Upon completion of step (10), CP/M has prompted the console for a command input. Test the disk write operation by typing

SAVE 1 X.COM

(recall that all commands must be followed by a carriage return).

CP/M should respond with another prompt (after several disk accesses):

A>

If it does not, debug your disk write functions and retry.

(12) Then test the directory command by typing

DIR

CP/M should respond with

A: X COM

(13) Test the erase command by typing

ERA X.COM

CP/M should respond with the A prompt. When you make it this far, you should have an operational system which will only require a bootstrap loader to function completely.

(14) Write a bootstrap loader which is similar to GETSYS, and place it on track \emptyset , sector l using PUTSYS (again using the test diskette, not the distribution diskette). See Sections 5 and 8 for more information on the bootstrap operation.

(15) Retest the new test diskette with the bootstrap loader installed by executing steps (11), (12), and (13). Upon completion of these tests, type a control-C (control and C keys simultaneously). The system should then execute a "warm start" which reboots the system, and types the A prompt.

(16) At this point, you probably have a good version of your customized CP/M system on your test diskette. Use GETSYS to load CP/M from your test diskette. Remove the test diskette, place the distribution diskette (or a legal copy) into the drive, and use PUTSYS to replace the distribution version by your customized version. Do not make this replacement if you are unsure of your patch since this step destroys the system which was sent to you from Digital Research.

(17) Load your modified CP/M system and test it by typing

DIR

CP/M should respond with a list of files which are provided on the initialized diskette. One such file should be the memory image for the debugger, called DDT.COM.

NOTE: from now on, it is important that you always reboot the CP/M system (ctl-C is sufficient) when the diskette is removed and replaced by another diskette, unless the new diskette is to be read only.

(18) Load and test the debugger by typing

DDT

(see the document "CP/M Dynamic Debugging Tool (DDT)" for operating procedures. You should take the time to become familiar with DDT, it will be your best friend in later steps.

(19) Before making further CBIOS modifications, practice using the editor (see the ED user's guide), and assembler (see the ASM user's guide). Then recode and test the GETSYS, PUTSYS, and CBIOS programs using ED, ASM, and DDT. Code and test a COPY program which does a sector-to-sector copy from one diskette to another to obtain back-up copies of the original diskette (NOTE: read your CP/M Licensing Agreement; it specifies your legal responsibilities when copying the CP/M system). Place the copyright notice

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on each copy which is made with your COPY program.

(20) Modify your CBIOS to include the extra functions for punches, readers, signon messages, and so-forth, and add the facilities for a additional disk drives, if desired. You can make these changes with the GETSYS and PUTSYS programs which you have developed, or you can refer to the following section, which outlines CP/M facilities which will aid you in the regeneration process.

You now have a good copy of the customized CP/M system. Note that although the CBIOS portion of CP/M which you have developed belongs to you, the modified version of CP/M which you have created can be copied for your use only (again, read your Licensing Agreement), and cannot be legally copied for anyone else's use.

It should be noted that your system remains file-compatible with all other CP/M systems, (assuming media compatiblity, of course) which allows transfer of non-proprietary software between users of CP/M.

3. SECOND LEVEL SYSTEM GENERATION

Now that you have the CP/M system running, you will want to configure CP/M for your memory size. In general, you will first get a memory image of CP/M with the "MOVCPM" program (system relocator) and place this memory image into a named disk file. The disk file can then be loaded, examined, patched, and replaced using the debugger, and system generation program. For further details on the operation of these programs, see the "Guide to CP/M Features and Facilities" manual.

Your CBIOS and BOOT can be modified using ED, and assembled using ASM, producing files called CBIOS.HEX and BOOT.HEX, which contain the machine code for CBIOS and BOOT in Intel hex format.

To get the memory image of CP/M into the TPA configured for the desired memory size, give the command:

MOVCPM xx *

where "xx" is the memory size in decimal K bytes (e.g., 32 for 32K). The response will be:

> CONSTRUCTING XXK CP/M VERS 2.0 READY FOR "SYSGEN" OR "SAVE 34 CPMXX.COM"

At this point, an image of a CP/M in the TPA configured for the requested memory size. The memory image is at location 0900H through 227FH. (i.e., The BOOT is at 0900H, the CCP is at 980H, the BDOS starts at 1180H, and the BIOS is at 1F80H.) Note that the memory image has the standard MDS-800 BIOS and BOOT on it. It is now necessary to save the memory image in a file so that you can patch your CBIOS and CBOOT into it:

SAVE 34 CPMxx.COM

The memory image created by the "MOVCPM" program is offset by a negative bias so that it loads into the free area of the TPA, and thus does not interfere with the operation of CP/M in higher memory. This memory image can be subsequently loaded under DDT and examined or changed in preparation for a new generation of the system. DDT is loaded with the memory image by typing:

DDT CPMxx.COM	Load DDT, then read the CPM
	image

DDT should respond with

NEXT	PC
2300	0100
-	

(The DDT prompt)

You can then use the display and disassembly commands to examine

portions of the memory image between 900H and 227FH. Note, however, that to find any particular address within the memory image, you must apply the negative bias to the CP/M address to find the actual address. Track 00, sector 01 is loaded to location 900H (you should find the cold start loader at 900H to 97FH), track 00, sec. r 02 is loaded into 980H (this is the base of the CCP), and so-forth through the entire CP/M system load. In a 20K system, for example, the CCP resides at the CP/M address 3400H, but is placed into memory at 980H by the SYSGEN program. Thus, the negative bias, denoted by n, satisfies

3400H + n = 980H, or n = 980H - 3400H

Assuming two's complement arithmetic, n = D580H, which can be checked by

3400H + D580H = 10980H = 0980H (ignoring high-order overflow).

Note that for larger systems, n satisfies

(3400H+b) + n = 980H, or n = 980H - (3400H + b), or n = D580H - b.

The value of n for common CP/M systems is given below

memc~y size	bias b	negative offset n
20K	NOOOH	D580H - 0000H = D580H
24K	1000H	D580H - 1000H = C580H
32K	зоорн	D580H - 3000H = A580H
4 Ø K	5000H	D580H - 5000H = 8580H
4 8K	7000H	D580H - 7000H = 6580H
5 6K	9000H	D580H - 9000H = 4580H
62K	ABØØH	D580H - A800H = 2D80H
64K	BØØØH	D580H - B000H = 2580H

Assume, for example, that you want to locate the address x within the memory image loaded under DDT in a 20K system. First type

Hx,n Hexadecimal sum and difference

and DDT will respond with the value of x+n (sum) and x-n (difference). The first number printed by DDT will be the actual memory address in the image where the data or code will be found. The input

H3400,D580

for example, will produce 980H as the sum, which is where the CCP is located in the memory image under DDT.

Use the L command to disassemble portions the BIOS located at (4A00H+b)-n which, when you use the H command, produces an actual address of 1F80H. The disassembly command would thus be

It is now necessary to patch in your CBOOT and CBIOS routines. The BOOT resides at location Ø900H in the memory image. If the actual load address is "n", then to calculate the bias (m) use the command:

H900,n Subtract load address from target address.

The second number typed in response to the command is the desired bias (m). For example, if your BOOT executes at 0080H, the command:

H900,80

will reply

0980 0880

Sum and difference in hex.

Therefore, the bias "m" would be Ø880H. To read-in the BOOT, give the command:

ICBOOT.HEX Input file CBOOT.HEX

Then:

Rm

Read CBOOT with a bias of m (=900H-n)

You may now examine your CBOOT with:

L900

We are now ready to replace the CBIOS. Examine the area at 1F80H where the original version of the CBIOS resides. Then type

ICBIOS.HEX Ready the "hex" file for loading

assume that your CBIOS is being integrated into a 20K CP/M system, and thus is origined at location 4A00H. In order to properly locate the CBIOS in the memory image under DDT, we must apply the negative bias n for a 20K system when loading the hex file. This is accomplished by typing

RD580 Read the file with bias D580H

Upon completion of the read, re-examine the area where the CBIOS has been loaded (use an "LIF80" command), to ensure that is was loaded properly. When you are satisfied that the change has been made, return from DDT using a control-C or "G0" command.

Now use SYSGEN to replace the patched memory image back onto a diskette (use a test diskette until you are sure of your patch), as shown in the following interaction

SYSGEN Start the SYSGEN program SYSGEN VERSION 2.0 Sign-on message from SYSGEN
SOURCE DRIVE NAME (OR RETURN TO SKIP)
Respond with a carriage return
to skip the CP/M read operation
since the system is already in
memory.
DESTINATION DRIVE NAME (OR RETURN TO REBOOT)
Respond with "B" to write the
new system to the diskette in
drive B.
DESTINATION ON B, THEN TYPE RETURN
Place a scratch diskette in
drive B, then type return.
FUNCTION COMPLETE
DESTINATION DRIVE NAME (OR RETURN TO REBOOT)

Place the scratch diskette in your drive A, and then perform a coldstart to bring up the new CP/M system you have configured.

Test the new CP/M system, and place the Digital Research copyright notice on the diskette, as specified in your Licensing Agreement:

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The following program provides a framework for the GETSYS and PUTSYS programs referenced in Section 2. The READSEC and WRITESEC subroutines must be inserted by the user to read and write the specific sectors.

> GETSYS PROGRAM - READ TRACKS Ø AND 1 TO MEMORY AT 3380H ; REGISTER USE ; (SCRATCH REGISTER) A ; B TRACK COUNT (0, 1) ; C SECTOR COUNT (1,2,...,26) ; (SCRATCH REGISTER PAIR) DE : LOAD ADDRESS HL. ; SET TO STACK ADDRESS SP ; SP,3380H ;SET STACK POINTER TO SCRATCH AREA H, 3380H ;SET BASE LOAD ADDRESS START: LXI SP,3380H LXI MVI B.Ø START WITH TRACK Ø RDTRK : ; READ NEXT TRACK (INITIALLY Ø) C,1 MVI ; READ STARTING WITH SECTOR 1 RDSEC: ;READ NEXT SECTOR CALL READSEC ;USER-SUPPLIED SUBROUTINE ; MOVE LOAD ADDRESS TO NEXT 1/2 PAGE LXI D,128 ; HL = HL + 128DAD D ;SECTOR = SECTOR + 1 INR C MOV A,C ;CHECK FOR END OF TRACK 27 CPI JC RDSEC ;CARRY GENERATED IF SECTOR < 27 ; ARRIVE HERE AT END OF TRACK, MOVE TO NEXT TRACK ; INR B NOV ;TEST FOR LAST TRACK A,B CPI 2 JC RDTRK ;CARRY GENERATED IF TRACK < 2 ; ARRIVE HERE AT END OF LOAD, HALT FOR NOW ; HLT ; USER-SUPPLIED SUBROUTINE TO READ THE DISK ; READSEC: ENTER WITH TRACK NUMBER IN REGISTER B, ; SECTOR NUMBER IN REGISTER C. AND ; ADDRESS TO FILL IN HL ; ; PUSH B ;SAVE B AND C REGISTERS PUSH H ;SAVE HL REGISTERS perform disk read at this point, branch to label START if an error occurs POP H ; RECOVER HL POP B ;RECOVER B AND C REGISTERS RET ; BACK TO MAIN PROGRAM END START

Note that this program is assembled and listed in Appendix C for reference purposes, with an assumed origin of 100H. The hexadecimal operation codes which are listed on the left may be useful if the program has to be entered through your machine's front panel switches.

The PUTSYS program can be constructed from GETSYS by changing only a few operations in the GETSYS program given above, as shown in Appendix D. The register pair HL become the dump address (next address to write), and operations upon these registers do not change within the program. The READSEC subroutine is replaced by a WRITESEC subroutine which performs the opposite function: data from address HL is written to the track given by register B and sector given by register C. It is often useful to combine GETSYS and PUTSYS into a single program during the test and development phase, as shown in the Appendix.

5. DISKETTE ORGANIZATION

The sector allocation for the standard distribution version of CP/M is given here for reference purposes. The first sector (see table on the following page) contains an optional software boot section. Disk controllers are often set up to bring track \emptyset , sector 1 into memory at a specific location (often location $\emptyset 0 \emptyset \theta H$). The program in this sector, called BOOT, has the responsibility of bringing the remaining sectors into memory starting at location 3400H+b. If your controller does not have a built-in sector load, you can ignore the program in track \emptyset , sector 1, and begin the load from track \emptyset sector 2 to location 3400H+b.

As an example, the Intel MDS-800 hardware cold start loader brings track 0, sector 1 into absolute address 3000H. Upon loading this sector, control transfers to location 3000H, where the bootstrap operation commences by loading the remainder of tracks 0, and all of track 1 into memory, starting at 3400H+b. The user should note that this bootstrap loader is of little use in a non-MDS environment, although it is useful to examine it since some of the boot actions will have to be duplicated in your cold start loader.

Track#	Sector#	Page#	Memory Address	CP/M Module name
ØØ	01		(boot address)	Cold Start Loader
00	Ø2	øø	3400H+b	CCP
••	03		348ØH+b	14 17
10 11	04	Ø1 "	3500H+b	
	05		3580H+b	
	Ø6	Ø2	3600H+b	
	Ø7 39		3680H+b 3700H+b	
	08 09	Ø 3	378ØH+b	
	10	Ø 4	3800H+b	5 0
	11		388ØH+b	
	12	Ø5	3900H+b	
н	13		3980H+b	
	14	ØG	3AØØH+b	
	15		3A8ØH+b	
	16	07	3BØØH+b	
ØØ	17		3B8ØH+b	CCP
 ØØ	18	 Ø8	3С00н+ь	BDOS
"	19	*	3C8ØH+b	"
••	20	Ø 9	3DØØH+b	••
**	21	**	3D8ØH+b	••
••	22	10	3EØØH+b	
••	23	••	3E8ØH+b	••
••	24	11	3FØØH+b	68
	25	"	3F8ØH+b	
	26	12	4000H+b	
Øl	01	6	4080H+b	
	Ø 2	13	4100H+b	
**	Ø3		418ØH+b	
	Ø 4	14	4200H+b	**
•*	Ø 5	•	428ØH+b	
••	Ø 6	15	4300H+b	
••	Ø7		438ØH+b	
	Ø8	16	4400H+b	
	Ø 9		4480H+b	
	10	17	4500H+b	и 4
	11		458ØH+b	
	12	18	4600H+b	
	13		4680H+b	
	14	19	4700H+b	
	15 16	20	4780H+b	
	17	20	4800H+b 4880H+b	
	18	21	4880H+D 4900H+b	
Øl	19	21	4980H+b	BDOS
ø1	20	22	4AØØH+b	BIOS
	21	11	4A8ØH+b	B105
64	23	23	4BØØH+b	
	24	2.5	4B8ØH+b	
**	25	24	4CØØH+b	54
Øl	26		4C8ØH+b	BIOS
Ø2-76	Ø1-26	** ** ** ** ** **		(directory and data)

(All Information Contained Herein is Proprietary to Digital Research.)

