User's Manual DISK JOCKEYTM I CONTROLLER

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User's Manual
DISK JOCKEYTM I

INTRODUCTION

The Morrow's Micro-Stuff DISK JOCKEY I (DJ) board features two distinct subsections:

- 1. A universal floppy disk controller capable of interfacing to a wide variety of floppy disk drives,
- 2. A serial interface that allows communication with a terminal device at TTY 20ma current loop or RS-232 levels.

The floppy disk controller section will interface with any floppy disk drive plug compatible with the Shugart 800 drive. Siemens, Remex, Memorex and MFE are some of the manufacturers of full-sized floppy disks which are plug compatible with the Shugart 800/801.

The DJ plugs into an S-100 bus slot in a system with an 8080 type CPU and has a cable connector for a flat cable to the first floppy disk drive. The controller can control a chain of up to four drives daisy chained on this cable. A second connector on the DJ is provided for the attachment of a terminal device.

The DJ uses memory mapped I/O. Device registers used to input from and output to the floppy disk and the serial port are accessed from the CPU board of the S-100 system by references to memory addresses. Some registers differ in function depending on whether they are being read from or written to.

Most users will not wish to use the hardware level registers directly. Instead, they can call standard disk and serial I/O subroutines contained in 512 bytes of PROM memory on the DJ board. This PROM occupies a 512 byte block of S-100 bus memory addresses. A 256 byte RAM is also provided which is used by the PROM firmware as a data buffer and for temporary computation.

Introduction

The actual addresses where the I/O registers, PROM, and RAM appear is controlled by another PROM, referred to as the address selection PROM. This PROM is supplied with standard addresses burned into it for these registers. If the standard addresses would conflict with some other device on the system, a PROM burned with non-standard addresses can be substituted.

The Disk Jockey uses 1024 bytes of memory starting at 340:0000 or E000H (standard version). The first 512 bytes are occupied by ROM, the next 256 bytes contain a RAM buffer. The last 256 bytes constitute the memory mapped I/O.

PROGRAMMING SPECIFICATION

Most users will wish to take advantage of the standard I/O subroutines supplied in PROM on the DJ.

The user should branch to the appropriate address in a jump table in the first few words of the system ROM. Since the subroutines end with a RET instruction, a CALL instruction should be used to branch to the subroutine.

The jump table contains jump instructions to the true address of the utility routines within the ROM. Having a jump table allows the individual routines to be updated and moved around within the ROM without having to change software that calls the routines. Let A represent the address of word Ø of the onboard ROM. In boards with standard address decoding PROMS, A = 340:000. The addresses to call for the utility routines are then:

Address	Standard Value	Symbolic Value	Function
A	340:000	DBOOT	DOS bootstrap routine Serial input Serial output Recalibrate (Seek to TRKØØ) Seek Select sector Set DMA address
A+3	340:003	TERMIN	
A+6	340:006	TRMOUT	
A+9	340:011	TKZERO	
A+12	340:014	TRKSET	
A+15	340:017	SECTOR	
A+18	340:022	DMA	
A+21	340:025	READ	Read a sector of disk data
A+24	340:030	WRITE	Write a sector of disk data

The specific function of each subroutine is described below.

A subroutine upon completion will execute a RET instruction. A disk subroutine that completes normally will return with a zero in the accumulator. A disk subroutine that detects an error condition will return a nonzero value in the accumulator. A program should execute an ANA A instruction after return from a disk utility subroutine to set the flags according to the contents of the accumulator and then branch if non-zero to an appropriate error handling routine.

Serial I/O

At the hardware level, a means is provided for the CPU program to determine whether the present level of the serial input is a MARK (1) or a SPACE (\emptyset). The user has a choice of connecting to the TTY or RS-232 serial input pin depending on the electrical requirements of the input device.

When a program in the CPU looks at bit 1 of the Read Status Register (see Hardware Registers below), this bit will be a \emptyset if the input line is at the SPACE level and a 1 if the input line is at the MARK level.

The hardware only provides the facility for the software to determine whether the device attached to the serial in line is sending a Ø or 1 signal. Decoding a pattern in time of zeroes and ones into meaningful information is the province of software.

Built-in Disk Jockey firmware provides the facility to receive or send a byte of data from the input device according to industry wide conventions for serial data transmission. These conventions are explained below.

Normally, the interface line will be at a 1 level. When an input device wishes to send a byte of data, it puts the input line to a \emptyset level for a length of time referred to as one bit time. This is called the start bit. At the conclusion of the start bit, the input device sends in order starting with the least significant bit and ending with the most significant bit, the eight data bits of the byte. To do this, the input device puts the interface line at a \emptyset or a 1 level for one bit time, depending on the value of the data bit to be sent. At the conclusion of the bit time, the next data bit is sent. After the final bit is sent, the line returns to the 1 level for at least one bit time (the stop bit).

After the final stop bit is sent, the input device is free to send the next character. If the next character is not yet available, for example, the next key on the input keyboard has not been pressed, the input device can wait any amount of time before lowering the line to a Ø to signal the beginning of the next character.

The format described above is the most common format used by terminal devices. There is some variation, however. Some devices use two stop bits, some transmit a parity bit, etc. Check the documentation for the specific terminal in use for details.

Different input devices can send their information at different baud rates. The relationship between baud rate and bit times is given below. Many terminal devices have switch selectable baud rates. The user will typically wish to select the highest baud rate that will not exceed the capacity of the transmission path between the terminal and the Disk Jockey.

Baud Rate	Bit Times
110	9.091 ms
300	3.333 ms
1200	.833 ms
1800	.555 ms
2400	.417 ms
4800	.208 ms

A program assembling an incoming character will typically wait one-half bit time from the leading edge of the start bit and then check to see if the start bit is still present. It will then wait for one bit time and sample the first data bit, etc. This technique samples each bit in the middle of its bit time, greatly improving the noise immunity of the system.

The firmware supplied with the Disk Jockey will assemble a byte received from the interface using this technique. A timing loop generates the delay of one-half bit time or one bit time as required. The timing constant used by this loop must be set by the user according to the baud rate required. The value of the constant also depends on the instruction execution time of the processor.

To set the timing constant, a program must deposit the correct value appropriate to the baud rate and processor speed at a fixed location in the Disk Jockey's onboard RAM. The address of this location has the symbolic name SCON in the firmware listing. Its value in units with a standard decoding ROM is 342:160. The value of the timing constant that should be stored for commonly used baud rates is given in the following table for a 2 Mhz clock system:

Baud Rate	SCON decimal	SCON octal	SCON hex
110	375	1:166	176
300	135	207	87
1200	33	41	21
1800	21	25	15
2400	16	20	. 10
4800	7	7	7

This timing constant must be set before the serial input or serial output subroutine is called.

The subroutine TERMIN can be called to wait for and assemble a character arriving over the serial input line. The character, once it is assembled, will be returned in CPU register A. The subroutine will not return until a character arrives.

The subroutine TRMOUT will transmit out the serial output port the character in register A. It will return once the character has been transmitted.

It is important to be aware of some of the disadvantages of serial communications as they affect programming technique. Since the CPU must sample the input line at the correct time and be listening to detect the start bit, the serial input subroutine must be called before the

start bit is received and must stay in its timing loops until a character is assembled. Once a character is received, all computation done by the program that called for the character must be done and control returned to the serial input subroutine before the start of the next character. This is normally not a problem due to the CPU instruction times being so short with respect to the bit times for commonly used baud rates, but attention must sometimes be paid to this requirement. Another disadvantage of this technique is the requirement to tie up the CPU to listen to see if a character is arriving. This prohibits a program that runs continuously until a key is pressed unless the program can guarantee that it can check for a start bit sufficiently frequently while doing the other computation that is necessary.