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6. THE BIOS ENTRY POINTS

The entry points into the BIOS from the cold start loader and BDOS are detailed below. Entry to the BIOS is through a "jump vector" located at 4A00H+b, as shown below (see Appendices B and C, as well). The jump vector is a sequence of 17 jump instructions which send program control to the individual BIOS subroutines. The BIOS subroutines may be empty for certain functions (i.e., they may contain a single RET operation) during regeneration of CP/M, but the entries must be present in the jump vector.

The jump vector at 4A00H+b takes the form shown below, where the individual jump addresses are given to the left:

4A00H+b	JMP BOOT	; ARRIVE HERE FROM COLD START LOAD
4AØ3H+b	JMP WBOOT	; ARRIVE HERE FOR WARM START
4A06H+b	JMP CONST	; CHECK FOR CONSOLE CHAR READY
4AØ9H+b	JMP CONIN	; READ CONSOLE CHARACTER IN
4AØCH+b	JMP CONOUT	; WRITE CONSOLE CHARACTER OUT
4A0FH+b	JMP LIST	; WRITE LISTING CHARACTER OUT
4A12H+b	JMP PUNCH	; WRITE CHARACTER TO PUNCH DEVICE
4A15H+b	JMP READER	; READ READER DEVICE
4A18H+b	JMP HOME	; MOVE TO TRACK 00 ON SELECTED DISK
4A1BH+b	JMP SELDSK	; SELECT DISK DRIVE
4AlEH+o	JMP SETTRK	; SET TRACK NUMBER
4A21H+b	JMP SETSEC	; SET SECTOR NUMBER
4A24H+b	JMP SETDMA	; SET DMA ADDRESS
4A27H+b	JMP READ	; READ SELECTED SECTOR
4A2AH+b	JMP WRITE	; WRITE SELECTED SECTOR
4A2DH+b	JMP LISTST	; RETURN LIST STATUS
4A3ØH+b	JMP SECTRAN	; SECTOR TRANSLATE SUBROUTINE

Each jump address corresponds to a particular subroutine which performs the specific function, as outlined below. There are three major divisions in the jump table: the system (re)initialization which results from calls on BOOT and WBOOT, simple character I/O performed by calls on CONST, CONIN, CONOUT, LIST, PUNCH, READER, and LISTST, and diskette I/O performed by calls on HOME, SELDSK, SETTRK, SETSEC, SETDMA, READ, WRITE, and SECTRAN.

All simple character I/O operations are assumed to be performed in ASCII, upper and lower case, with high order (parity bit) set to zero. An end-of-file condition for an input device is given by an ASCII control-z (IAH). Peripheral devices are seen by CP/M as "logical" devices, and are assigned to physical devices within the BIOS.

In order to operate, the BDOS needs only the CONST, CONIN, and CONOUT subroutines (LIST, PUNCH, and READER may be used by PIP, but not the BDOS). Further, the LISTST entry is used currently only by DESPOOL, and thus, the initial version of CBIOS may have empty subroutines for the remaining ASCII devices.

The characteristics of each device are

- CONSOLE The principal interactive console which communicates with the operator, accessed through CONST, CONIN, and CONOUT. Typically, the CONSOLE is a device such as a CRT or Teletype.
- LIST The principal listing device, if it exists on your system, which is usually a hard-copy device, such as a printer or Teletype.
- PUNCH The principal tape punching device, if it exists, which is normally a high-speed paper tape punch or Teletype.
- READER The principal tape reading device, such as a simple optical reader or Teletype.

Note that a single peripheral can be assigned as the LIST, PUNCH, and READER device simultaneously. If no peripheral device is assigned as the LIST, PUNCH, or READER device, the CBIOS created by the user may give an appropriate error message so that the system does not "hang" if the device is accessed by PIP or some other user program. Alternately, the PUNCH and LIST routines can just simply return, and the READER routine can return with a IAH (ctl-Z) in reg A to indicate immediate end-of-file.

For added flexibility, the user can optionally implement the "IOBYTE" function which allows reassignment of physical and logical devices. The function creates a mapping of logical to IOBYTE physical devices which can be altered during CP/M processing (see the STAT command). The definition of the IOBYTE function corresponds to the Intel standard as follows: a single location in memory (currently location 0003H) is maintained, called IOBYTE, which defines the logical to physical device mapping which is in effect at a particular time. The mapping is performed by splitting the IOBYTE into four distinct fields of two bits each, called the CONSOLE, READER, PUNCH, and LIST fields, as shown below:

		most signi	ficant	least	significant
IOBYTE AT	ØØØ3H	LIST	PUNCH	READER	CONSOLE
		bits 6,7	bits 4,5	bits 2,3	bits Ø,1

The value in each field can be in the range $\emptyset-3$, defining the assigned source or destination of each logical device. The values which can be assigned to each field are given below

CONSOLE field (bits 0,1) Ø - console is assigned to the console printer device (TTY:) 1 - console is assigned to the CRT device (CRT:) 2 - batch mode: use the READER as the CONSOLE input, and the LIST device as the CONSOLE output (BAT:) 3 - user defined console device (UCl:) READER field (bits 2.3) Ø - READER is the Teletype device (TTY:) 1 - READER is the high-speed reader device (RDR:) 2 - user defined reader # 1 (UR1:) 3 - user defined reader # 2 (UR2:) PUNCH field (bits 4,5) Ø - PUNCH is the Teletype device (TTY:) 1 - PUNCH is the high speed punch device (PUN:) 2 - user defined punch # 1 (UP1:) 3 - user defined punch # 2 (UP2:) LIST field (bits 6,7) 0 - LIST is the Teletype device (TTY:) 1 - LIST is the CRT device (CRT:) 2 - LIST is the line printer device (LPT:) 3 - user defined list device (UL1:)

> Note again that the implementation of the IOBYTE is optional, and affects only the organization of your CBIOS. No CP/M systems use the IOBYTE (although they tolerate the existence of the IOBYTE at location 0003H), except for PIP which allows access to the and STAT which allows physical devices. logical-physical assignments to be made and/or displayed (for more information, see the "CP/M Features and Facilities Guide"). In any case, the IOBYTE implementation should be omitted until your basic CBIOS is fully implemented and tested; then add the IOBYTE to increase your facilities.

> Disk I/O is always performed through a sequence of calls on the various disk access subroutines which set up the disk number to access, the track and sector on a particular disk, and the direct memory access (DMA) address involved in the I/O operation. After all these parameters have been set up, a call is made to the READ or WRITE function to perform the actual I/O operation. Note that there is often a single call to SELDSK to select a disk drive, followed by a number of read or write operations to the selected disk before selecting another drive for subsequent operations. Similarly, there may be a single call to set the DMA address, followed by several calls which read or write from the selected DMA address before the DMA address is changed. The track and sector subroutines are always called before the READ or WRITE operations are performed.

Note that the READ and WRITE routines should perform several retries (10 is standard) before reporting the error condition to the BDOS. If the error condition is returned to the BDOS, it will report the error to the user. The HOME subroutine may or may not actually perform the track 00 seek, depending upon your controller characteristics; the important point is that track 00 has been selected for the next operation, and is often treated in exactly the same manner as SETTRK with a parameter of 00.

The exact responsibilites of each entry point subroutine are given below:

BOOT

The BOOT entry point gets control from the cold start loader and is responsible for basic system initialization, including sending a signon message (which can be omitted in the first version). If the IOBYTE function is implemented, it must be set at this point. The various system parameters which are set by the WBOOT entry point must be initialized, and control is transferred to the CCP at 3400H+b for further processing. Note that reg C must be set to zero to select drive A.

WBOOT The WBOOT entry point gets control when a warm start occurs. A warm start is performed whenever a user program branches to location ØØØ0H, or when the CPU is reset from the front panel. The CP/M system must be loaded from the first two tracks of drive A up to, but not including, the BIOS (or CBIOS, if you have completed your patch). System parameters must be initialized as shown below:

> location 0,1,2 set to JMP WBOOT for warm starts (0000H: JMP 4A03H+b) location 3 set initial value of IOBYTE, if implemented in your CBIOS location 5,6,7 set to JMP BDOS, which is the primary entry point to CP/M for transient programs. (0005H: JMP 3C06H+b)

(see Section 9 for complete details of page zero use) Upon completion of the initialization, the WBOOT program must branch to the CCP at 3400H+b to (re)start the system. Upon entry to the CCP, register C is set to the drive to select after system initialization.

CONST Sample the status of the currently assigned console device and return ØFFH in register A if a character is ready to read, and ØØH in register A if no console characters are ready.

CONIN Read the next console character into register A, and

set the parity bit (high order bit) to zero. If no console character is ready, wait until a character is typed before returning.

- CONOUT Send the character from register C to the console output device. The character is in ASCII, with high order parity bit set to zero. You may want to include a time-out on a line feed or carriage return, if your console device requires some time interval at the end of the line (such as a TI Silent 700 terminal). You can, if you wish, filter out control characters which cause your console device to react in a strange way (a control-z causes the Lear Seigler terminal to clear the screen, for example).
- LIST Send the character from register C to the currently assigned listing device. The character is in ASCII with zero parity.
- PUNCH Send the character from register C to the currently assigned punch device. The character is in ASCII with zero parity.
- READER Read the next character from the currently assigned reader device into register A with zero parity (high order bit must be zero), an end of file condition is reported by returning an ASCII control-z (1AH).
- HOME Return the disk head of the currently selected disk (initially disk A) to the track ØØ position. If your controller allows access to the track Ø flag from the drive, step the head until the track Ø flag is detected. If your controller does not support this feature, you can translate the HOME call into a call on SETTRK with a parameter of Ø.
- Select the disk drive given by register C for further SELDSK operations, where register C contains Ø for drive A, 1 for drive B, and so-forth up to 15 for drive P (the standard CP/M distribution version supports four drives). On each disk select, SELDSK must return in HL the base address of a 16-byte area, called the Disk Parameter Header, described in the Section 10. For standard floppy disk drives, the contents of the header and associated tables does not change, and thus the program segment included in the sample CBIOS performs this operation automatically. If there is an attempt to select a non-existent drive, SELDSK returns HL=0000H as an error indicator. Although SELDSK must return the header address on each call, it is advisable to postpone the actual physical disk select operation until an I/O function (seek, read or write) is actually performed, since disk selects often occur without utimately performing any disk I/O, and many controllers will unload the head of the current disk

before selecting the new drive. This would cause an excessive amount of noise and disk wear.

- SETTRK Register BC contains the track number for subsequent disk accesses on the currently selected drive. You can choose to seek the selected track at this time, or delay the seek until the next read or write actually occurs. Register BC can take on values in the range Ø-76 corresponding to valid track numbers for standard floppy disk drives, and Ø-65535 for non-standard disk subsystems.
- SETSEC Register BC contains the sector number (1 through 26) for subsequent disk accesses on the currently selected drive. You can choose to send this information to the controller at this point, or instead delay sector selection until a read or write operation occurs.
- Register BC contains the DMA (disk memory access) SETDMA address for subsequent read or write operations. For example, if B = 00H and C = 80H when SETDMA is called, then all subsequent read operations read their data 80H through ØFFH, and all subsequent write into operations get their data from 80H through ØFFH, until The initial DMA the next call to SETDMA occurs. address is assumed to be 80H. Note that the controller need not actually support direct memory access. If, for example, all data is received and sent through I/O ports, the CBIOS which you construct will use the 128 byte area starting at the selected DMA address for the memory buffer during the following read or write operations.
- READ Assuming the drive has been selected, the track has been set, the sector has been set, and the DMA address has been specified, the READ subroutine attempts to read one sector based upon these parameters, and returns the following error codes in register A:
 - Ø no errors occurred
 - 1 non-recoverable error condition occurred

Currently, CP/M responds only to a zero or non-zero value as the return code. That is, if the value in register A is Ø then CP/M assumes that the disk operation completed properly. If an error occurs, however, the CBIOS should attempt at least 10 retries to see if the error is recoverable. When an error is reported the BDOS will print the message "BDOS ERR ON x: BAD SECTOR". The operator then has the option of typing <cr>> to ignore the error, or ctl-C to abort.

WRITE Write the data from the currently selected DMA address to the currently selected drive, track, and sector. The data should be marked as "non deleted data" to

maintain compatibility with other CP/M systems. The error codes given in the READ command are returned in register A, with error recovery attempts as described above.

- LISTST Return the ready status of the list device. Used by the DESPOOL program to improve console response during its operation. The value ØØ is returned in A if the list device is not ready to accept a character, and ØFFH if a character can be sent to the printer. Note that a ØØ value always suffices.
- SECTRAN Performs sector logical to physical sector translation in order to improve the overall response of CP/M. Standard CP/M systems are shipped with a "skew factor" 6, where six physical sectors are skipped between of each logical read operation. This skew factor allows enough time between sectors for most programs to load their buffers without missing the next sector. In particular computer systems which use fast processors, memory, and disk subsystems, the skew factor may be changed to improve overall response. Note, however, that you should maintain a single density IBM compatible version of CP/M for information transfer into and out of your computer system, using a skew factor of 6. In general, SECTRAN receives a logical sector number in BC, and a translate table address in The sector number is used as an index into the DE. translate table, with the resulting physical sector number in HL. For standard systems, the tables and indexing code is provided in the CBIOS and need not be changed.

(All Information Contained Herein is Proprietary to Digital Research.)