The performance of the serial input system can be improved when it becomes desirable to do so either by adding a more complex serial interface including a UART chip that will assemble an entire character without requiring intervention from the CPU or by using a parallel interface with a flag bit.

In spite of the disadvantages listed, serial communications offers the most economical way to attach a terminal to a computer. Virtually every terminal device manufactured offers a serial data path for connection to its CPU.

Disk I/O

To understand the significance of the disk utility subroutines, it is necessary to say a few words about how data is organized on the disk.

Information on the disk is organized into a number of concentric tracks:

Full-sized floppy 77 tracks
Mini-floppy 35 tracks

The disk read/write head can be moved to any track by a series of step in or step out commands. A step in command moves the read/write head one track towards the center of the disk. A step out command moves the head one track away from the center of the disk. It is the responsibility of software to keep track of what track number the disk is currently at and to calculate how many step in or step out commands are necessary to move the head to a desired new position.

Once the read/write head has been moved to the desired track, the rotation of the disk will move a circle of magnetic material beneath the head. Within this circle of material, data is recorded in distinct regions called sectors. The sector is the smallest amount of information that can be separately read or written from the disk. There are twenty-six sectors/track on a standard and sixteen sectors/track on a mini-floppy. Each sector contains 128 data bytes.

In the header field of each sector, the track and sector number is recorded. During read or write commands, this header is read before data transfer takes place. If the DJ firmware detects a discrepancy between the track number it thinks it is at and the number recorded on the disk, it reports an error. To recover from this error, a program must do a recalibrate operation. The disk drive has a sensor that reports when the disk is physically positioned at track $\emptyset\emptyset$. A series of step out commands must be issued until this status line comes on. This operation will always position the disk to the same physical track. The recalibrate sequence is a standard utility subroutine supplied with the disk firmware.

Transferring a sector of disk data between memory and the disk therefore involves the following steps, each corresponding to a subroutine call to Disk Jockey firmware:

Position the read/write head to the desired track.

Specify the sector number to be involved in the data transfer.

Specify the memory address that disk data is to be written from or read to.

Actually perform the read or write operation.

Check for error conditions.

<u>Subroutines</u>

- TRKSET Given a one byte track address between 0 and 76 in register C of the CPU, this subroutine verifies that it is a valid track address. It then performs a seek operation to position the read/write head to the desired track. It does not read the header fields of the track to verify that the seek took place correctly -- this is doen when an attempt is made to read or write a sector.
- SECTOR Given a one byte sector number between 1 and 26 in register C, this subroutine checks that it is a valid sector number. It then records it for use in a later read or write.
- DMA Given a memory address in the B-C register pair, this subroutine checks that it is not an address within the Disk Jockey's address space. It then records the address in its RAM. Any subsequent read or write operations will transfer data to or from this memory address (and the 127 successively higher numbered memory cells). No disk operations actually take place.
- READ 128 bytes of data are read off the current track from the sector specified by the last SECTOR subroutine call to the address in memory specified by the last DMA call.
- WRITE 128 bytes of data are written on the current track into the sector specified by the last sector subroutine call from the address in memory specified by the last DMA call.

- TKZERO Calling this subroutine will position the read/write head to track 00 using the track 00 sensor as a reference. As discussed above, a call to this subroutine is the appropriate response to a seek error -- that is, the software detecting that the track number recorded in the header field of a sector does not agree with the count of track number being maintained in software.
- DBOOT Branching to this routine will result in a bootstrap load operation from the floppy disk. The serial constant is initialized to 1200 baud. Sector 1 of track ØØ of disk 1 will be read into memory location 200Q (80H). A branch will then be performed to this location. Branching to this subroutine from the front panel is the typical way to initialize a disk system.
- DISKETTE INITIALIZATION Before a new diskette can be successfully used, it must be initialized. Most diskettes are sold pre-initialized. However, it is sometimes necessary to reinitialize a diskette. The initialization process involves writing the header field of every sector onto the disk. None of the subroutines described above can be used to write these header fields. This is a safety measure to ensure that an erroneous branch to the firmware PROM cannot reinitialize a disk, destroying all the data recorded on it. The initialization function for diskettes is typically provided by a command included in the Disk Operating System. CP/M diskettes furnished by Thinker Toys contains a command INTLIZE.COM. The source code for this command, INTLIZE.ASM is also included. The Disk/ATE diskette has a command file called DSKINT which initializes diskettes and places a bootstrap loader on track zero sector one.

SOFTWARE SPECIFICATIONS

The Bootstrap

At location 340:000Q (E000H), there is a bootstrap routine which is intended to make the disk come up automatically. This routine selects drive number 1, does a home to track ØØ and reads sector 1 of track ØØ into locations 200Q through 377Q. After reading this sector into memory, the routine branches to location 200Q. The first sector of track ØØ should have a loader which will bring in the rest of the user's system. Customers who purchase CP/M with The Disk Jockey receive a diskette which has this loader. By setting the program counter of the CPU to 340:000Q (E000H) and pressing RUN, CP/M automatically loads itself into memory and starts.

Full-Sized Floppy/Mini-Floppy Hardware Configuration

The Disk Jockey controller has two slide switches which configure the board hardware to control either an 8" drive or a 5" drive.

8" Full-Sized Floppy Drives: BOTH slide switches must be in the LEFT-most position;

5" Mini-Floppy Drives: BOTH slide switches must be in the RIGHT-most position.

<u>Selecting Drives</u>

If decoders are installed in the disk drives themselves, up to eight drives can be connected in a daisy chain fashion to the Disk Jockey. Without decoders, four drives can be controlled by the Disk Jockey. It is important to note that none of the firmware with the exception of the boot loader ever changes the selection of drives. Before sending commands via the firmware, it is the programmer's responsibility to select the proper drive by storing the drive selection code in two places in memory.

Drive selection codes for drives 1, 2, 3 and 4 are as follows:

Drive	Drive Sele	ction Code
1	2100	88H
2	011Q	09H
3	050Q	28H
4	0120	OAH

The appropriate drive selection code must be stored at memory locations 342:1630 (E273H) and 343:002 (E302H).

Utilizing the Disk Jockey Firmware

In order to transfer information to and from a disk drive under the control of the Disk Jockey, the firmware functions must be used in the proper sequence. A seek function should always precede a set sector function. A read or write function should be preceded by a DMA function, a set sector function and a seek function.

Data Transfer Examples

READ:

Suppose sectors 5, 6, 7 and 8 of track 12, drive 1, need to be read into memory starting at location 7:200Q (780H). The following program will do this:

Octal

Hex

4B

CD 09 E0

100 102	076 210 062 163 342	READ	MVI	A,210Q	
102	062 003 345			DRIVE FUNCTN	
110	062 002 343 061 150 342		LXI	SD STACK	INITIALIZE STACK POINTER
113	315 011 340			TKZERO	
116	016 014		MVI	C,12	
120	315 014 340		CALL	TSEEK	TRACK 12
123	001 005 004		LXI		SECTOR COUNT & NUMBER
126	041 200 007				DMA ADDRESS
131	305	LOOP	PUSH	В	SAVE SECTOR & COUNT
132	345			Н	
133	315 017 340		CALL	SECTOR	
136	301		POP	В	GET DMA ADDRESS
137	305		PUSH	В	SAVE IT AGAIN
140	315 022 340		CALL	DMA	
143	315 025 340		CALL	DISKR	
146	247		ana	Α	TEST FOR ERRORS
147	302 100 000		JNZ	READ	START OVER IF ERROR
152	341		POP	H	RECOVER DMA ADDRESS
153	301		POP	В	RECOVER SECTOR & COUNT
154	005		DCR	В	DECREMENT COUNT
155	310		RZ	^	RETURN IF DONE
156	014		INR	C	INCREMENT SECTOR
157	021 200 000		LXI	D,200Q	INCREMENT DMA
162	031		DAD		ADDRESS BY 200Q
163	303 131 000		JMP	LOOP	READ NEXT SECTOR
40	3E 88	START	MVI	A,88H	
42	32 73 E2			DRIVE	
45	32 02 E3			FUNCTN	
48 -	31 68 E2		LXI	SP,STACK	

TKZERO

CALL

Soft Ware Specifications

4E	OE 1	0C			MVI	C,12
50	ÇD 1	OC.	E0		CALL	TSEEK
53	01	05	04		LXI	B,405H
56	21	80	07		LXI	H,780H
59	C5			LOOP	PUSH	В
5A	E5				PUSH	H
5B	CD	0F	E0		CALL	SECTOR
5E	C1				POP	В
5F	C5				PUSH	В
60	CD	12	E 0		CALL	DMA
63	ÇD	15	E0		CALL	DISKR
66	A7				ANA	Α
67	C2	40	00		JNZ	START
6A	E1				POP	H
6B	C1				POP	В
6C	05				DCR	В
6D	¢8				RZ	
6E	00				INR	С
6F	11	80	00		LXI	D,80H
72	19				DAD	D
73	C3	59			J M P	L00P

WRITE:

The following program writes from memory starting at location 200:000Q (8000H) onto tracks 4, 5, and 6 of disk drive 1.

Octal

200 202 205 210 213 216 220 223 224 227 235 236	076 210 062 163 342 062 002 343 061 150 342 315 011 340 076 004 062 305 000 117 315 014 340 001 001 032 041 000 200 305 345	WRITE TLOOP SLOOP	MVI STA STA LXI CALL MVI STA MOV CALL LXI LXI PUSH PUSH	B,32:001Q H,200:000 B H	RECALIBRATE DISK HEAD SET UP TRACK REGISTER SEEK TO NEXT TRACK SECTOR COUNT& NUMBER DMA ADDRESS SAVE SECTOR & COUNT SAVE DMA ADDRESS
237 242 243 246 251 254 255 260 261 262 265	315 017 340 301 305 315 022 340 315 030 340 247 302 255 000 341 301 021 200 000 031	STALL	CALL POP PUSH CALL CALL ANA JNZ POP POP LXI DAD	SECTOR B DMA DISKW A STALL H B D,200Q D	SET CURRENT SECTOR RECOVER DMA ADDRESS & SAVE SET DMA ADDRESS WRITE A SECTOR TEST FOR WRITE PROTECTED DISKETTE RECOVER DMA ADDRESS RECOVER SECTOR & COUNT INCREMENT DMA ADDRESS

Soft Ware Specifications

	266 267 270 273 276 277 301 302 305	014 005 302 235 0 072 305 0 074 376 007 310 303 220 0	000	TEMP:	INR DCR JNZ LDA INR CPI RZ JMP DB	C B SLOOP TEMP A 7 TLOOP	INCREMENT SECTOR DECREMENT COUNT READ ANOTHER SECTOR GET OLD TRACK & INCREMENT TEST FOR DONE SEEK TO NEXT TRACK TEMPORARY TRACK REG
Нех							
	80 82 85 88 8B 8E	3E 88 32 73 E2 32 02 E3 31 68 E2 CD 09 E0 3E 04		WRITE	MVI STA STA LXI CALL MVI	A,88H DRIVE FUNCTN SP,STACK TKZERO A,4	
	90	32 C5 00		TL00P	STA MOV	TEMP C,A	
	93 94	4F CD OC EO			CALL	TSEEK	
	97	01 01 1A			LXI	B,1A01H	
	9A	21 00 80		CL OOD	LXI	H,8000H	
	9D 9E	C5		SL00P	PUSH PUSH	B H	
	9E 9F	E5 CD OF 34			CALL	SECTOR	
	A2	CD OI 34			POP	В	
	A3	C5			PUSH	B	
	A6	CD 12 34			CALL	DMA	
	A9	CD 18 34			CALL	DISKW	
	AC	A7			ANA		
	AD	C2 AD 00		STALL	JNZ	STALL	
	B0 B1	E1 C1			P0P P0P	H B	
	B2	11 80 00			LXI	B,80H	
	B5	19			DAD	D	
	B6	ŌC			INR	C	
	B7	05			DCR	В	
	B8	C2 9D 00			JNZ	SLOOP	
	BB	3A C5 00	t		LDA	TEMP	
	BE BF	3C FE 07			INR CPI	A 7	
	C1	C8			RZ	•	
	C2	C3 90 00)		JMP	TLOOP	
	C5	00		TEMP:	DB	0	

The two examples above will appear to run rather slowly. This is due to reads and writes being done on consecutive sectors. Since there is overhead error checking done in the disk read and write firmware, the next sector has been passed over by the time the subsequent read or write operation commences. Therefore, the disk has to go almost a full revolution before the right sector is under the head again. Thus, to write twenty-six sectors, the diskette revolves roughly twenty-six revolutions instead of one. The method used to overcome this problem is called reading skewed sectors. That is, instead of reading, say, sectors 1, 2, 3, 4, 5 ..., 26, in consecutive order, it is more efficient to read in the order 1, 6, 11, 16, 21, 26, 5, 10, 15, 20, 25, 4, 9, 14, 19, 24, 3, 8, 13, 18, 23, 2, 7, 12, 17, 22. This will read an entire track in five revolutions instead of twenty-six. For an example of how to read and write sectors with a skew of five, see the software listings for patching CP/M* to the Disk Jockey.

^{*}CP/M is a trade mark of Digital Research, Pacific Grove, CA.

DISK SYSTEM SOFTWARE

An assembled Disk Jockey is part of a DISCUS I system and is also accompanied by a copy of Disk/ATEtm. Disk/ATE and CP/M* are also available at additional cost to those who have purchased only the controller. Both Disk/ATE and Disk Jockey CP/M are tailored to the I/O of the Disk Jockey controller. Both expect that a serial TTY/RS232 terminal set for 1200 baud is connected to J2 (serial port) of the Disk Jockey. Both are supplied on a write protected diskette (notch open) which should be kept that way. DO NOT COVER THE NOTCH ON THE DISKETTE. Finally, both systems are designed to self load when the disk is in place in drive A and a branch to 340:000Q (E000H) is made. For the CP/M user, a series of manuals accompanies the diskette which describes how to back-up the CP/M diskette. The only precaution is that when drive B is to be used for back-up purposes, it must be "logged on" (e.g., DIR:B) before the back-up process is started.