7. A SAMPLE BIOS

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The program shown in Appendix C can serve as a basis for your first BIOS. The simplest functions are assumed in this BIOS, so that you can enter it through the front panel, if absolutely necessary. Note that the user must alter and insert code into the subroutines for CONST, CONIN, CONOUT, READ, WRITE, and WAITIO subroutines. Storage is reserved for user-supplied code in these regions. The scratch area reserved in page zero (see Section 9) for the BIOS is used in this program, so that it could be implemented in ROM, if desired.

Once operational, this skeletal version can be enhanced to print the initial sign-on message and perform better error recovery. The subroutines for LIST, PUNCH, and READER can be filled-out, and the IOBYTE function can be implemented.

8. A SAMPLE COLD START LOADER

The program shown in Appendix D can serve as a basis for your cold start loader. The disk read function must be supplied by the user, and the program must be loaded somehow starting at location 0000. Note that space is reserved for your patch so that the total amount of storage required for the cold start loader is 128 bytes. Eventually, you will probably want to get this loader onto the first disk sector (track Ø, sector 1), and cause your controller to load it into memory automatically upon system start-up. Alternatively, you may wish to place the cold start loader into ROM, and place it above the CP/M In this case, it will be necessary to originate the program system. at a higher address, and key-in a jump instruction at system start-up which branches to the loader. Subsequent warm starts will not require this key-in operation, since the entry point 'WBOOT' gets control, thus bringing the system in from disk automatically. Note also that the skeletal cold start loader has minimal error recovery, which may be enhanced on later versions.

9. RESERVED LOCATIONS IN PAGE ZERO

Main memory page zero, between locations ØØH and ØFFH, contains several segments of code and data which are used during CP/M processing. The code and data areas are given below for reference purposes.

Locations	Contents
from to 0000H - 0002H	Contains a jump instruction to the warm start entry point at location 4A03H+b. This allows a simple programmed restart (JMP 0000H) or manual restart from the front panel.
0003H - 0003H	Contains the Intel standard IOBYTE, which is optionally included in the user's CBIOS, as described in Section 6.
0004H - 0004H	Current default drive number (\emptyset =A,,15=P).
0005H — 0007H	Contains a jump instruction to the BDOS, and serves two purposes: JMP 0005H provides the primary entry point to the BDOS, as described in the manual "CP/M Interface Guide," and LHLD 0006H brings the address field of the instruction to the HL register pair. This value is the lowest address in memory used by CP/M (assuming the CCP is being overlayed). Note that the DDT program will change the address field to reflect the reduced memory size in debug mode.
0008H - 0027H	(interrupt locations 1 through 5 not used)
0030H - 0037H	(interrupt location 6, not currently used - reserved)
ØØ38H - ØØ3AH	Restart 7 - Contains a jump instruction into the DDT or SID program when running in debug mode for programmed breakpoints, but is not otherwise used by CP/M.
003BH - 003FH	(not currently used - reserved)
0040H — 004FH	<pre>16 byte area reserved for scratch by CBIOS, but is not used for any purpose in the distribution version of CP/M</pre>
0050H - 005BH	(not currently used - reserved)
005CH - 007CH	default file control block produced for a transient program by the Console Command Processor.
007DH - 007FH	Optional default random record position

ØØ80H - ØØFFH default 128 byte disk buffer (also filled with the command line when a transient is loaded under the CCP).

Note that this information is set-up for normal operation under the CP/M system, but can be overwritten by a transient program if the BDOS facilities are not required by the transient.

If, for example, a particular program performs only simple I/O and must begin execution at location \emptyset , it can be first loaded into the TPA, using normal CP/M facilities, with a small memory move program which gets control when loaded (the memory move program must get control from location $\emptyset l \vartheta \theta H$, which is the assumed beginning of all transient programs). The move program can then proceed to move the entire memory image down to location \emptyset , and pass control to the starting address of the memory load. Note that if the BIOS is overwritten, or if location \emptyset (containing the warm start entry point) is overwritten, then the programmer must bring the CP/M system back into memory with a cold start sequence.

10. DISK PARAMETER TABLES.

Tables are included in the BIOS which describe the particular characteristics of the disk subsystem used with CP/M. These tables can be either hand-coded, as shown in the sample CBIOS in Appendix C, or automatically generated using the DISKDEF macro library, as shown in Appendix B. The purpose here is to describe the elements of these tables.

In general, each disk drive has an associated (16-byte) disk parameter header which both contains information about the disk drive and provides a scratchpad area for certain BDOS operations. The format of the disk parameter header for each drive is shown below

					Disk		Para	ameter	Header					
1	XLT	1	0000	1	0000	1	0000	DIRBUF	DPB	1	CSV	1	ALV	1
	16b		16b		16b		16b	16b	16b		16b		16b	

where each element is a word (16-bit) value. The meaning of each Disk Parameter Header (DPH) element is

- XLT Address of the logical to physical translation vector, if used for this particular drive, or the value 0000H if no sector translation takes place (i.e, the physical and logical sector numbers are the same). Disk drives with identical sector skew factors share the same translate tables.
- 0000 Scratchpad values for use within the BDOS (initial value is unimportant).
- DIRBUF Address of a 128 byte scratchpad area for directory operations within BDOS. All DPH's address the same scratchpad area.
- DPB Address of a disk parameter block for this drive. Drives with identical disk characteristics address the same disk parameter block.
- CSV Address of a scratchpad area used for software check for changed disks. This address is different for each DPH.
- ALV Address of a scratchpad area used by the BDOS to keep disk storage allocation information. This address is different for each DPH.

Given n disk drives, the DPH's are arranged in a table whose first row of 16 bytes corresponds to drive \emptyset , with the last row corresponding to drive n-1. The table thus appears as

DPBASE:

ØØ	XLT	ØØI	0000	1	0000	1	0000	DIRBUF DBP	ØØICSV	ØØ ALV	ØØI
Ø1	XLT	011	0000	1	0000	1	0000	DIRBUF DBP	Ø1 CSV	Ø1 ALV	Ø1
(and so-forth through)											
n-]	XLT	n-1	0000	1	0000	Ι	0000	DIRBUF DBP1	n-1 CSV1	n-1 ALV	n-1

where the label DPBASE defines the base address of the DPH table.

A responsibility of the SELDSK subroutine is to return the base address of the DPH for the selected drive. The following sequence of operations returns the table address, with a 0000H returned if the selected drive does not exist.

NDISKS	EQU	4 ;NUMBE	ER OF DISK DRIVES
SELDSK:			
	;SEL	ECT DISK GIV	EN BY BC
	LXI	н,0000н	;ERROR CODE
	MOV	A,C	;DRIVE OK?
	CPI	NDISKS	;CY IF SO
	RNC		;RET IF ERROR
	;NO	ERROR, CONTI	INUE
	MOV	L,C	;LOW(DISK)
	MOV	H,B	;HIGH(DISK)
	DAD	Н	;*2
	DAD	H	;*4
	DAD	Н	;*8
	DAD	Н	;*16
	LXI	D, DPBASE	FIRST DPH
	DAD	D	;DPH (DISK)
	RET		 A second support of the second support of support 12 mg/s

The translation vectors (XLT 00 through XLTn-1) are located elsewhere in the BIOS, and simply correspond one-for-one with the logical sector numbers zero through the sector count-1. The Disk Parameter Block (DPB) for each drive is more complex. A particular DPB, which is addressed by one or more DPH's, takes the general form

1	SPT	BSH	BLM	EXM	DSM	1	DRM	ALØ	AL1	CKS	1	OFF	1
	16b	8b	8b	8b	16b		16b	8b	8b	16b		16b	

where each is a byte or word value, as shown by the "8b" or "16b" indicator below the field.

SPT is the total number of sectors per track

BSH is the data allocation block shift factor, determined by the data block allocation size.

- EXM is the extent mask, determined by the data block allocation size and the number of disk blocks.
- DSM determines the total storage capacity of the disk drive
- DRM determines the total number of directory entries which can be stored on this drive ALØ,AL1 determine reserved directory blocks.
- CKS is the size of the directory check vector
- OFF is the number of reserved tracks at the beginning of the (logical) disk.

The values of BSH and BLM determine (implicitly) the data allocation size BLS, which is not an entry in the disk parameter block. Given that the designer has selected a value for BLS, the values of BSH and BLM are shown in the table below

BLS	BSH	BLM
1,024	3	7
2,048	4	15
4,096	5	31
8,192	6	63
16.384	7	127

where all values are in decimal. The value of EXM depends upon both the BLS and whether the DSM value is less than 256 or greater than 255, as shown in the following table

BLS	DSM <	256	DSM	>	255
1,024	Ø		ľ	N/A	4
2,048	1			Ø	
4,096	3			1	
8,192	7			3	
16,384	15			7	

The value of DSM is the maximum data block number supported by this particular drive, measured in BLS units. The product BLS times (DSM+1) is the total number of bytes held by the drive and, of course, must be within the capacity of the physical disk, not counting the reserved operating system tracks.

The DRM entry is the one less than the total number of directory entries, which can take on a 16-bit value. The values of ALØ and AL1, however, are determined by DRM. The two values ALØ and AL1 can together be considered a string of 16-bits, as shown below.

I							A	LØ								I						A	L1							I
1	1		1	-			. –	1		1		1	-	1		1	1	_	I		1		1		1			1		I
	-	ø 1		ø	2	ø	3		4	ø	5	ø	6	ø	7	ø	8	Øg)	10		.1	1	2	1	3	1	4	1!	5

where position 00 corresponds to the high order bit of the byte labelled AL0, and 15 corresponds to the low order bit of the byte labelled AL1. Each bit position reserves a data block for number of directory entries, thus allowing a total of 16 data blocks to be assigned for directory entries (bits are assigned starting at 00 and filled to the right until position 15). Each directory entry occupies 32 bytes, resulting in the following table

BLS	Dire	Entries				
1,024	32	times	#	bits		
2,048	64	times	#	bits		
4,096	128	times	#	bits		
8,192	256	times	#	bits		
16,384	512	times	#	bits		

Thus, if DRM = 127 (128 directory entries), and BLS = 1024, then there are 32 directory entries per block, requiring 4 reserved blocks. In this case, the 4 high order bits of AL0 are set, resulting in the values AL0 = 0F0H and AL1 = 00H.

The CKS value is determined as follows: if the disk drive media is removable, then CKS = (DRM+1)/4, where DRM is the last directory entry number. If the media is fixed, then set CKS = Ø (no directory records are checked in this case).

Finally, the OFF field determines the number of tracks which are skipped at the beginning of the physical disk. This value is automatically added whenever SETTRK is called, and can be used as a mechanism for skipping reserved operating system tracks, or for partitioning a large disk into smaller segmented sections.

To complete the discussion of the DPB, recall that several DPH's can address the same DPB if their drive characteristics are identical. Further, the DPB can be dynamically changed when a new drive is addressed by simply changing the pointer in the DPH since the BDOS copies the DPB values to a local area whenever the SELDSK function is invoked.

Returning back to the DPH for a particular drive, note that the two address values CSV and ALV remain. Both addresses reference an area of uninitialized memory following the BIOS. The areas must be unique for each drive, and the size of each area is determined by the values in the DPB.

The size of the area addressed by CSV is CKS bytes, which is sufficient to hold the directory check information for this particular drive. If CKS = (DRM+1)/4, then you must reserve (DRM+1)/4 bytes for directory check use. If CKS = 0, then no storage is reserved.

The size of the area addressed by ALV is determined by the maximum number of data blocks allowed for this particular disk, and is computed as (DSM/8)+1.

The CBIOS shown in Appendix C demonstrates an instance of these tables for standard 8" single density drives. It may be useful to examine this program, and compare the tabular values with the definitions given above.

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11. THE DISKDEF MACRO LIBRARY.

A macro library is shown in Appendix F, called DISKDEF, which greatly simplifies the table construction process. You must have access to the MAC macro assembler, of course, to use the DISKDEF facility, while the macro library is included with all CP/M 2.0 distribution disks.

A BIOS disk definition consists of the following sequence of macro statements:

MACLIB DISKDEF DISKS n DISKDEF Ø,... DISKDEF 1,... DISKDEF n-1 ENDEF

where the MACLIB statement loads the DISKDEF.LIB file (on the same disk as your BIOS) into MAC's internal tables. The DISKS macro call follows, which specifies the number of drives to be configured with your system, where n is an integer in the range 1 to 16. A series of DISKDEF macro calls then follow which define the characteristics of each logical disk, Ø through n-1 (corresponding to logical drives A through P). Note that the DISKS and DISKDEF macros generate the in-line fixed data tables described in the previous section, and thus must be placed in a non-executable portion of your BIOS, typically directly following the BIOS jump vector.

The remaining portion of your BIOS is defined following the DISKDEF macros, with the ENDEF macro call immediately preceding the END statement. The ENDEF (End of Diskdef) macro generates the necessary uninitialized RAM areas which are located in memory above your BIOS.

The form of the DISKDEF macro call is

DISKDEF dn,fsc,lsc,[skf],bls,dks,dir,cks,ofs,[0]

where

dn is the logical disk number, Ø to n-1 fsc is the first physical sector number (0 or 1) lsc is the last sector number skf is the optional sector skew factor bls is the data allocation block size is the number of directory entries dir is the number of "checked" directory entries cks ofs is the track offset to logical track ØØ [Ø] is an optional 1.4 compatibility flag

The value "dn" is the drive number being defined with this DISKDEF

macro invocation. The "fsc" parameter accounts for differing sector numbering systems, and is usually Ø or 1. The "lsc" is the last numbered sector on a track. When present, the "skf" parameter defines the sector skew factor which is used to create a sector translation table according to the skew. If the number of sectors is less than 256, a single-byte table is created, otherwise each translation table element occupies two bytes. No translation table is created if the skf parameter is omitted (or equal to Ø). The "bls" parameter specifies the number of bytes allocated to each data block, and takes on the values 1024, 2048, 4096, 8192, or 16384. Generally, performance increases with larger data block sizes since there are fewer directory references and logically connected data records are physically close on the disk. Further, each directory entry addresses more data and the BIOS-resident ram space is reduced. The "dks" specifies the total disk size in "bls" units. That is, if the bls 2048 and dks = 1000, then the total disk capacity is 2,048,000 bytes. If dks is greater than 255, then the block size parameter bls must be The value of "dir" is the total number of greater than 1024. directory entries which may exceed 255, if desired. The "cks" parameter determines the number of directory items to check on each directory scan, and is used internally to detect changed disks during system operation, where an intervening cold or warm start has not occurred (when this situation is detected, CP/M automatically marks the disk read/only so that data is not subsequently destroyed). As stated in the previous section, the value of cks = dir when the media is easily changed, as is the case with a floppy disk subsystem. If the disk is permanently mounted, then the value of cks is typically Ø, since the probability of changing disks without a restart is quite low. The "ofs" value determines the number of tracks to skip when this particular drive is addressed, which can be used to reserve additional operating system space or to simulate several logical drives on a single large capacity physical drive. Finally, the $[\emptyset]$ parameter is included when file compatibility is required with versions of 1.4 which have been modified for higher density disks. This parameter ensures that only 16K is allocated for each directory record, as was the case for previous versions. Normally, this parameter is not included.