Backing Up Disk/ATE

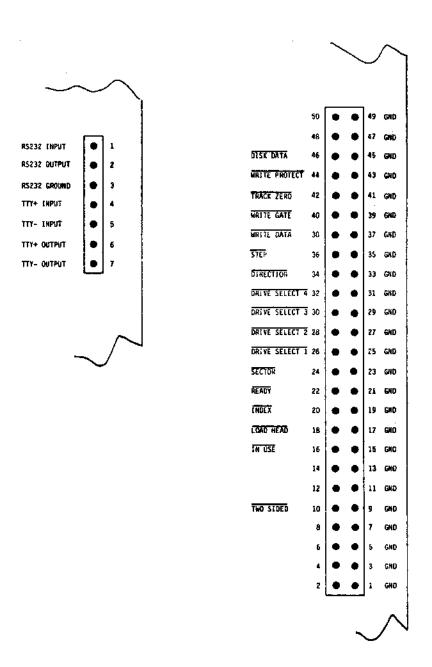
To make a back-up copy of Disk/ATE, load Disk/ATE and have a blank diskette which is <u>not</u> write protected (notch covered). Follow the steps outlined below:

- a. Type: B16 followed by a carriage return. This command forces ATE to express numbers and addresses in hexidecimal radix.
- b. Type: L IOTBL <T> followed by a carriage return. This command loads the I/O driver symbol table from the disk.
- c. Type: ? SYSIO...IOEND followed by a carriage return. This is the beginning and ending addresses of the ATE I/O driver software expressed in hexidecimal radix. Make a note of these two values.
- d. Type: L ATETBL <T> followed by a carriage return. This command loads a selected subset of ATE's symbol table from the disk.
- e. Type: ? BEGIN...END
 followed by a carriage return. This is thebeginning and ending
 address of ATE expressed in hexidecimal radix. Make a note of these
 two values.
- f. Type: GO DSKINT followed by a carriage return. This is the load and go command for the diskette initialization program that formats a diskette to the IBM soft-sectored standard and also places a boot strap loader on track zero sector one.

- g. DSKINT will ask that the diskette be placed in the drive and then will ask which drive the diskette is in. If the diskette is write protected or the diskette is not in the drive or if the drive door is not closed or if the indicated drive is not connected to the system, DSKINT will detect these conditions and start all over until all the conditions necessary to write on a diskette are satisfied. After completing the initialization, DSKINT automatically returns to ATE.
- h. IO and ATE must be saved on the new diskette. IO must be saved first.
- i. Using the values for SYSIO and IOEND obtained in step c, type: S IO (SYSIO value here)H...(IO END value here)H followed by a carriage return. The "H" suffix is necessary to force ATE to interpret the preceeding digits as a number in hexidecimal radix.
- j. Using the values for BEGIN and END obtained in step e, type: S ATE (BEGIN value here)H...(END value here)H followed by a carriage return.
- k. Disk/ATE has now been copied on the fresh diskette. Files may now be transferred from the original diskette as needed.

I/O CONNECTORS J1 AND J2

Illustrated below are the details of the pin connections of J1 and J2. In both illustrations, the top of the circuit board is shown as the straight line on the right side of the connector.



PATCHES FOR CP/M*

General

This section is included for those users of the Disk Jockey who have purchased a copy of CP/M Vers. 1.4 from a source <u>other</u> than Thinker Toys. Copies of CP/M sold through Thinker Toys have the necessary I/O routines to interface CP/M to the Disk Jockey disk controller and serial I/O port.

At the end of this section are two listings which are designed to allow the Disk Jockey to be interfaced with the Digital Research CP/M operating system. This can be done with a minimum of effort.

The first listing is the so called "cold start loader" which is used to bring CP/M in from the disk. It also has code which will allow the user to easily write a modified version of CP/M out on the disk. There is even a small routine which loads the "cold start loader" itself on sector 1 of track 0.

The second listing is CBIOS software (Custom Basic Input-Output System) which is the interface between CP/M and the Disk Jockey controller. The general idea is to key in the cold start loader, use the loader to bring CP/M in from a diskette, enter the CBIOS code and, finally, use the cold start loader to save everything out on a clean diskette.

The "Cold Start Loader"

There are three parts to the cold start loader. LOAD is at address 200Q (80H) and is designed to read CP/M into memory from location 51:000 through 77:377 (290H through 3FFFH). After loading CP/M, the LOAD routine branches to location 76:000 which is a routine that initializes several memory locations, prints a sign-on message and then branches to CP/M proper.

SAVE is at location 232Q (9AH) and is the reverse of LOAD. SAVE writes out on the disk starting at track \emptyset sector 2 all memory locations between 51:000 and 77:377. After performing this operation, SAVE comes to a dynamic halt at STALL 261 (B1H).

INTLZ is a short routine which writes locations 200 through 377 on sector 1 of track \emptyset . Thus, once the cold start loader is keyed into memory, it can save itself at the right location on the disk.

CBIOS

The standard version of CP/M is designed to run with the Intel MDS development system a floppy disk interface. Most of the CP/M system software is completely independent of the particular 8080 hardware environment that it is running in. However, there is a certain part which must be tailored to the hardware of the host system. This hardware-dependent software is completely contained on pages 76 and 77 of CP/M (assuming the standard 16K version). CP/M can be made to run on different hardware by changing the software on pages 76 and 77 (3E00H through 3FFFH). The CBIOS software which is supplied with the Disk Jockey is designed to let CP/M run when an 8" full-sized floppy disk is attached to the Disk Jockey controller which is plugged into an S-100 main frame.

Patching CP/M*

Before actually performing any of the steps below, the Disk Jockey should be plugged into the mainframe S-100 bus, and an 8" disk drive should be connected to the controller. You should have on hand two diskettes: one with CP/M and one blank. A copy of CP/M which will run on the Disk Jockey will be constructed on the blank disk before any changes are attempted on the original CP/M disk.

Step I:

Plug in the controller. Connect the disk to the controller and turn on the CPU and disk drive. Do NOT put a diskette in the drive at this point.

Step II:

Enter the "cold start loader" into main memory starting at location 200 (80 H). The instructions will extend from 200 to 377 filling all of the second half of page 0.

Step III:

Load location 342:163 with 210 (E273 hex with 88 hex).

Load location 343:002 with 210 (E302 with 88 hex). Location 343:002 is in actuality an I/O device (memory mapped I/O) so values read from 343:002 and the values written into 343:002 will be independent.

Step IV:

Set the program counter of the CPU to location 272 (BA hex) but do NOT start the CPU yet.

Step V:

Insert the <u>blank</u> diskette into the drive and close the door. Be sure that the diskette is NOT write protected. (Write protected 8" diskettes have a notch near the corner of the diskette diagonally opposite the labeled corner.) Be sure the diskette is inserted right side up. On a Discus I system, the label should be on top and the diskette should be inserted by holding the label of the diskette between the thumb and fore-finger when it is inserted into the drive.

Patches for CP/M*

Step VI:

Start the computer. After the CPU is started at 272, the activity light (if one is present) should come on, the head should load and step out to track 0 (if not already there). After sixteen revolutions of the diskette, the head will unload and the activity light will go off.

Step VII:

Stop the CPU. It should be executing the instruction JMP DONE. -- 303 312 000 (C3 CA 00 hex). The cold start loader should be on sector 1 of track 0.

Step VIII:

Remove the diskette from the disk drive.

Step IX:

Change locations 222 and 223 from 000 and 076 to 261 and 000 respectively (change 92 to B1 and 93 to 00 hex).

Step X:

Initialize the program counter of the CPU to 200 (80 H). Do NOT start the CPU.

Step XI:

Insert the CP/M disk and be sure that it is write protected before it is put into the drive. Close the drive door securely.

Step XII:

Start the CPU. The head should load and after a second or two the head should step to track 1. Wait for the head to unload and the activity light to go off. CP/M has been loaded into memory between 51:000Q and 77:377Q.

Step XIII:

Enter the CBIOS code starting at location 76:000 through 77:072 (3E00 through 3F3A hex). Verify that the code has been entered correctly.

Patches for CP/M*

Step XIV:

Initialize the program counter of the CPU to 232 (9A hex) but do NOT start the CPU.

Step XV:

Insert the blank (except for the "cold start loader") diskette into the drive and close the door securely.

Step XVI:

Start the CPU. The head should load, return to track 0 and write the better part of tracks 0 and 1 before the head unloads. After the head unloads, do NOT remove the diskette but stop the CPU. The CPU should be executing the JMP STALL instruction.

Step XVII:

Connect a terminal to the serial port of the Disk Jockey and adjust the baud rate to 1200 baud.

Step XVIII:

Initialize the program counter of the CPU to location 340:000 (E000 hex) and start the CPU (the blank disk should still be in the drive). After several seconds, the terminal should print:

16k CP/M VERS/1.4

After five seconds or so, the prompt should appear:

A>

A Disk Jockey version of CP/M is now up and running. After this new version of CP/M has been tested (see CP/M documentation), Steps I through XVII can be used to alter the original CP/M diskette if desired.

CP/M* "COLD START LOADER"

	342:156 342:164 340:017 340:025 340:030 340:014 340:011 340:000	DMAADOR SECTOR READ WRITE SEEK HOME	EQU EQU EQU EQU EQU	342:156Q 342:164Q 340:017Q 340:025Q 340:030Q 340:030Q 340:014Q 340:011Q 340:000Q	
21	000:200 315 315 000 203 000 204 305 205 315 017 340 210 315 025 340 213 247 214 302 000 340 217 301 220 005 221 312 000 076 224 315 327 000 227 303 203 000	LOAD RLOOP	CALL NOP PUSH CALL CALL ANA JNZ POP DCR JZ CALL JMP	SETUP B SECTOR READ A BOOYSTRAP S B 76:000Q NSDMA RLOOP	INITIALIZE SECTOR COUNT & DMA ADDRESS SAVE SECTOR & COUNT INITIALIZE SECTOR READ DATA TEST FOR ERROR RESTART IF ERROR RECOVER SECTOR & COUNT DECREMENT & TEST FOR CP/M LOADED CALCULATE NEXT SECTOR & DMA ADDR
	232 061 156 342 235 315 011 340 240 315 315 000 243 000 244 305 245 315 017 340 250 315 030 340 253 247 254 302 261 000 267 302 261 000 264 315 327 000 267 303 243 000	STALL	LXI CALL NOP PUSH CALL CALL ANA JNZ POP DCR JZ CALL JMP	SP,STACK HOME SETUP B SECTOR WRITE A STALL B STALL NSDMA WLOOP	INITIALIZE STACK POINTER RETURN HEAD TO TRACK Ø INITIALIZE SECTOR, COUNT & OMA ADDR SAVE SECTOR & COUNT GET SECTOR ADDRESS WRITE DATA TEST FOR ERROR RECOVER SECTOR & COUNT DEGREMENT COUNT DYNAMIC HALT CALCULATE NEXT SECTOR & OMA ADDR
	272 061 156 342 275 315 011 340 300 041 200 000 303 042 164 342 306 315 030 346 311 000 312 303 312 000) - 	LXI CALL LXI SHLO CALL NOP JMP	H,200Q DMAADDR	INITIALIZE STACK POINTER RETURN HEAD TO TRACK Ø INITIALIZE DMA ADDRESS WRITE 2000 THROUGH 3770 ON TRACK Ø SECTR 1 DYNAMIC HALT

"Cold Start Loader"

320 041 000 053	
327 076 005 NSDMA MVI A,5 SECTOR SKEW 331 201 ADD C ADD TO CURRENT 332 117 MOV C,A SECTOR 333 336 033 SB! 27 TEST FOR 3340 074 INR A ADJUST 341 117 MOV C,A SECTOR OVERFLOW 340 074 INR A ADJUST 341 117 MOV C,A SECTOR OVERFLOW 342 021 200 365 LXI D,365:200Q DMA ADJUSTMENT 345 376 001 CPI 1 TEST FOR SECTOR 1 347 302 362 000 JNZ OK+3 NOT THROUGH W/ TRACK 9 352 305 PUSH B THROUGH W/ TRACK 9 353 315 014 340 CALL SEEK MOVE TO TRACK 1 366 301 POP B RECOVER SECTOR & COUNTY 357 021 200 002 OK LXI D,2:200Q DMA ADJUSTMENT 366 031 DAO D CALCULATE NEN 366 042 164 342 SHLD DMAADDR GET OLD DMA ADDR 366 042 164 342 SHLD DMAADDR TEST FOR 371 076 077 MVI A,77Q TEST FOR 372 224 SUB H ADDRESS OVERFLOW 374 372 327 000 JM NSDMA & ADJUST IF 377 311 RET NECESSARY	

Hex Listings (see octal listings for comments)

CP/M* "COLD START LOADER"

	E26E	STACK	EQU	342:156Q
	E274	OMAADDR	EQU	342:1640
	EOOF	SECTOR	ะี้ดับ	340:0170
				340:0250
	E015	READ	EQU	
	EQ18	WRITE	EQU	340:030Q
	E00C	SEEK	EQU	340:0140
	E009	HOME	EOU	340:0110
	£000	BOOTSTRA		340:0000
	£000	BOOTSTI	ii EQU	•10/2004
1080	CD CD CO	LOAD	CALL	SETUP
	00	RLOOP	NOP	
0083	C5	KEOO	PUSH	В
0084			CALL	SECTOR
0085	CD OF EO			
38800	CD 15 EO		CALL	READ
B880	A7		ANA	A
0080	C2 00 E0		JNZ	BOOTSTRAP
008F	Cl		POP	В
	05		DCR	8
0090			JZ	75:0000
0091	CA 00 3E			
0094	CD 07 00		CALL	NSDMA
0097	C3 83 00		JMP	RL00P
2005	31 6E E2	SAVE	LXI	SP,STACK
009A		SWAE	CALL	HOME
009D	CD 09 E0			
0000	CD CD 00		CALL	SETUP
00A3	00	WLOOP	N05	
00A4	C5		PUSH	8
00A5	CD OF EO		CALL	SECTOR
BAGO	CD 18 E0		CALL	WRITE
	A7		ANA	A
BAOO			JNZ	STALL
OOAC	C2 B1 00			
3A00	C1		POP	В
0080	05		OCR	В
0081	CA B1 00	STALL	JŽ	STALL
0084	CD D7 00		CALL	NS DMA
0087	C3 A3 00		JMP	WLOOP
0007	C3 H3 00		•	
OOBA	31 6E E2	INTLZ	LX1	SP,STACK
0080	CD 09 E0		CALL	HOME
	21 80 00		LXI	H,2000
0000			SHLD	DMAADDR
0003	22 74 E2			WRITE
0006	CD 18 EO		CALL	#KIIE
0009	QÛ		NOP	
OOCA	C3 CA 00	DONE	JMP	DONE
	•			

"Cold Start Loader"

0000 0003 0006	01 06 2E 21 00 2B 22 74 E2 C9	SETUP	LXI LXI SHLD RET	B,56:006Q H,53:000Q DMAADDR
0007 0009 000A 000B 0000 00E0 00E1	3E 05 81 4F DE 1B FA EF 00 3C 4F	AMOZN	MVI ADD MOV SBI JM INR MOV	A.5 C.A C.A 27 OK A C.A
00E2 00E2 00E7 00EA 00EB 00EF 00F2 00F5 00F6 00FB 00FC	11 80 F5 FE 01 C2 F2 00 C5 CD 0C E0 C1 11 80 02 2A 74 E2 19 22 74 E2 3E 3F 94 FA D7 00 C9	0K	CYL CPI JNZ PUSH CALL POP LHLD OAD SHLD MVI SUB JM RET	D,365:200Q 1 OK+3 B SEEK B D,2:200Q DMAADOR DMAADDR A,77Q H NSDMA