For convenience and economy of table space, the special form

DISKDEF i,j

gives disk i the same characteristics as a previously defined drive j. A standard four-drive single density system, which is compatible with version 1.4, is defined using the following macro invocations:

DISKS 4 DISKDEF Ø,1,26,6,1024,243,64,64,2 DISKDEF 1,0 DISKDEF 2,0 DISKDEF 3,0 ENDEF

with all disks having the same parameter values of 26 sectors per track (numbered 1 through 26), with 6 sectors skipped between each access, 1024 bytes per data block, 243 data blocks for a total of 243k byte disk capacity, 64 checked directory entries, and two operating system tracks.

The DISKS macro generates n Disk Parameter Headers (DPH's), starting at the DPH table address DPBASE generated by the macro. Each disk header block contains sixteen bytes, as described above, and correspond one-for-one to each of the defined drives. In the four drive standard system, for example, the DISKS macro generates a table of the form:

DPBASE	EQU	Ş
DPEØ:	DW	XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV0,ALV0
DPE1:	DW	XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV1,ALV1
DPE2:	DW	XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV2,ALV2
DPE3:	DW	XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV3,ALV3

where the DPH labels are included for reference purposes to show the beginning table addresses for each drive Ø through 3. The values contained within the disk parameter header are described in detail in the previous section. The check and allocation vector addresses are generated by the ENDEF macro in the ram area following the BIOS code and tables.

Note that if the "skf" (skew factor) parameter is omitted (or equal to \emptyset), the translation table is omitted, and a $\emptyset\emptyset\emptyset\emptyset$ H value is inserted in the XLT position of the disk parameter header for the disk. In a subsequent call to perform the logical to physical translation, SECTRAN receives a translation table address of DE = $\emptyset\emptyset\emptyset\emptyset$ H, and simply returns the original logical sector from BC in the HL register pair. A translate table is constructed when the skf parameter is present, and the (non-zero) table address is placed into the corresponding DPH's. The table shown below, for example, is constructed when the standard skew factor skf = 6 is specified in the DISKDEF macro call:

XLTØ: DB 1,7,13,19,25,5,11,17,23,3,9,15,21 DB 2,8,14,20,26,6,12,18,24,4,10,16,22

Following the ENDEF macro call, a number of uninitialized data areas are defined. These data areas need not be a part of the BIOS which is loaded upon cold start, but must be available between the BIOS and the end of memory. The size of the uninitialized RAM area is determined by EQU statements generated by the ENDEF macro. For a standard four-drive system, the ENDEF macro might produce

4C72 =	BEGDAT EQU Ş
	(data areas)
$4DB\emptyset =$	ENDDAT EQU \$
Ø13C =	DATSIZ EQU \$-BEGDAT

which indicates that uninitialized RAM begins at location 4C72H, ends at 4DB0H-1, and occupies 013CH bytes. You must ensure that these addresses are free for use after the system is loaded.

After modification, you can use the STAT program to check your drive characteristics, since STAT uses the disk parameter block to decode the drive information. The STAT command form

STAT d:DSK:

decodes the disk parameter block for drive d (d=A,...,P) and displays the values shown below:

r: 128 Byte Record Capacity
k: Kilobyte Drive Capacity
d: 32 Byte Directory Entries
c: Checked Directory Entries
e: Records/ Extent
b: Records/ Block
s: Sectors/ Track
t: Reserved Tracks

Three examples of DISKDEF macro invocations are shown below with corresponding STAT parameter values (the last produces a full 8-megabyte system).

DISKDEF 0,1,58,,2048,256,128,128,2 r=4096, k=512, d=128, c=128, e=256, b=16, s=58, t=2

DISKDEF 0,1,58,,2048,1024,300,0,2 r=16384, k=2048, d=300, c=0, e=128, b=16, s=58, t=2

DISKDEF Ø,1,58,,16384,512,128,128,2 r=65536, k=8192, d=128, c=128, e=1024, b=128, s=58, t=2

12. SECTOR BLOCKING AND DEBLOCKING.

Upon each call to the BIOS WRITE entry point, the CP/M BDOS includes information which allows effective sector blocking and deblocking where the host disk subsystem has a sector size which is a multiple of the basic 128-byte unit. The purpose here is to present a general-purpose algorithm which can be included within your BIOS which uses the BDOS information to perform the operations automatically.

Upon each call to WRITE, the BDOS provides the following information in register C:

Ø = normal sector write
1 = write to directory sector
2 = write to the first sector
of a new data block

Condition Ø occurs whenever the next write operation is into a previously written area, such as a random mode record update, when the write is to other than the first sector of an unallocated block, or when the write is not into the directory area. Condition 1 occurs when a write into the directory area is performed. Condition 2 occurs when the first record (only) of a newly allocated data block is written. In most cases, application programs read or write multiple 128 byte sectors in sequence, and thus there is little overhead involved in either operation when blocking and deblocking records since pre-read operations can be avoided when writing records.

Appendix G lists the blocking and deblocking algorithms in skeletal form (this file is included on your CP/M disk). Generally, the algorithms map all CP/M sector read operations onto the host disk through an intermediate buffer which is the size of the host disk sector. Throughout the program, values and variables which relate to the CP/M sector involved in a seek operation are prefixed by "sek," while those related to the host disk system are prefixed by "hst." The equate statements beginning on line 29 of Appendix G define the mapping between CP/M and the host system, and must be changed if other than the sample host system is involved.

The entry points BOOT and WBOOT must contain the initialization code starting on line 57, while the SELDSK entry point must be augmented by the code starting on line 65. Note that although the SELDSK entry point computes and returns the Disk Parameter Header address, it does not physically selected the host disk at this point (it is selected later at READHST or WRITEHST). Further, SETTRK, SETTRK, and SETDMA simply store the values, but do not take any other action at this point. SECTRAN performs a trivial trivial function of returning the physical sector number.

The principal entry points are READ and WRITE, starting on lines 110 and 125, respectively. These subroutines take the place of your previous READ and WRITE operations.

The actual physical read or write takes place at either WRITEHST or READHST, where all values have been prepared: hstdsk is the host

disk number, hsttrk is the host track number, and hstsec is the host sector number (which may require translation to a physical sector number). You must insert code at this point which performs the full host sector read or write into, or out of, the buffer at hstbuf of length hstsiz. All other mapping functions are performed by the algorithms.

This particular algorithm was tested using an 80 megabyte hard disk unit which was originally configured for 128 byte sectors, producing approximately 35 megabytes of formatted storage. When configured for 512 byte host sectors, usable storage increased to 57 megabytes, with a corresponding 400% improvement in overall response. In this situation, there is no apparent overhead involved in deblocking sectors, with the advantage that user programs still maintain the (less memory consuming) 128-byte sectors. This is primarily due, of course, to the information provided by the BDOS which eliminates the necessity for pre-read operations to take place.

APPENDIX A: THE MDS COLD START LOADER

			MDS-800	Cold Sta	art Load	er for CP/M 2.0
		;				
		;	Version	2.Ø Augu	ist, 1979	9
		;		~		
0000		false	equ	Ø		
ffff Ø000		true testing	equ	not fals false	se	
0000		;	equ	Laise		
		,	if	testing		
		bias	equ	Ø3400h		
			endif			
			if	not test	ting	
0000	=	bias	equ	0000h		
			endif	hinn		there of dealord
ØØØØ Ø8Ø6		cpmb bdos	equ	bias 806h+bia	20	;base of dos load ;entry to dos for calls
1880		bdose	equ equ	1880h+b:		;end of dos load
1600		boot	equ	1600h+b:		; cold start entry point
1603		rboot	equ	boot+3	140	;warm start entry point
		;	- 1	2.2.2.2.2		
3000			org	3000h	;loaded	here by hardware
		;				
1880		bdosl	equ	bdose-c	omb	
0002		ntrks	egu	2 bdog1/1	20	;tracks to read ;# sectors in bdos
ØØ31 ØØ19		bdoss bdosØ	egu egu	bdos1/1: 25	20	;# on track Ø
0018		bdosl	equ	bdoss-b	dosØ	;# on track 1
0010		;	CHU	20000 2		/
£800	=	mon8Ø	equ	Ø£8ØØh	; intel n	monitor base
ffØf		rmon8Ø	equ	ØffØfh		t location for mon80
0078		base	equ	Ø78h		used by controller
0079		rtype	equ	base+1	;result	
007b 007f		rbyte	equ	base+3		
00/1		reset	equ	base+7	;reset (controller
0078	=	dstat	equ	base	disk s	tatus port
0079		ilow	equ	base+1	sensitive and a second second second	pb address
ØØ7a	=	ihigh	egu	base+2		opb address
ØØff	=	bsw	equ	Øffh	;boot s	witch
0003		recal	equ	3h		brate selected drive
0004		readf	egu	4h		ead function
0100	=	stack	equ	100h	;use en	d of boot for stack
		; rstart:				
3000	310001	istart.	lxi	sp.stac	kin cas	e of call to mon8Ø
		;		isk state		
3003	db79		in	rtype	an 1127	
3005	db7b		in	rbyte	32.03 AV4 4	
		1		f boot st	witch is	off
2007	abee	coldsta		h m s		
	dbff		in ani	bsw Ø2b	·cwitch	0.02
3008	820730		ani jnz	čőldsta	rt ^{switch}	011.

clear the controller ; 300e d37f out reset ;logic cleared ; ; 3010 0602 b, ntrks ; number of tracks to read mvi 3012 214230 lxi h.iopbØ ; start: ; read first/next track into cpmb ; 3Ø15 7d a.1 mov 3016 d379 ilow out 3018 7c a,h mov 3019 d37a out ihigh 301b db78 wait0: in dstat 3014 e60430 ani 4 waitØ ; check disk status ; 3022 db79 in rtype 3024 e603 11b ani 3026 fe02 2 cpi ; if testing ;go to monitor if 11 or 10 cnc rmon80 endif if not testing 3028 d20030 jnc rstart ; retry the load endif ; 302b db7b ; i/o complete, check status in rbvte if not ready, then go to mon80 ; 302d 17 ral 302e dc0fff CC rmon80 ;not ready bit set 3Ø31 1f ;restore rar 3Ø32 e61e ani 11110b ;overrun/addr err/seek/crc ; if testing cnz rmon8Ø ;go to monitor endif if not testing 3034 c20030 jnz rstart ; retry the load endif ; ; 3037 110700 lxi d, iopbl ; length of iopb 3Ø3a 19 dad d ;addressing next iopb 303b 05 dcr b ; count down tracks 303c c21530 jnz start ; ; jmp boot, print message, set-up jmps ; 303f c30016 jmp boot ; parameter blocks ;

6110

3042	80	iopb0:	db	8Øh	;iocw, no update
3043	Ø 4		db	readf	;read function
3044	19		db	bdosØ	;# sectors to read trk Ø
3045	ØØ		db	Ø	;track Ø
3046	Ø 2		db	2	;start with sector 2, trk Ø
3047	0000		dw	cpmb	;start at base of bdos
0007	=	iopbl	egu	\$-iopb0	
		;			
3049	80	iopbl:	db	80h	
304a	Ø4		đb	readf	
304b	18		db	bdosl	;sectors to read on track 1
304c	Ø1		đb	1	;track 1
304d	Øl		db	1	;sector 1
304e	800c		dw	cpmb+bdo	psØ*128 ;base of second rd
3050			end	8 7 .)	