TYPE B:TTCBIOS.PRN

		;CBIOS	DRIVERS	FOR CP/M	
		1			
2900	=	CPM	EQU	29008	
3106	=	ENTRY	EQU	Ç₽M+806H	
0004	=	CDISK	EQU	4	
E000	=	ORIGIN	EQU	HCDOBO	
E003	T	THPUT	EQU	ORIGEN+3	
E006		OUTPUT	EQU	ORIGIN+8	
£009	=	TKZERO	EQU	ORIGIN+9H	
EOOC		TSEEK	5QU	ORIGIN+OC	
EDOF	E	SECTOR	ΕQU	ORIGIN+OF	
E012		DMA	EQU	ORIGIN+12	
E015	=	DISKA	EGN	ORIGIN+15	
E018	=	DISKW	EQU	ORIGIN+18	
ES9E	=	STACK	EQU	ORIGIN+25	
E273	=	DRIVE	EQU	ORIGIN+27	
E274	=	DMAADR	EQU	ORIGIN+27	
E277	=	TRACK	EON	ORIGIN+27	
£302	=	STATUS	EQU	ORIGIN+30	
E1F3	•	DELAY	EQU	ORIGIN+1F	
E270		SCON	EQU	ORIGIN+27	UH
0003	=	TOBYTE	EQU	ЗН	
0000	=	INTION		0	
	2	RDYMSK	EQU	SOH	
0001	=	WPTCT	EQU	1	
		:			
		•			
3E00			ORG	CPM+1500H	
32.00		Ė			
3E00	C32D3E	START	JMP	300T	
3E03	C36D3E	D1	JMP	WBOOT	
3E06	C3973E		JMP	CONST	
3E09	C3C53E		JMP	CONIN	
3EOC	C3D73E	CPOUT	9 K L	CONQUE	
BOF	C3F23E		JMP	LIST	
3E12			JMP	PUNCH	
3E 15	C3E73E C3DD3E		JMP	READER	
3E18	C3913E		JMP	HOME	
3E 18	C3343F		JMP	SELDISK	
3E,1E	C3BD3E		9ML	SETTRK	
3E21	C3823F		JMP	SETSEC	
3E24	C312E0		1Mb	DHA	
3E27	C38C3F		JMP	READ	
3E2A	C3A13F		JMP	WRITE	

3E2D 316E 3E3D CD61 3E33 21CC 3E39 AF 3E3A 3204 3E3D 328B 3E3D 328B 3E4D 3EC3 3E4B 3EC3 3E4B 2103 3E4B 2205 3E5E 2206 3E5E 2206 3E5E 4300 3E5E 4300 3E5E 4300 3E5E 4300 3E5E 4300 3E5E 4300 3E5E 4300 3E5E 4300	3E 3F 3F 00 GOCP E2 300 13E 000 631 000	LXI GALL LXRA STA STA LXI SHLD HVI SHLD STA LXI SHLD STA LXI SHLD STA LXI SHLD STA LXI SHLD JMP	SP, STACK TINIT M, PROMPT MESSG A CDISK CDISKA H, BOH DMAADR A, OC3H O H, START+3 1 5 H, ENTRY 6 CDISK CPM
3E61 3E00 3E63 3203 3E66 2100 3E69 2100 3E6C C9	300 100	T MVI STA LXI LXI RET	A,INTIOBY IOBYTE H,O H,O
3E6D 3166 3E70 0180 3E73 CD13 3E76 3E88 3E78 CD23 3E78 CD23 3E78 CD23 3E81 2292 3E86 32C6 3E86 32C6 3E88 32C6 3E88 32C7 3E88 32C7	000 EE0 SE2 SE0 3E 000 500 500	T LXI LXI CALL MYI STA CALL LXI SHLD MVI STA MVI STA MVI STA JMP	SP, STACK B, 200Q DMA A, 210Q DRIVE ORIGIN+52Q H, GOCPH 222Q A, 52Q 317Q A, (CPM+1400H)/100H 372Q 200Q SELDSK
3E91 CD39 3E94 C309		JMP	TKZERO

'	3E97 0604 3E99 3A0ZE3 3E9C 1F	CONST	MVI I, DA AAR RAR	B,4 STATUS	3EF5 3EF8 3EF9		LIST	LXI LDA PAR RAB JMP	H, LTBLE TOBYTE PNCH1
	3E9D 1F 3E9E D2A73E		JNC	STATUS	366 M	C3E03E		JHE	146411
	3EA1 05		DCR	В	3EFD	293F	ĆITBLE	D₩	CITTY
	3EA2 C2993E		JNZ	CONST+2	3 EFF			DW	CICRT
	3EA5 AF		XAA	A	3F01			DW	READER
	3EA6 C9		RET	TARIT	3F03		COTOLE	DW DW	CIUCI
	3EA7 CDO3EO	STATUS1	CPI	INPUT 19	3F05 3F07		COTBLE	DW	COCRT
	3EAA FE13 3EAC C2B13E		JNZ	STATUS4	3809			DW	LIST
	3EAF AF		ARX	A	3F0B			DW	COUC 1
	3EBO C9		RET	•	3FOD		LTBLE	D₩	CCTTY
	ĴEB1 FĒOD	STATUS4		ODH	3F OF			₽₩	COCRT
	3EB3 C2A73E		JNZ	STATUS1	3F11			DW	COLPT
	3EB6 32BC3E		STA	STATUS5	3F13 3F15		PTBLE	DW DW	COUL 1 COTTY
	3E89 F6FF		ORI RET	OFFH	3F 17		FEBLE	D₩	COPTP
	3E89 C9 3E8C OO	STATUS5		0	3F 19			DW	COUP1
	3500 00	:	75		3F 1 B			DW	COUP2
	3EBD C5	SETTRK	PUSH	8	3F 1D		RTBLE	DW	CITTY
	3E8E CD393F		CALL	SELDSK	3F 1F			D₩	CIPTR
Ņ	3EC1 C1		POP	Ð	3F21			DW	CIURI
4-	3EC2 C3GCEO		JMP	TSEEK	3F23 3F25		COTTY	E O U	CIUR2 \$
		;			3F 25		COCRT	EQU	i
	3EC5 21FD3E	CONIN	LXI	H, CITBLE	3F 25		COUCI	EQU	Š
	3EC8 3A0300	CONINI	L DA BAL	IOBYTE	3825		COLPT	FQU	\$
	3ECB 17 3ECC E606	SELDEV	ANI	6н	3825		COUL 1	EQU	\$
	3ECE 1600	JELDER	MVI	D.0	3F 25		COPTP	EQU	\$
	3EDO SF		MOV	E,A	3F25		COUPI	EQU	\$
	3ED1 19		DAD	D,	3F25		COUP2	EQU	3
	3ED2 7Ē		MOV	A,M	3F25	C306E0		MOV JMP	A,C OUTPUT
	3ED3 23		INX	¥	,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,	0,0040			001101
	3ED4 66		MOV MOV	H, H	3F 29	=	CITTY	EQU	\$
	3ED5 6F 3ED6 E9		PCHL	L,A	3F 29	3	CICRT	EQU	3
	3200 29	•	1 4112		3F 29		CIPTR	EQU	\$
	3ED7 21053F	CONOUT	LXI	H,COTBLE	3F29	=	CIUCI	EQU	*
	3EDA C3C93E		JMP	CÓNIN1	3F29 3F29		CIUR1 CIUR2	EQU	\$ \$
	32277	1				21BC3E	CIORE	LXI	H.STATUS5
	3EDD 211D3F	READER		H, RTBLE	3F2C			MÔV	A, M
	3EE0 3A0300	001001	LDA	IOBYTÉ	3F2D			ORA	A
	3EE3 1F	READRI	RAR Jmp	SELDEV	3F2E	CAGSEO		JZ	INPUT:
	3664 C3CC3E		JHF	SELDET	3F 3 1			MVI	M.O
	3EE7 21153F	PUNCH	LXI	H.PTBLE	3F 33	C9		RET	
	3EEA 3A0300		LDA	IOBYTE			i		
	3EED 1F	PNCH1	BAR		3F 34		SELDISK		A,C
	3888 1F		AAR			323D3F		STA	NEWDRV+1
	BEEF CBEBBE		JMP	READR 1	3F38	C 9		RET	
		;							