APPENDIX B: THE MDS BASIC I/O SYSTEM (BIOS)

	;			vers for cp/m 2.0 (le density version)
	;	version	2.0 augu	ast, 1979
ØØ14 =	vers	equ	20	;version 2.0
	;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;	digital box 579	nt (c) 19 research , pacific nia, 9395	c grove
$\begin{array}{rcl} 4a00 \\ 3400 = \\ 3c06 = \\ 1600 = \\ 002c = \\ 0002 = \\ 0004 = \\ 0080 = \\ 000a = \\ 000a = \\ \end{array}$	cpmb bdos cpml nsects offset cdisk buff retry	org egu egu egu egu egu egu	\$-cpmb	;base of bdos in 20k system ;length (in bytes) of cpm system 3;number of sectors to load ;number of disk tracks used by cp
	, , , , , , , , , , , , , .	boot wboot (boot and const conin conout list	<pre>cold sta warm sta nd wboot console reg-a = reg-a = console console list out punch ou paper ta</pre>	art (save i/o byte) are the same for mds)
	, , , , , , , , , , , , , , , , , , , ,	mds, wh seldsk settrk setsec	ich is us select o set trac set sec	calls set-up the io parameter bloc sed to perform subsequent reads an disk given by reg-c (0,1,2) ck address (0,76) for sub r/w tor address (1,,26) sequent dma address (initially 80h
	;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;	read/wr read write	read tra	ne previous calls to set i/o parms ack/sector to preset dma address rack/sector from preset dma addres
4a00 c3b34a 4a03 c3c34a 4a06 c3614b 4a09 c3644b 4a0c c36a4b		jmp	ctor for boot wboot const conin conout	indiviual routines

4a0f c36d4b 4a12 c3724b 4a15 c3754b 4a18 c3784b 4a1b c37d4b 4a1e c3a74b 4a21 c3ac4b 4a21 c3ac4b 4a24 c3bb4b 4a27 c3c14b 4a2a c3ca4b 4a2d c3704b 4a30 c3b14b		jmp jmp jmp jmp jmp jmp jmp jmp jmp jmp	list punch reader home seldsk settrk setsec setdma read write listst sectran	;list	status
4a33+= 4a33+824a00 4a37+000000 4a3b+6e4c73 4a3f+0d4dee 4a43+824a00 4a47+000000 4a4b+6e4c73 4a4f+3c4d1d 4a53+824a00 4a55+6e4c73 4a5f+6b4d4c 4a63+824a00 4a67+000000 4a6b+6e4c73 4a6f+9a4d7b	dpel: dpe2:	maclib disks equ dw dw dw dw dw dw dw dw dw dw dw dw dw	4	; four ; base Øh ØØh pbØ Ø Ø b 0 b 0 b 0 b 2 Ø h ØØh pb2 2 Ø h ØØh pb3	<pre>the disk definition library disks of disk parameter blocks ;translate table ;scratch area ;dir buff,parm block ;check, alloc vectors ;translate table ;scratch area ;dir buff,parm block ;check, alloc vectors</pre>
4a73+= 4a73+1a00 4a75+03 4a76+07 4a77+00 4a77+00 4a78+f200 4a7c+c0 4a7c+c0 4a7c+c0 4a7c+1000 4a80+0200 4a80+0200 4a82+= 4a82+01 4a83+07 4a84+0d 4a85+13 4a86+19 4a87+05 4a88+0b 4a89+11 4a8a+17 4a8b+03	dpb0 xlt0	diskdef equ dw db db db db db db db db db db db db db	Ø,1,26,6 \$ 26 3 7 Ø 242 63 192 Ø 16 2 \$ 1 7 13 19 25 5 11 17 23 3	,1024,	,243,64,64,offset ;disk parm block ;sec per track ;block shift ;block mask ;extnt mask ;disk size-1 ;directory max ;allocØ ; 1loc1 ;check size ;offset ;translate table

1-0-100		21-	0	
4a8c+09		db	9	
4a8d+Øf		db	15	
4a8e+15		db	21	
4a8f+02		db	2	
4a90+08		đb	8	
4a91+0e		db	14 2Ø	
4a92+14		db db	26	
4a93+1a 4a94+06		db	6	
4a94+06 4a95+0c		db	12	
4a95+bC 4a96+12		db	18	
4a97+18		db	24	
4a98+04		db	4	
4a99+0a		db	10	
4a9a+10		db	16	
4a9b+16		db	22	
4430110		diskdef		
4a73+=	dpbl	equ	dpbØ	;equivalent parameters
001f+=	alsl	egu	alsØ	;same allocation vector size
ØØ1Ø+=	cssl	equ	cssØ	;same checksum vector size
4a82 + =	xltl	equ	xltØ	same translate table
		diskdef		
4a73+=	dpb2	egu	dpbØ	;equivalent parameters
00lf + =	als2	egu	alsØ	;same allocation vector size
ØØ1Ø+=	css2	equ	cssØ	;same checksum vector size
4a82+=	xlt2	egu	xltØ	;same translate table
		diskdef		
4a73+=	dpb3	equ	dpbØ	;equivalent parameters
001f+=	als3	equ	alsØ	;same allocation vector size
0010+=	css3	equ	cssØ	;same checksum vector size
4a82+=	xlt3	equ	xltØ	;same translate table
	;	ender o	ccurs at	end of assembly
	;	7 6		
	;			er - independent code, the remaini
	7			the particular operating environm
	;	be alte	rea for a	any system which differs from the
		the fol	lowing a	ode assumes the mds monitor exists
				o subroutines within the monitor
		and use	5 the 1/	o subloatines within the monitor
	,	we also	assume	the mds system has four disk drive
ØØfd =	revrt	equ	Øfdh	; interrupt revert port
ØØfc =	intc	equ	Øfch	; interrupt mask port
ØØ£3 =	icon	equ	Øf3h	; interrupt control port
007e =	inte	equ		10b; enable rst 0(warm boot), rst 7
	;	el 20		
	;	mds mon	itor equa	ates
f800 =	mon8Ø	egu	Ø£800h	;mds monitor
fføf =	rmon8Ø	equ	ØffØfh	;restart mon80 (boot error)
f8Ø3 =	ci	equ	Ø£8Ø3h	; console character to reg-a
£806 =	ri	equ	Øf8Ø6h	;reader in to reg-a
f809 =	co	equ	Øf809h	; console char from c to console o
f80c =	ро	equ	Øf8Øch	; punch char from c to punch devic
f80f =	10	equ	Øf8Øfh	;list from c to list device
f812 =	csts	equ	Ø£812h	;console status 00/ff to register

; disk ports and commands ; ;base of disk command io ports 0078 =78h base equ base :disk status (input) 0078 =dstat equ ;result type (input) base+1 0079 =equ rtype ;result byte (input) base+3 007b =rbyte equ ; iopb low address (output) 0079 =ilow eau base+1 007a =ihigh equ base+2 ; iopb high address (output) ; 4h ;read function 0004 = readf equ 6h :write function 0006 =writf equ 3h :recalibrate drive 0003 =recal equ 0004 =iordy equ 4h ;i/o finished mask Ødh 000d =;carriage return cr equ lf Øah ;line feed 000a = equ ; signon: ;signon message: xxk cp/m vers y.y 4a9c ØdØaØa db cr,lf,lf 4a9f 3230 db '20' ;sample memory size 'k cp/m vers ' 4aal 6b2043f db vers/10+'0','.',vers mod 10+'0' 4aad 322e30 db 4abØ ØdØaØØ đb cr,lf,Ø boot: ; print signon message and go to ccp (note: mds boot initialized iobyte at 0003h) ; 4ab3 310001 lxi sp,buff+80h 4ab6 219c4a lxi h,signon 4ab9 cdd34b call prmsq ;print message 4abc af xra ;clear accumulator a 4abd 320400 sta cdisk ;set initially to disk a 4acØ c3Øf4b jmp gocpm ; go to cp/m ; ; wboot:; loader on track Ø, sector 1, which will be skippe read cp/m from disk - assuming there is a 128 byt ; start. ; ; 4ac3 318000 lxi sp, buff ; using dma - thus 80 thru ff ok f ; 4ac6 ØeØa mvi c, retry ; max retries 4ac8 c5 push b wboot0: ;enter here on error retries 4ac9 010034 lxi b,cpmb ;set dma address to start of disk 4acc cdbb4b call setdma 4acf ØeØØ mvi C,0 ;boot from drive Ø seldsk 4adl cd7d4b call 4ad4 ØeØØ mvi C.0 4ad6 cda74b call settrk ;start with track Ø 4ad9 ØeØ2 mvi c,2 ;start reading sector 2 4adb cdac4b call setsec ; read sectors, count nsects to zero ; 4ade cl ;10-error count pop b 4adf Ø62c mvi b,nsects

r.,	dsec: ;read ne	ext secto	r
4ael c5	push		;save sector count
4ae2 cdcl4b		read	
4ae5 c2494b		영상가 이 맛 안 봐야 한 것이다.	retry if errors occur
4ae8 2a6c4c	-		; increment dma address
4aeb 118000	lxi	d,128	;sector size
4aee 19	dad	d	; incremented dma address in hl
4aef 44	mov	b,h	
4afØ 4d			;ready for call to set dma
4afl cdbb4b	call	setdma	
4af4 3a6b4c			;sector number just read
4af7 fela	cpi	26	;read last sector?
4af9 da054b	jc	rdl	· · · · · · · · · · · · · · · · · · ·
;			6, zero and go to next track
4afc 3a6a4c	lda	iot	;get track to register a
4aff 3c	inr	a	mander for an 11
4b00 4f			;ready for call
4b0l cda74b 4b04 af	call	settrk	aloor costor number
	dl: inr	a a	clear sector number; to next sector
4b06 4f	mov		; ready for call
4b07 cdac4b	call	setsec	, icady for call
4b0a cl	gog		;recall sector count
4bøb Ø5	der	b	;done?
4b0c c2el4a	jnz	rdsec	 Relation to the set
;	-		
;	done wit		ad, reset default buffer address
g			om cold start boot)
;		stØ and	rst7
4bØf f3	di	(12) - 202	
4b10 3e12	mvi	a,12h	; initialize command
4b12 d3fd	out	revrt	
4bl4 af	xra	a	
4b15 d3fc	out		;cleared
4b17 3e7e			;rst0 and rst7 bits on
4bl9 d3fc 4blb af	out	intc	
4blc d3f3	xra out	a icon	;interrupt control
		icon	, incertape concroi
;		ault buff	er address to 80h
4ble 018000	lxi	b,buff	
4b21 cdbb4b	call	setdma	
;	G Register		
;	reset mo	onitor en	try points
4b24 3ec3	mvi	a,jmp	Selector Contraction (Contraction)
4b26 320000	sta	Ø	
4b29 21034a	lxi	h, wboote	
4b2c 220100	shld	1	;jmp wboot at location 00
4b2f 320500	sta	5	
4b32 21063c	lxi	h,bdos	
4b35 220600	shld	6	;jmp bdos at location 5
4b38 323800	sta	7*8	;jmp to mon80 (may have been chan
4b3b 2100f8	lxi	h,mon8Ø	
4b3e 223900	shld	7*8+1	
;	Teave 10	obyte set	

previously selected disk was b, send parameter to cdisk ;last logged disk number 4b41 3a0400 lda ;send to ccp to log it in 4b44 4f mov c,a 4b45 fb ei cpmb 4b46 c30034 jmp ; error condition occurred, print message and retry ; booterr: 4b49 cl b :recall counts pop dcr 4b4a Ød С 4b4b ca524b jΖ booterØ try again ; 4b4e c5 push b 4b4f c3c94a jmp wbootØ booter0: otherwise too many retries ; h,bootmsg 4b52 215b4b lxi 4b55 cdd34b call prmsq rmon80 ;mds hardware monitor 4b58 c30fff jmp bootmsg: 4b5b 3f626f4 '?boot',Ø db ; ; ; console status to reg-a const: (exactly the same as mds call) ; 4b61 c312f8 jmp csts ; ; console character to reg-a conin: 4b64 cd03f8 call ci 7fh 4b67 e67f ani ; remove parity bit 4b69 c9 ret ; conout: ; console character from c to console out 4b6a c309f8 jmp CO list: ;list device out (exactly the same as mds call) ; 4b6d c30ff8 jmp 10 listst: ;return list status 4b70 af xra a 4b71 c9 ret ;always not ready ; ; punch device out punch: (exactly the same as mds call) ; 4b72 c30cf8 jmp po ; reader: ;reader character in to reg-a (exactly the same as mds call) ; 4b75 c306f8 jmp ri home: ; move to home position

treat as track 00 seek ; 4b78 ØeØØ mvi c.0 4b7a c3a74b jmp settrk seldsk: ;select disk given by register c 4b7d 210000 lxi h.0000h ;return 0000 if error 4b8Ø 79 mov a.c 4b81 feØ4 cpi ndisks ;too large? 4b83 dØ ; leave hl = 0000rnc ; 4b84 e602 ani 10b ;00 00 for drive 0,1 and 10 10 fo 4b86 32664c :to select drive bank dbank sta 4b89 79 ;00, 01, 10, 11 mov a,c 4b8a e601 ani lb ;mds has 0,1 at 78, 2,3 at 88 ;result 00? 4b8c b7 ora а 4b8d ca924b İΖ setdrive 4b90 3e30 a,00110000b ;selects drive 1 in bank mvi setdrive: 4b92 47 mov b,a ;save the function 4b93 21684c lxi h,iof ; io function 4b96 7e mov a,m 4b97 e6cf ani 11001111b ;mask out disk number 4b99 bØ :mask in new disk number ora b 4b9a 77 ;save it in iopb m,a mov 4892 5800 h:6 moy :hl=disk number ;*2 4b9e 29 dad h ;*4 4b9f 29 dad h 4baØ 29 ;*8 dad h 4bal 29 dad ;*16 h 4ba2 11334a lxi d,dpbase 4ba5 19 dad d ;hl=disk header table address 4ba6 c9 ret ; ; settrk: ;set track address given by c 4ba7 216a4c lxi h.iot 4baa 71 mov m,c 4bab c9 ret ; setsec: ;set sector number given by c 4bac 216b4c lxi h,ios 4baf 71 mov m,c 4bb0 c9 ret sectran: ;translate sector bc using table at de 4bbl 0600 mvi b,Ø ;double precision sector number i 4bb3 eb xchq ;translate table address to hl 4bb4 Ø9 dad b ;translate(sector) address 4bb5 7e mov ;translated sector number to a a,m 4bb6 326b4c sta ios 4669 6f 1,a ;return sector number in 1 moy ; setdma: ;set dma address given by regs b.c

4bbb 69 mov 1.c h.b 4bbc 60 mov 4bbd 226c4c shld iod 4bc0 c9 ret ; ;read next disk record (assuming disk/trk/sec/dma read: c, readf ; set to read function 4bcl ØeØ4 mvi call setfunc 4bc3 cde04b ;perform read function waitio 4bc6 cdfØ4b call ;may have error set in reg-a 4bc9 c9 ret ; ; ;disk write function write: c.writf 4bca ØeØ6 mvi 4bcc cdeØ4b call setfunc ; set to write function 4bcf cdfØ4b call waitio 4bd2 c9 ret ;may have error set ; ; utility subroutines ; ;print message at h,l to Ø prmsq: 4bd3 7e mov a,m 4bd4 b7 ora а ;zero? 4bd5 c8 rz more to print ; 4bd6 e5 push h 4bd7 4f c,a mov 4bd8 cd6a4b call conout 4bdb el h pop 4bdc 23 inx h 4bdd c3d34b jmp prmsq ; setfunc: set function for next i/o (command in reg-c) ; 4beØ 21684c h,iof ; io function address lxi ;get it to accumulator for maskin 4be3 7e mov a,m 11111000b 4be4 e6f8 ani ; remove previous command 4be6 bl ;set to new command ora C mov 4be7 77 ;replaced in iopb m,a the mds-800 controller reg's disk bank bit in sec ; mask the bit from the current i/o function ; 4be8 e620 ;mask the disk select bit 00100000b ani 4bea 216b4c 1xi ;address the sector selec h,ios 4bed b6 ;select proper disk bank ora m 4bee 77 mov ;set disk select bit on/o m,a 4bef c9 ret ; waitio: 4bfØ ØeØa mvi c, retry ; max retries before perm error rewait: start the i/o function and wait for completion ; 4bf2 cd3f4c call intype ; in rtype 4bf5 cd4c4c call inbyte ; clears the controller ; 4bf8 3a664c lda dbank ;set bank flags

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• • • m-+++

; zero if drive 0,1 and nz 4bfb b7 ora а a.iopb and Øffh ;low address for iopb 4bfc 3e67 mvi 4bfe Ø64c ; high address for iopb mvi b, iopb shr 8 ;drive bank 1? iodrl 4c00 c20b4c jnz 4cØ3 d379 ilow ;low address to controlle out 4cØ5 78 mov a,b 4cØ6 d37a out ihiqh ; high address 4cØ8 c31Ø4c ; to wait for complete jmp waitØ iodrl: ;drive bank 1 4c0b d389 out ilow+10h :88 for drive bank 10 4cØd 78 a,b mov 4c0e d38a ihigh+10h out ; call ;wait for completion 4cl0 cd594c wait0: instat 4cl3 e604 ani iordy ;ready? 4c15 cal04c waitØ jΖ ; check io completion ok ; ; must be io complete (00) 4c18 cd3f4c call intype 00 unlinked i/o complete, Øl linked i/o comple ; 10 disk status changed ll (not used) ; 4clb fe02 10b ;ready status change? cpi 4cld ca324c jz wready ; must be 00 in the accumulator ; 4c20 b7 ora а 4c21 c2384c jnz werror ; some other condition, re ; check i/o error bits ; 4c24 cd4c4c call inbyte 4c27 17 ral 4c28 da324c jc wready ;unit not ready 4c2b 1f rar 4c2c e6fe ani 11111110b ; any other errors? 4c2e c2384c jnz werror ; read or write is ok, accumulator contains zero ; 4c31 c9 ret : wready: ;not ready, treat as error for now 4c32 cd4c4c call inbyte ;clear result byte 4c35 c3384c jmp trycount ; werror: ; return hardware malfunction (crc, track, seek, e the mds controller has returned a bit in each pos ; of the accumulator, corresponding to the conditio ; Ø - deleted data (accepted as ok above) ; - crc error 1 ; 2 - seek error ; 3 - address error (hardware malfunction) ; 4 - data over/under flow (hardware malfunct ; 5 - write protect (treated as not ready) ; 6 - write error (hardware malfunction) ; 7 - not ready ;

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(accumulator bits are numbered 7 6 5 4 3 2 1 Ø) ; ; it may be useful to filter out the various condit ; but we will get a permanent error message if it i ; recoverable. in any case, the not ready conditio ; treated as a separate condition for later improve ; trycount: register c contains retry count, decrement 'til z ; 4c38 Ød dcr C 4c39 c2f24b jnz rewait ; for another try ; cannot recover from error ; 4c3c 3eØl ;error code mvi a.1 4c3e c9 ret ; intype, inbyte, instat read drive bank 00 or 10 4c3f 3a664c intype: lda dbank 4c42 b7 ora a 4c43 c2494c jnz intypl ;skip to bank 10 4c46 db79 in rtype 4c48 c9 ret 4c49 db89 intypl: in rtype+10h ;78 for Ø,1 88 for 2,3 4c4b c9 ret 4c4c 3a664c inbyte: lda dbank 4c4f b7 ora а 4c50 c2564c inz inbytl 4c53 db7b in rbyte 4c55 c9 ret 4c56 db8b inbytl: in rbyte+10h 4c58 c9 ret 4c59 3a664c instat: 1da dbank 4c5c b7 ora a 4c5d c2634c jnz instal 4c60 db78 in dstat 4c62 c9 ret 4c63 db88 instal: in dstat+10h 4c65 c9 ret ; ; ; data areas (must be in ram) : 4c66 00 dbank: db Ø ;disk bank 00 if drive 0.1 10 if drive 2.3 ; iopb: ; io parameter block 4c67 80 db 80h ;normal i/o operation 4c68 Ø4 iof: db readf ; io function, initial read 4c69 Ø1 ion: db 1 ;number of sectors to read 4c6a Ø2 iot: db offset ;track number 4c6b Ø1 ios: db 1 ;sector number 4c6c 8000 iod: dw buff ; io address ; ; ; define ram areas for bdos operation

4c6e+=	begdat	endef equ	Ş			11
4c6e+	dirbuf:	ds	128	;directory	access	buffer
4cee+	alvØ:	ds	31			
4dØd+	csvØ:	ds	16			
4dld+	alvl:	ds	31			
4d3c+	csvl:	ds	16			
4d4c+	alv2:	ds	31			
4d6b+	csv2:	ds	16			
4d7b+	alv3:	ds	31			
4d9a+	csv3:	ds	16			
4daa+=	enddat	equ	\$			
Ø13c+=	datsiz	equ	\$-begdat	1		
4daa		end				

APPENDIX C: A SKELETAL CBIOS

	skeletal cbios for first level of cp/m 2.0 altera
,	Skeletal CDIOS IOI IIISt level of CD/m 2.0 altera
ØØ14 = msize	equ 20 ;cp/m version memory size in kilo
	"bias" is address offset from 3400h for memory sy
;	than 16k (referred to as "b" throughout the text)
;	
0000 = bias	egu (msize-20)*1024
3400 = ccp	egu 3400h+bias ;base of ccp
$3c\emptyset6 = bdos$	equ ccp+806h ;base of bdos
4a00 = bios	equ ccp+1600h ;base of bios
0004 = cdisk	equ ØØØ4h ;current disk number Ø=a,,15=p
0003 = iobyte	equ 0003h ;intel i/o byte
4a00	org bios ;origin of this program
ØØ2c = nsects	equ (\$-ccp)/128 ;warm start sector count
;	
;	jump vector for individual subroutines
4a00 c39c4a	jmp boot ;cold start
4a03 c3a64a wboote:	
4a06 c3114b 4a09 c3244b	jmp const ;console status jmp conin ;console character in
4a0c c3374b	jmp conout ; console character in
4a0f c3494b	jmp list ;list character out
4a12 c34d4b	jmp punch ;punch character out
4al5 c34f4b	jmp reader ;reader character out
4a18 c3544b	jmp home ;move head to home positi
4alb c35a4b	jmp seldsk ;select disk
4ale c37d4b	jmp settrk ;set track number
4a21 c3924b	jmp setsec ;set sector number
4a24 c3ad4b 4a27 c3c34b	jmp setdma ;set dma address
4a27 C3C34b 4a2a c3d64b	jmp read ;read disk jmp write ;write disk
4a2d c34b4b	jmp listst ;return list status
4a30 c3a74b	jmp sectran ;sector translate
;	
;	fixed data tables for four-drive standard
;	ibm-compatible 8" disks
1-22 724-00 2-4	disk parameter header for disk ØØ
4a33 734a00 dpbase:	
4a37 000000 4a3b f04c8d	dw 0000h,0000h dw dirbf,dpblk
4a3f ec4d70	dw dirbf,dpblk dw chk00,all00
;	disk parameter header for disk Øl
4a43 734a00	dw trans,0000h
4a47 000000	dw 0000h,0000h
4a4b fØ4c8d	dw dirbf,dpblk
4a4f fc4d8f	dw chk01,all01
4a53 734a00	disk parameter header for disk Ø2
4a55 754a00 4a57 000000	dw trans,0000h dw 0000h,0000h
4a5b f04c8d	dw dirbf,dpblk
4a5f Øc4eae	dw chk02,all02
	,

; 4a63 734a00 4a67 000000 4a6b f04c8d 4a6f lc4ecd	disk pa dw dw dw dw dw	rameter header fo trans,0000h 0000h,0000h dirbf,dpblk chk03,all03	or disk Ø3
; 4a73 01070d t 4a77 190506 4a75 170309 4a7f 150208 4a83 141a06 4a87 121804 4a8b 1016	sector db db db db db db db db db	translate vector 1,7,13,19 25,5,11,17 23,3,9,15 21,2,8,14 20,26,6,12 18,24,4,10 16,22	sectors 1,2,3,4 sectors 5,6,7,8 sectors 9,10,11,12 sectors 13,14,15,16 sectors 17,18,19,20 sectors 21,22,23,24 sectors 25,26
; d 4a8d 1a00 4a8f 03 4a90 07 4a91 00 4a92 f200 4a94 3f00 4a96 c0 4a97 00 4a98 1000 4a9a 0200	lpblk: ;disk p dw db db db dw dw dw db db dw dw dw	arameter block, o 26 3 7 Ø 242 63 192 Ø 16 2	common to all disks ;sectors per track ;block shift factor ;block mask ;null mask ;disk size-1 ;directory max ;alloc Ø ;alloc 1 ;check size ;track offset
; ; 4a9c af 4a9d 320300 4aa0 320400 4aa3 c3ef4a	individ		to perform each function st perform parameter initi ;zero in the accum ;clear the iobyte ;select disk zero ;initialize and go to cp/
4aa6 318000 4aa9 0e00 4aab cd5a4b 4aae cd544b	lxi mvi call call	sp,80h	ad the disk until all sect ;use space below buffer f ;select disk Ø ;go to track ØØ
4abl 062c 4ab3 0e00 4ab5 1602 ; 4ab7 210034	mvi mvi mvi note th		;b counts # of sectors to ;c has the current track ;d has the next sector to eading track Ø, sector 2 s loader, which is skipped ;base of cp/m (initial lo
		ne more sector b ;save so d ;save no h ;save du c,d ;get se setsec ;set se	ector count, current track ext sector to read ma address ctor address to register c ctor address from register dma address to b,c

;replace on stack for later recal 4ac2 c5 push b 4ac3 cdad4b call setdma ;set dma address from b,c : drive set to Ø, track set, sector set, dma addres ; 4ac6 cdc34b call read ØØh 4ac9 fe00 cpi ; any errors? ; retry the entire boot if an erro 4acb c2a64a jnz wboot ; no error, move to next sector ; ;recall dma address 4ace el pop h d,128 4acf 118000 lxi :dma=dma+128 4ad2 19 ;new dma address is in h,l dad d 4ad3 d1 pop d ;recall sector address 4ad4 cl ;recall number of sectors remaini b pop 4ad5 Ø5 dcr b ;sectors=sectors-1 4ad6 caef4a ;transfer to cp/m if all have bee jz gocpm ; more sectors remain to load, check for track chan ; 4ad9 14 inr d 4ada 7a a,d ;sector=27?, if so, change tracks mov 4adb felb 27 cpi 4add daba4a jc loadl ;carry generated if sector<27 ; end of current track, go to next track ; 4aeØ 16Ø1 mvi ; begin with first sector of next d,1 4ae2 Øc inr C ;track=track+1 ; save register state, and change tracks ; 4ae3 c5 push b d 4ae4 d5 push 4ae5 e5 push h 4ae6 cd7d4b ;track address set from register call settrk 4ae9 el h pop d 4aea d1 pop 4aeb cl pop b 4aec c3ba4a loadl : for another sector jmp ; end of load operation, set parameters and go to c ; gocpm: 4aef 3ec3 a,Øc3h ;c3 is a jmp instruction mvi 4afl 320000 sta Ø ; for jmp to wboot 4af4 21034a lxi h.wboote ;wboot entry point 4af7 220100 shld 1 ;set address field for jmp at Ø ; 4afa 320500 5 sta ; for jmp to bdos 4afd 21063c lxi h,bdos ;bdos entry point 4b00 220600 shld 6 ; address field of jump at 5 to bd ; 4b03 018000 lxi b,80h ;default dma address is 80h 4b06 cdad4b call setdma ; 4b09 fb ei ;enable the interrupt system 4b0a 3a0400 1da cdisk ;get current disk number 4bød 4f mov c,a ;send to the ccp 4b0e c30034 jmp CCP ;go to cp/m for further processin

; ; simple i/o handlers (must be filled in by user) ; in each case, the entry point is provided, with s ; to insert your own code ; ; ; console status, return Øffh if character ready, const: 4b11 ;space for status subroutine ds lØh 4b21 3e00 mvi a.00h 4b23 c9 ret ; ; console character into register a conin: 4b24 lØh ;space for input routine ds 4b34 e67f 7fh strip parity bit ani 4b36 c9 ret ; conout: ; console character output from register c 4b37 79 mov a,c ;get to accumulator 4b38 ds lØh ;space for output routine 4b48 c9 ret list: ;list character from register c 4b49 79 mov a,c ; character to register a 4b4a c9 ret ;null subroutine listst: ;return list status (Ø if not ready, l if ready) 4b4b af ;Ø is always ok to return xra a 4b4c c9 ret ; ; punch character from register c punch: 4b4d 79 ; character to register a mov a,c 4b4e c9 ;null subroutine ret ; ; reader: ; read character into register a from reader devic 4b4f 3ela mvi a,lah ;enter end of file for now (repla 4b51 e67f ani 7fh ; remember to strip parity bit 4b53 c9 ret ; ; ; i/o drivers for the disk follow for now, we will simply store the parameters away ; in the read and write subroutines ; ; ;move to the track 00 position of current drive home: translate this call into a settrk call with param ; 4b54 ØeØØ mvi C.Ø ;select track Ø 4b56 cd7d4b call settrk 4b59 c9 ret ;we will move to 00 on first read ; seldsk: ;select disk given by register c 4b5a 210000 lxi h,0000h ;error return code 4b5d 79 mov a,c 4b5e 32ef4c sta diskno 4b61 feØ4 cpi ;must be between Ø and 3 4