3F30 21 3F3C 3E 3F3F 77 3F40 C8 3F41 4F		SELDSK New Drv	HOV HOV RZ HOV	H,CDISKA A,O M M,A C,A
3F45 F5 3F46 E6	7352	i SELDSK1	LDA PUSH ANI LXI	DRIVE PSW OF7H H,DTABLE
3F4B E5 3F4C BE 3F4D CA 3F50 23 3F51 23		SLOOP	PUSH CMP JZ INX INX	H H H
3F52 C3 3F55 23 3F56 3A 3F59 77	403F 77E2	FOUND	JMP INX LDA MOV	H Sloop H Track M,a
3F5A £1 3F5B 79 3F5C 87 3F5D 85 3F5E 6F 3F5F D2 3F62 24	633F		POP MOV ADD ADD MOV JNC INR	H A,C A L L,A HISOK H
3F64 E6 3F64 E6 3F66 B6 3F67 32 3F6A 32 3F6D 23		HIŜOK	POP ANI ORA STA STA INX	PSW B M STATUS DRIVE H
3F72 AF 3F73 11	77E2 02E3 F3E1	DTABLĒ	MOV STA XRA LXI JMP DB DB DB DB DB	A,M TRACK A D,STATUS DELAY BO 1 0 20H 0
3F7F 02 3F80 00 3F81 00		CDISKA	DB DB DB	0 0 5
3F82 79 3F83 32 3F86 C9	883F	SETSEC	MOV Sta Ret	A,C SECTORA+1
3F87 OE 3F89 C3	00 OFEO	SECTORA ;	HVI JMP	C,O SECTOR

3F8C 3F8F 3F92 3F94 3F95 3F98 3F9A 3F9A 3F9C 3F9F 3FAO	0E0A C5 CD15E0 C1 A7 C5 OD C2943F 2F	READ 1 READ 1 READ 1	CALL CALL MVI PUSH CALL POP ANA BZ DCB JNZ CMA RET	SELDSK SECTORA C,10 8 DISKA B A C READ1
3FA1 3FA4 3FA9 3FAA 3FAD 3FAE 3FAF	OD CO	₩RITE DSKW	CALL CALL MVI PUSH CALL POP ANA RZ DCR	SELDAK SECTORA C,10 B DISKW B A
3F81 3F84 3F87 3F89 3F8C 3F8F 3FC0	3A02E3 E601 CA9F3F 21E23F 7E	PROTCT Messc	JNZ LDA ANI JZ LXI MOV	DSKW STATUS WPTCT READY H, PTCTMSG A, M
3FC1 3FC2 3FC3 3FC4 3FC7 3FC8 3FC9	E5 4F CDOC3E E1 23		FNA RZ PUSH MOV CALL POP INX JMP	A H C, A CPOUT H M MESSG
3FCC 3FCC 3FCE 3FD7 3FDC 3FDF	ODOA 3136482043 5645525320 312834 ODOA		DS DB DB DB DB	0 00H,0AH '16K CP/M ' 'VERS ' '1.4' 00H,0AH
3FE1 3FE2 3FE4 3FEB 3FED		PTCTMSG	08 08 08 08 08	O ODH, GAH 'PROTECT' ODH, OAH O

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HARDWARE LEVEL REGISTERS

Users desiring a greater level of control over the floppy disk or serial interface may wish to directly refer to the I/O device registers on the DJ from their 8080 program. There are eight one-byte registers, four of them are read only, and four of them are write only. The registers have four memory addresses on the S-100 bus with a different register being selected during a read operation and during a write operation. This gives the total of eight registers.

To make use of the registers, the first operation is to write to the write function register a bit pattern that selects one of the disk drives of the floppy disk system. All the other command and status bits will then refer to the operation of this selected drive.

Readable Registers

Register 0 - Disk Data Register (location 343:000Q or E300H standard version):

Reading from this register will transfer one byte of information from the floppy disk to the CPU. Note that the correct use of this register requires programming sequences involving setting bits in other control registers and adhering to certain timing rules, or the data read will be meaningless.

Register 1 - Disk Read Mark Register (loc. 343:0010 or E301H standard version):

Reading from this register puts the CPU into a wait state until the DJ detects a field it recognizes as an address mark on the floppy disk. (See floppy disk data formats.) The bit pattern of the mark then appears as the contents of the register, and the CPU is allowed to leave the wait state and read in the register value.

If no valid data is found on the disk, the CPU will hang and an external reset will be necessary. It is therefore important for a disk to be in the drive, the door closed, the drive selected, and the head loaded and positioned to a valid track before a read mark is done.

Register 2 - Read Status Register (Ioc. 343:002Q or E302H standard version):

This register contains bits that identify the current status of the DJ and the currently selected drive.

BIT 7 6 5 4 3 2 1 Ø

READY SECTOR 2SIDED INDEX TRKØØ HD-LOADED SER-IN WRT-PROT

*Bits marked with an asterisk reflect the current state of status lines from the floppy disk drive. See the documentation for the floppy disk drive actually in use for a detailed specification of these signals.

Hardware Level Registers

- READY This bit is a 1 when the currently selected disk drive is powered on, the door is closed, and a diskette is present.
- SECTOR This line relects the status of the sector status line from the floppy disk drive.
- 2SIDED This line is a 1 when the status line from the floppy disk drive indicates that it is equipped to read two-sided floppy disks.
- INDEX This line reflects the status of the INDEX line from the floppy disk drive. It goes to a 1 once per revolution of the floppy disk.
- TRKØØ This bit is a 1 when the read/write head in the floppy disk is positioned on track ØØ.
- HD-LOADED This bit is a 1 when the DJ is directing the floppy disk drive to press the read/write head against the recording medium.
- SER-IN This bit reflects the current status of the TTY or RS-232 serial input line.
- WRT-PROT This bit is a 1 when the diskette in the drive is write protected, that is, there is a notch in the jacket for 8" drives or there is no notch for 5" drives.

Register 3 - Load Head Register (loc. 340:003Q or E303H standard version):

Reading a byte of data from the address of this register causes the currently selected floppy disk drive to load the read/write head, that is, operate the solenoid to bring the head into contact with the medium. The byte of data that will actually be read from this memory location is meaningless. The head will remain loaded for sixteen revolutions after a load head command and will then automatically unload unless bit 3 of the write function register is a \emptyset .