;no carry if 4,5,... 4b63 dØ rnc disk number is in the proper range ; ;space for disk select 4b64 ds 10 compute proper disk parameter header address ; 4b6e 3aef4c lda diskno ;l=disk number 0,1,2,3 4b71 6f mov 1,a 4b72 2600 ; high order zero h,Ø mvi ;*2 4b74 29 dad h :*4 h 4b75 29 dad 4b76 29 dad h :*8 4b77 29 ;*16 (size of each header) dad h d,dpbase 4b78 11334a lxi ;hl=.dpbase(diskno*16) 4b7b 19 dad d 4b7c c9 ret ; settrk: ;set track given by register c 4b7d 79 a.c mov sta track 4b7e 32e94c lØh ;space for track select 4b81 ds 4b91 c9 ret ; setsec: ;set sector given by register c 4b92 79 mov a,c 4b93 32eb4c sta sector 4b96 lØh ;space for sector select ds 4ba6 c9 ret ; sectran: ;translate the sector given by bc using the ;translate table given by de 4ba7 eb xchq ;hl=.trans 4ba8 Ø9 dad b ; hl=.trans(sector) 4ba9 6e 1,m ;1 = trans(sector) mov 4baa 2600 ;hl= trans(sector) mvi h,Ø 4bac c9 ret ;with value in hl ; setdma: ;set dma address given by registers b and c 4bad 69 mov 1,c ;low order address 4bae 60 mov h,b ; high order address 4baf 22ed4c shld dmaad ;save the address 4bb2 ds 10h ;space for setting the dma addres 4bc2 c9 ret ; read: ;perform read operation (usually this is similar so we will allow space to set up read command, th ; common code in write) ; 4bc3 ds lØh ;set up read command 4bd3 c3e64b jmp waitio ; to perform the actual i/o ;perform a write operation write: 4bd6 lØh ;set up write commanu ds : waitio: ;enter here from read and write to perform the ac operation. return a 00h in register a if the ope ; properly, and Ølh if an error occurs during the r ;

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4be6 4ce6 4ce8		;;;;;	ds mvi ret	256 a,1	<pre>have saved the disk number in 'd the track number in 'track' (0-76 the sector number in 'sector' (1- the dma address in 'dmaad' (0-655 ;space reserved for i/o drivers ;error condition ;replaced when filled-in</pre>
		2			the cbios is reserved uninitiali loes not need to be a part of the
					nage (the space must be available,
		;			n "begdat" and "enddat").
		;			
4ce9		track:	ds	2	;two bytes for expansion
4ceb		sector:	ds	2 2	;two bytes for expansion
4ced		dmaad:	ds	2	;direct memory address
4cef		diskno:	ds	1	;disk number 0-15
		;			
		1			a for bdos use
4cfØ	=	begdat	egu	\$;beginning of data area
4cfØ		dirbf:	ds	128	;scratch directory area
4d7Ø		al100:	ds	31	;allocation vector Ø
4d8f		allØl:	ds	31	;allocation vector 1
4dae		al102:	ds	31	;allocation vector 2
4dcd		al103:	ds	31	;allocation vector 3
4dec		chk00:	ds	16	;check vector Ø
4dfc		chkØ1:	ds	16	; check vector 1
4eØc 4elc		chk02:	ds	16	; check vector 2
4erc		chkØ3:	ds	16	;check vector 3
4e2c	=	enddat	equ	Ş	;end of data area
Ø13c		datsiz	equ		t;size of data area
4e2c			end	,, u u	
			2.4993.000 		

1	APPENDIX	D: A SK	ELETAL G	ETSYS/PUTSYS P	ROGRAM
	;;			and putsys pro ms at the base	grams from Sec 4. of the TPA
0100		org	0100h		
0014 =	msize	egu	20	; size	of cp/m in Kbytes
	; "bias" ;			o add to addre "b" throughout	sses for > 20k the text)
0000 = 3400 = 3c00 = 4a00 = 4a000 0 = 4a0000000000	bias ccp bdos bios	equ equ	(msize-2 3400h+bi ccp+0800 ccp+1600	as Ih	
	; ;	getsys p 3880h +		tracks Ø and l	to memory at
	;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;	register a b c d,e h,l sp		usage (scratch regis track count (Ø sector count ((scratch regis load address set to stack a	76) 126) ter pair)
Ø100 318033 Ø103 218033 Ø106 Ø600	gstart: rd\$trk:	lxi mvi	sp,ccp-0 h,ccp-00 b,0		; start of getsys ; convenient plac ; set initial loa ; start with trac ; read next track
Ø108 Øe01	rd\$sec:	mvi	c,1		; each track star
010a cd0003 010d 118000 0110 19 0111 0c 0112 79 0113 felb 0115 da0a01	call lxi dad inr mov cpi jc	lxi dad inr mov cpi	read\$sec d,128 d c a,c 27 rdsec	2	<pre>; get the next se ; offset by one s ; (h1=h1+128) ; next sector ; fetch sector nu ; and see if la ; <, do one more</pre>
	; arriv	e here at	end of	track, move to	next track
0118 04 0119 78 011a fe02 011c da0801		mov cpi	b a,b 2 rd\$trk		<pre>; track = track+1 ; check for last ; track = 2 ? ; <, do another</pre>
	; arriv	e here at	end of	load, halt for	lack of anything b
Øllf fb Øl20 76		ei hlt			

-

U 1<u>2/212</u> 01 0

putsys program, places memory image starting at 3880h + bias back to tracks 0 and 1 start this program at the next page boundary

0200 org (\$+0100h) and 0ff00h

;;;

0200 318033 0203 218033 0206 0600		sp,ccp-0080h h,ccp-0080h b,0	; convenient plac ; start of dump ; start with trac		
Ø208 Øe01	mvi wr\$sec:	c,1	; start with sect		
020a cd0004 020d 118000 0210 19 0211 0c 0212 79 0213 felb 0215 da0a02	call lxi dad inr mov cpi	write\$sec d,128 d c a,c 27 wr\$sec	<pre>; write one secto ; length of each ; <hl>=<hl> + 128 ; <c> = <c> + 1 ; see if ; past end of t ; no, do another</c></c></hl></hl></pre>		
	; arrive here at end of track, move to next track				
0218 04 0219 78 021a fe02 021c da0802	inr mov cpi jc	b a,b 2 wr\$trk	; track = track+1 ; see if ; last track ; no, do another		
	; done w:	ith putsys, halt for lack	of anything bette		
Ø21f fb Ø22Ø 76	ei hlt				
	; user supplied	d subroutines for sector	read and write		
	; move to	o next page boundary			
0300	org	(\$+0100h) and 0ff00h			
	; track ; secto	the next sector (in <b), or in <c> ddr in <hl></hl></c></b), 			
0300 c5 0301 e5	push push	b h			
0302	; user defined ds	read operation goes here 64			
Ø342 el Ø343 cl	pop	h b			

Ø344 c9	ret
0400	org (\$+0100h) and 0ff00h ; another page bo
	write\$sec:
	; same parameters as read\$sec
Ø400 c5	push b
Ø401 e5	push h
0402	; user defined write operation goes here ds 64
Ø442 el	pop h
Ø443 cl	pop b
Ø444 c9	ret
	; end of getsys/putsys program
0445	end

; this is a sample cold start loader which, when modified ; resides on track 00, sector 01 (the first sector on the ; diskette). we assume that the controller has loaded : this sector into memory upon system start-up (this pro-; gram can be keyed-in, or can exist in read/only memory ; beyond the address space of the cp/m version you are ; running). the cold start loader brings the cp/m system ; into memory at "loadp" (3400h + "bias"). in a 20k ; memory system, the value of "bias" is 0000h, with large ; values for increased memory sizes (see section 2). afte ; loading the cp/m system, the clod start loader branches ; to the "boot" entry point of the bios, which begins at ; "bios" + "bias." the cold start loader is not used un-; til the system is powered up again, as long as the bios ; is not overwritten. the origin is assumed at 0000h, an ; must be changed if the controller brings the cold start ; loader into another area, or if a read/only memory area ; is used.

0000		org	Ø	;	base of ram in cp/m
ØØ14 = n	msize	equ	20	;	min mem size in kbytes
3400 = 0 4a00 = 0 0300 = 0 4a00 = 0 1900 = 0	bias ccp bios biosl boot size sects	egu egu egu egu egu egu	(msize-20)*1024 3400h+bias ccp+1600h 0300h bios bios+biosl-ccp size/128	;;;;;;	offset from 20k system base of the ccp base of the bios length of the bios size of cp/m system # of sectors to load
;	;	begin th	ne load operatior	ì	
0000 010200 0003 1632 0005 210034	cold:	lxi mvi lxi	b,2 d,sects h,ccp	;	<pre>b=Ø, c=sector 2 d=# sectors to load base transfer address</pre>
:	lsect:	; load t	the next sector		
		incort i	pline code at th		noint to

; insert inline code at this point to ; read one 128 byte sector from the ; track given in register b, sector ; given in register c, ; into the address given by <hl> ; ; branch to location "cold" if a read error occurs

; * ; * user supplied read operation goes here ... ; * ; ; 0008 c36b00 past\$patch ; remove this when patche jmp 000b ds 60h past\$patch: ; go to next sector if load is incomplete ØØ6b 15 dcr d ; sects=sects-1 006c ca004a iΖ boot ; head for the bios ; more sectors to load ; we aren't using a stack, so use <sp> as scratch registe ; to hold the load address increment ; 006f 318000 1 x i sp,128 ; 128 bytes per sector 0072 39 dad ; $\langle hl \rangle = \langle hl \rangle + 128$ sp 0073 Øc inr C ; sector = sector + 1 0074 79 mov a,c 0075 felb ; last sector of track? 27 cpi 0077 da0800 jc lsect ; no, go read another ; end of track, increment to next track 007a 0e01 c,1 ; sector = 1mvi Ø07c Ø4 inr b ; track = track + 1 007d c30800 ; for another group jmp lsect 0080 end ; of boot loader

APPENDIX F: CP/M DISK DEFINITION LIBRARY 1: ; CP/M 2.0 disk re-definition library 2: ; 3: : Copyright (c) 1979 4: ; Digital Research 5: ; Box 579 6: ; Pacific Grove, CA 7: ; 93950 8: ; CP/M logical disk drives are defined using the 9: ; 10: ; macros given below, where the sequence of calls 11: ; is: 12: ; 13: ; disks 11 14: ; diskdef parameter-list-Ø 15: ; diskdef parameter-list-l 16: : 17: : diskdef parameter-list-n 18: : endef 19: ; 20: ; where n is the number of logical disk drives attached to the CP/M system, and parameter-list-i defines the 21: ; 22: ; characteristics of the ith drive (i=0,1,...,n-1) 23: ; 24: ; each parameter-list-i takes the form 25: : dn,fsc,lsc,[skf],bls,dks,dir,cks,ofs,[0] 26: ; where is the disk number Ø,1,...,n-1 27: ; dn is the first sector number (usually Ø or 1) 28: ; fsc 29: ; lsc is the last sector number on a track is optional "skew factor" for sector translate 30: : skf 31: ; bls is the data block size (1024,2048,...,16384) 32: ; dks is the disk size in bls increments (word) 33: ; dir is the number of directory elements (word) 34: : is the number of dir elements to checksum cks is the number of tracks to skip (word) 35:; ofs is an optional Ø which forces 16K/directory en 36: ; $[\emptyset]$ 37: ; 38: ; for convenience, the form 39: ; dn.dm 40: ; defines disk dn as having the same characteristics as 41: ; a previously defined disk dm. 42: : 43: : a standard four drive CP/M system is defined by 44: ; disks 45: ; diskdef 0,1,26,6,1024,243,64,64,2 46: ; dsk set Ø 47: ; 3 rept 48: ; dsk dsk+1 set 49: ; diskdef %dsk,Ø 50: ; endm 51: ; endef 52: ; the value of "begdat" at the end of assembly defines t 53: :

54: ; beginning of the uninitialize ram area above the bios, 55: ; while the value of "enddat" defines the next location 56: ; following the end of the data area. the size of this 57: ; area is given by the value of "datsiz" at the end of t 58: ; assembly. note that the allocation vector will be qui 59: ; large if a large disk size is defined with a small blo 60: : size. 61: ; 62: dskhdr macro dn 63: ;; define a single disk header list 64: dpe&dn: dw xlt&dn.0000h ;translate table 65: 0000h.0000h dw ;scratch area 66: dw dirbuf,dpb&dn ;dir buff,parm block 67: dw csv&dn,alv&dn ;check, alloc vectors 68: endm 69: ; 70: disks macro nd define nd disks 71: ;; 72: ndisks ;; for later reference set nd 73: dpbase S equ ; base of disk parameter blocks 74: ;; generate the nd elements 75: dsknxt set Ø 76: nd rept 77: dskhdr %dsknxt 78: dsknxt dsknxc+1 set 79: endm 80: endm 81: ; 82: dpbhdr macro dn 83: dpb&dn equ S ;disk parm block 84: endm 85: ; 86: ddb macro data, comment 87: ;; define a db statement 88: db data comment 89: endm 90: ; 91: ddw macro data, comment 92: ;; define a dw statement 93: dw data comment 94: endm 95: ; 96: gcd macro m,n 97: ;; greatest common divisor of m,n 98: ;; produces value gcdn as result 99: ;; (used in sector translate table generation) 100: gcdm set m ;;variable for m 101: gcdn set n ;;variable for n 102: gcdr set Ø ;;variable for r 103: 65535 rept 104: gcdx set gcdm/gcdn 105: gcdr gcdm - gcdx*gcdn set 106: if $gcdr = \emptyset$ 107: exitm 108: endif

109: gcdm set qcdn 110: gcdn set gcdr 111: endm 112: endm 113: ; 114: diskdef macro dn,fsc,lsc,skf,bls,dks;dir,cks,bfs,kl6 115: ;; generate the set statements for later tables 116: if nul lsc 117: ;; current disk dn same as previous fsc 118: dpb&dn dpb&fsc ;equivalent parameters eau 119: als&dn equ als&fsc ;same allocation vector size 120: css&dn equ css&fsc ;same checksum vector size 121: xlt&dn equ xlt&fsc ;same translate table 122: else 123: secmax lsc-(fsc) ;;sectors Ø...secmax set secmax+l;;number of sectors 124: sectors set 125: als&dn set (dks)/8 ;;size of allocation vector 126: if ((dks) mod δ) ne Ø 127: als&dn set als&dn+1 endif 128: 129: css&dn set (cks)/4 ;;number of checksum elements 130: ;; generate the block shift value 131: blkval set bls/128 ;;number of sectors/block ;;counts right 0's in blkval 132: blkshf set Ø 133: blkmsk set Ø ;;rills with 1's from right ;;once for each bit position 134: rept 16 135: if blkval=1 136: exitm 137: endif 138: ;; otherwise, high order 1 not found yet 139: blkshf set blkshf+1 140: blkmsk set (blkmsk shl 1) or 1 141: blkval set blkval/2 142: endm 143: ;; generate the extent mask byte 144: blkval set bls/1024 ;;number of kilobytes/block 145: extmsk set ;;fill from right with l's 0 146: rept 16 147: if blkval=1 148: exitm 149: endif 150: ;; otherwise more to shift 151: extmsk set (extmsk shl 1) or 1 152: blkval set blkval/2 153: endm 154: ;; may be double byte allocation 155: (dks) > 256 if 156: extmsk set (extmsk shr 1) 157: endif 158: ;; may be optional [0] in last position 159: if not nul k16 160: extmsk set k16 161: endif now generate directory reservation bit vector 162: ;; 163: dirrem set dir ;;# remaining to process

bls/32 ;;number of entries per block 164: dirbks set ;;fill with l's on each loop 165: dirblk Ø set 16 166: rept 167: if dirrem=Ø 168: exitm 169: endif not complete, iterate once again 170: ;; shift right and add 1 high order bit 171: ;; 172: dirblk (dirblk shr i) or 8000h set 173: if dirrem > dirbks 174: dirrem set dirrem-dirbks 175: else Ø 176: dirrem set endif 177: 178: endm 179: dpbhdr dn ;;generate equ \$ 180: ddw %sectors,<;sec per track> 181: ddb %blkshf,<;blcck shift> %blkmsk,<;block mask> 182: ddb 183: ddb %extmsk,<;extnt mask> 184: ddw %(dks)-l,<;uisk size-l> 185: udw %(dir)-l,<;airectory max> 186: dàb %dirblk shr 8,<;alloc0> 187: ddb %dirblk and Øffh,<;allocl> %(cks)/4,<;check size> 188: ddw 189: ddw %ofs,<;offset> 190: ;; generate the translate table, if requested 191: if nul skf 192: xlt&dn equ Ø ;no xlate table 193: else 194: $skf = \emptyset$ if 195: xlt&dn equ Ø ;no xlate table 196: else 197: ;; generate the translate table 198: nxtsec Ø ;;next sector to fill set 199: nxtbas Ø ;; moves by one on overflow set 200: gcd %sectors,skf 201: ;; gcdn = gcd(sectors,skew) 202: neltst set sectors/gcdn 203: ;; neltst is number of elements to generate 204: ;; before we overlap previous elements 205: nelts set neltst ;;counter 206: xlt&dn \$ equ ;translate table 207: rept sectors ;;once for each sector 208: if sectors < 256 209: ddb %nxtsec+(fsc) 210: else 211: ddw %nxtsec+(fsc) 212: endif 213: nxtsec set nxtsec+(skf) 214: if nxtsec >= sectors 215: nxtsec set nxtsec-sectors 216: endif 217: nelts nelts-1 set 218: if nelts = Ø

22Ø: 221: 222:	nxtbas nxtsec nelts	set set endif	nxtbas+i nxtbas neltst
223: 224: 225: 226:		endm endif endif endm	;;end of nul fac test ;;end of nul bls test
227: 228: 229: 230: 231:	defds lab:	macro ds endm	lab,space space
232: 233: 234:	lds	macro äefds endm	lb,dn,val lb&dn,%val&dn
	endef	macro	
237: 238:		generate egu	e the necessary ram data areas \$
	dirbuf: dsknxt		<pre>128 ;directory access buffer Ø</pre>
241: 242: 243:	donnac	rept lås lås	ndisks ;;once for each disk alv,%dsknxt,als csv,%dsknxt,css
244: 245:	dsknxt	set endm	dsknxt+1
246:	enddat datsiz ;;	egu egu	\$ Ş-begdat this point forces hex record

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 $\mathcal{A} \in \mathcal{A}$

1: :***** 2: :* 3: ;* Sector Deblocking Algorithms for CP/M 2.0 + 4: ;* 6: : utility macro to compute sector mask 7: ; 8: smask hblk macro compute log2(hblk), return @x as result 9: ;; 10: ;; (2 ** @x = hblk on return)11: @y hblk set 12: @x set 13: ;; count right shifts of @y until = 1 14: rept 8 15: if @y = 116: exitm 17: endif 18: ;; @y is not 1, shift right one position 19: @y set @y shr l 20: @x 0x + 1set 21: endm 22: endm 23: ; 25: ;* * 26: :* CP/M to host disk constants * 27: ;* 29: blksiz 2048 :CP/M allocation size equ 30: hstsiz equ 512 ;host disk sector size 31: hstspt equ 20 ;host disk sectors/trk 32: hstblk hstsiz/128 ;CP/M sects/host buff equ hstblk * hstspt ;CP/M sectors/track 33: cpmspt equ 34: secmsk hstblk-1 equ ;sector mask 35: smask hstblk ;compute sector mask 36: secshf equ 0x ;log2(hstblk) 37: ; 39: ;* * 40: ;* * BDOS constants on entry to write 41: ;* Ø 43: wrall equ ;write to allocated 44: wrdir 1 equ ;write to directory 45: wrual equ 2 ;write to unallocated 46: ; 48: ;* 49: ;* The BDOS entry points given below show the * 50: ;* code which is relevant to deblocking only. 51: ;* 53: ;

54: ; DISKDEF macro, or hand coded tables go here ;disk param block base 55: dpbase equ \$ 56: ; 57: boot: 58: wboot: 59: ;enter here on system boot to initialize 60: xra a ;Ø to accumulator hstact 61: ;host buffer inactive sta 62: unacnt ;clear unalloc count sta 63: ret 64: ; 65: seldsk: 66: ;select disk 67: mov ;selected disk number a,c 68: sta sekdsk ;seek disk number 69: ;disk number to HL mov 1,a 70: mvi h,Ø 71: ; multiply by 16 rept 4 72: dad h 73: endm 74: lxi d,dpbase ;base of parm block 75: dad d ;hl=.dpb(curdsk) 76: ret 77:; 78: settrk: 79: ;set track given by registers BC 80: mov h,b 81: mov 1,c 82: shld sektrk ;track to seek 83: ret 84: ; 85: setsec: 86: ;set sector given by register c 87: mov a,c 88: sta seksec ;sector to seek 89: ret 90: ; 91: setdma: 92: ;set dma address given by BC 93: mov h,b 94: mov 1,c 95: dmaadr shld 96: ret 97: ; 98: sectran: 99: ;translate sector number BC 100: mov h,b 101: mov 1,c 102: ret 103: ;

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* 105: :* * The READ entry point takes the place of 106: :* * 107: ;* the previous BIOS definition for READ. * 108: :* 110: read: 111: :read the selected CP/M sector 112: mvi a.1 113: sta readop ;read operation 114: sta rsflag ;must read data 115: mvi a, wrual 116: sta wrtype ;treat as unalloc 117: jmp ; to perform the read rwoper 118: ; 120: ;* * 121: ;* * The WRITE entry point takes the place of 122: :* * the previous BIOS defintion for WRITE. + 123: ;* 125: write: 126: ;write the selected CP/M sector 127: xra ;Ø to accumulator a 128: sta readop ;not a read operation 129: mov ;write type in c a,c 130: sta wrtype 131: cpi wrual ;write unallocated? 132: chkuna :check for unalloc jnz 133: ; 134: ; write to unallocated, set parameters 135: a,blksiz/128 mvi ;next unalloc recs 136: sta unacnt 137: lda sekdsk ;disk to seek 138: sta unadsk ;unadsk = sekdsk 139: lhld sektrk 140: shld unatrk ;unatrk = sectrk 141: lda seksec 142: sta unasec ;unasec = seksec 143: ; 144: chkuna: 145: ; check for write to unallocated sector 146: lda unacnt ; any unalloc remain? 147: ora a 148: alloc jz ;skip if not 149: ; 150: : more unallocated records remain 151: dcr a ;unacnt = unacnt-1 152: sta unacnt 153: 1da sekdsk ;same disk? 154: lxi h, unadsk 155: cmp m ;sekdsk = unadsk? 156: jnz alloc ;skip if not 157: ; 158: ; disks are the same

159: lxi h,unatrk ;sektrk = unatrk? 160: call sektrkcmp 161: ;skip if not jnz alloc 162: ; 163: ; tracks are the same 164: lda seksec ;same sector? 165: lxi h,unasec ;seksec = unasec? 166: Cmp m 167: alloc ;skip if not jnz 168: ; match, move to next sector for future ref 169: ; 170: inr ;unasec = unasec+1 m 171: mov a,m ;end of track? 172: cpi cpmspt ;count CP/M sectors 173: jc noovf ;skip if no overflow 174: ; 175: ; overflow to next track 176: mvi m,Ø ; unasec = \emptyset 177: lhld unatrk 178: inx h 179: shld unatrk ;unatrk = unatrk+1 180: ; 181: noovf: 182: ;match found, mark as unnecessary read ;Ø to accumulator 183: а xra 184: rsflag ;rsflag = Ø sta 185: jmp rwoper ; to perform the write 186: ; 187: alloc: ;not an unallocated record, requires pre-read 188: 189: ;0 to accum xra a 190: sta unacnt ; unacht = \emptyset 191: inr ;1 to accum a 192: rsflag sta ;rsflag = 1 193: ; * 195: ;* * 196: ;* Common code for READ and WRITE follows 197: ;* 199: rwoper: 200: ;enter here to perform the read/write 201: xra ;zero to accum a 202: sta erflag ;no errors (yet) 203: lda seksec ; compute host sector 204: rept secshf 205: ora a ;carry = Ø 206: rar ;shift right 207: endm 208: sta sekhst ;host sector to seek 209: ; 210: ; active host sector? 211: lxi h,hstact ;host active flag 212: mov a,m 213: mvi m,1 ;always becomes 1

a filhst 214: ora ;was it already? 215: :fill host if not jz 216: ; 217: ; host buffer active, same as seek buffer? 218: lda sekdsk lxi cmp jnz 219: h,hstdsk ;same disk? 220: ;sekdsk = hstdsk? m 221: nomatch 222: ; 223: ; same disk, same track? 224: lxi h,hsttrk 225: call sektrkcmp ;sektrk = hsttrk? 226: jnz nomatch 227: ; same disk, same track, same buffer? 228: ; 229: lda sekhst 230: lxi h,hstsec ;sekhst = hstsec? 231: cmp m 232: jz match ;skip if match 233: ; 234: nomatch: 235: ;proper disk, but not correct sector 236: lda hstwrt ;host written? 237: ora а 238: cnz writehst :clear host buff 239: ; 240: filhst: 241: ;may have to fill the host buffer lda sekdsk 242: sta 243: sta hstdsk lhld sektrk shld hsttrk lda sekhst sta hstsec lda rsflag ora a hstdsk 244: 245: 246: 247: 248: ;need to read? ora a cnz readhst xra a sta hstwrt 249: 250: ;yes, if l 251: ;0 to accum 252: ;no pending write 253: ; 254: match: 255: ; copy data to or from buffer 256: lda seksec ;mask buffer number 257: ani secmsk ;least signif bits 258: mov ;ready to shift l,a mvi h,Ø 259: ;double count 260: rept dad 7 ;shift left 7 261: h 262: endm 263: ; hl has relative host buffer address 264: lxi d,hstbuf 265: dad d ;hl = host address 266: xchq ;now in DE lhld 267: dmaadr ;get/put CP/M data 268: mvi c,128 ;length of move

;which way? 269: lda readop 270: ora a ;skip if read 271: jnz rwmove 272: ; 273: ; write operation, mark and switch direction 274: a,1 mvi 275: sta hstwrt ; hstwrt = 1276: ;source/dest swap xchq 277: ; 278: rwmove: ;C initially 128, DE is source, HL is dest 279: 280: ;source character ldax d 281: inx d 282: mov m,a ;to dest 283: inx h 284: ;loop 128 times dcr C 285: jnz rwmove 286: ; data has been moved to/from host buffer 287: ; 288: lda ;write type wrtype 289: wrdir cpi ;to directory? 290: lda erflag ; in case of errors 291: rnz ;no further processing 292: ; 293: ; clear host buffer for directory write 294: ora ;errors? a 295: ;skip if so rnz 296: xra ;0 to accum a 297: hstwrt ; buffer written sta 298: call writehst 299: lda erflag 300: ret 301: ; * 303: ;* * 304: ;* Utility subroutine for 16-bit compare 305: ;* * 307: sektrkcmp: 308: ;HL = .unatrk or .hsttrk, compare with sektrk 309: xchq 310: lxi h.sektrk 311: ldax d ;low byte compare 312: CMp m ;same? 313: rnz ;return if not 314: ; low bytes equal, test high 1s 315: inx d 316: inx h 317: ldax d 318: cmp m ;sets flags 319: ret 320: ;

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322: ;* * * 323: :* WRITEHST performs the physical write to * 324: :* the host disk, READHST reads the physical * 325: ;* disk. 326: ;* * 328: writehst: 329: ;hstdsk = host disk #, hsttrk = host track #, 330: ;hstsec = host sect #. write "hstsiz" bytes 331: ; from hstbuf and return error flag in erflag. ;return erflag non-zero if error 332: 333: ret 334: ; 335: readhst: 336: ;hstdsk = host disk #, hsttrk = host track #, 337: ;hstsec = host sect #. read "hstsiz" bytes 338: ; into hstbuf and return error flag in erflag. 339: ret 340: ; * 342: :* * 343: ;* Unitialized RAM data areas 344: ;* * 346: ; 347: sekdsk: ds 1 ;seek disk number 348: sektrk: ds 2 ;seek track number 349: seksec: ds 1 ;seek sector number 350: ; 351: hstdsk: ds 1 ;host disk number 352: hsttrk: ds 2 ;host track number 353: hstsec: ds 1 ;host sector number 354: ; 355: sekhst: ds 1 ;seek shr secshf 356: hstact: ds 1 ;host active flag 357: hstwrt: ds 1 ;host written flag 358: ; 359: unacht: ds 1 ;unalloc rec cnt 360: unadsk: ds 1 ;last unalloc disk 2 361: unatrk: ds ;last unalloc track 362: unasec: ds 1 ;last unalloc sector 363: ; 364: erflag: ds 1 ;error reporting 365: rsflag: ds 1 ;read sector flag 366: readop: ds 1 ;1 if read operation ;write operation type 367: wrtype: ds 1 ;last dma address 368: dmaadr: ds 2 369: hstbuf: ds hstsiz ;host buffer 370: ;