Write Only Registers

Register \emptyset - Disk Data Register (loc. 340:000Q or E300H standard version):

The byte of data to be written to the disk is written to this register. As with the Read Disk Data Register, a proper sequence of setup commands is necessary for a write to this register to be effective.

Register 1 - Disk Write Mark Register (loc. 340:001Q or E301H standard vers.):

Writing a byte of data to this register, when done at the proper time and place in a programming sequence, can write one of the special synchronization marks onto the disk. These synchronization marks are recorded on the disk in a way that is immediately distinguishable from data since certain clock pulses are omitted in accordance with industry standards. The DJ automatically generates the correct pattern of missing clock pulses to be recorded with the mark written.

The standard marks and the hex value of the byte of data that should be written to this register to get them are:

- FB Data Mark
- F8 Deleted Data Mark (rarely used)
- FE Sector Mark (only used during initialization of disk)
- FC Index Mark (only used during initialization of disk)

Register 2 - Write Function Register (loc. 340:0020 or E302H standard vers.):

This register contains bits that activate various control signals to the floppy disk drive(s) and to the internal circuitry of the DJ.

- DSI-DS4 These bits are passed on to the output flat cable to the floppy disk drives as DRIVE SELECT 1 thru DRIVE SELECT 4. Each disk drive has a means of setting an internal jumper to determine which drive select line in the cable it will respond to. Only the disk drive whose drive select line is on will respond to the commands output by the program to the writable registers and present drive status in the readable registers. Likewise, only the selected drive will participate in data transfers. Only one drive should be selected at a time.
- IN-USE This line is normally not used. Setting it to a 1 forces the LED on the front of the Shugart compatible floppy disk drive to be on regardless of the state of other activity at the drive.
- STEP This bit controls the step line to the floppy disk drive. Setting this bit from a 1 to a Ø causes the stepping motor in the floppy disk drive to move the read/write head one track in the direction specified by the DIR bit.

Hardware Level Registers

DIR - This bit indicates to the floppy disk drive the direction it should move the read/write head in response to a step command:

$$\emptyset$$
 = out 1 = in

E-A-U Enable Auto Unload. This bit being a 1 activates the automatic head unload feature. To prolong medium life, it is important to minimize the amount of time the read/write head is pressed against the disk. When this bit is set to a 1, the DJ will automatically unload the read/write head sixteen revolutions after the last load head command was issued. This bit should always to set to a 1 except for hardware debugging or data transfer optimization purposes.

Register 3 - Write Serial Register (loc. 340:003Q or E303H standard version):

Bit 5 of this register will appear at an appropriate voltage (or current) level at the RS-232 (and TTY) serial output connector. The other bits of this register have no significance.

PARTS LIST

7	180Ω ¼ wat	t resistors	(brown-grey-brown)	R1, 8, 9, 10, 11, 12, 13
3	240Ω ½ wat	t resistors	(red-yellow-brown)	R21, 27, 28
1	47 0Ω 1 wat	t resistor	(yellow-purple-brown)	R20
4	510Ω ½ wat	t resistors	(green-brown-brown)	R22, 23, 25, 26
4	1kΩ	n .	(brown-black-red)	R3, 16, 17, 24
1	1.5kΩ	ţr	(brown-green-red)	R14
3	3.3kΩ	II .	(orange-orange-red)	R4, 7, 15
1	3.9kΩ	n	(orange-white-red)	R2
1	4.7kΩ	H	(yellow-purple-red)	R19
2	27kΩ	П	(red-purple-orange)	R5, 6
1	47k Ω	tt	(yellow-purple-orange)R18
1	.01 µfd di	sk capacitor		C1
17	by-pass car	pacitors*		
1	.82 µfd ta	ntulum capacit	or	C3
2	$2.7~\mu fd$	II .		C5, 6
2	$39~\mufd$	u		C2, 4
1	1N914/4820	-0201 signal d	iode	CR1
2	1N4742 12	volt zener di	odes	CR2, 3
1	2N3904 NP	N transistor		Q2
2	2N3906 PN	P transistors		Q1, 3
2	slide swit	ches		
1	heat sink			
1	set machin	e nut and scre	W	
1	4Mhz cryst	al		
1	50 pin rig	ht angle flat	cable connector	J1
1	7 pin righ	t angle connec	tor w/ matching header	and pins J2

Parts List

1	8 pin low profile socket	
16	14 pin low profile sockets	
14	16 pin low profile sockets	
5	20 pin low profile sockets	
1	74LS00 quad 2-input NAND gate	3C
1	74LS02 quad 2-input NOR gate	10B
1	74LS08 quad 2-input AND gate	5B
1	74LS10 tri 3-input NAND gate	3B
1	74LS14 hex Schmitt-trigger inverter	5C
2	74LS20 dual 4-input NAND gate	12C, 13A
1	74LS32 quad 2-input OR gate	8B
8	74LS74 dual D-type flip flop	9A, 10A, 1B, 2B, 12B, 13B, 1C, 2C
1	74LS124 dual voltage-controlled oscillator	8A
1	74LS155 dual 1 of 4 decoder	4C
2	74LS161 synchronous 4-bit counter	4B, 13C
1	74165 parallel-load 8-bit shift register	9B
1	74LS174 hex D-type latch with clear	7A
1	74LS241 octal tri-state** bus driver	10C
1	74LS273 octal D-type latch with clear	5A
1	74LS299 8-bit bidirectional shift/storage regi	ister 11C
2	74366/368 hex 2-4 tri-state** inverting bus dr	rivers 3A, 6A
1	DM8090 quad inverter, dual 2-input NAND gate	118
1	DM81LS95/97 octal tri-state** buffer	9C
1	MMI 6301/82S129 4x256 bit tri-state** PROM	8C
2	MMI 6306/82S131 4x512 bit tri-state** PROM	6C, 7C
1	DM81LS96/98 octal tri-state** inverting buffe	r 4A

Parts List

2	2112 4x256 bit RAM with tri-state** output	6B, 7B
1	7805 5 volt 1 amp monolithic regulator	1A
1	741 operation amplifier	11A
1	5" x 10" printed circuit board	
1	glossy photograph	

^{*}by-pass capacitors will vary in value from .01 μfd to .1 μfd depending upon currently available supplies.

^{**}tri-state is a trademark of National Semiconductor Corp.

ASSEMBLY INSTRUCTIONS

DO NOT INSTALL OR SOLDER ANY PARTS UNTIL YOU HAVE READ THESE INSTRUCTIONS SEVERAL TIMES AND HAVE FULLY DIGESTED THE INFORMATION!

CAUTION -- DO NOT SOLDER OR CLIP COMPONENT LEADS WITHOUT USING SAFETY GLASSES!

INSPECTION

Use the Parts List to make sure that there are no missing items in your kit. Please notify us of any shortages. Be sure to check for missing parts <u>before</u> you start to assemble.

COMPONENT LEAD WIDTHS

Bend the leads with the plastic bending block in your kit. Be sure that the leads are bent to the proper width <u>before</u> the part is inserted. Properly bent components solder easier and give the finished kit a professional appearance. See the right hand column of the Parts Installation for proper component widths.

SOCKETS

A socket is furnished for every integrated circuit. It is important that you use the sockets; otherwise, a defective part will be extremely difficult to replace.

NO REPAIR OR SERVICE WILL BE PERFORMED ON A KIT WHICH HAS HAD INTEGRATED CIRCUITS SOLDERED TO THE CIRCUIT BOARD.

PARTS ORIENTATION

In all references throughout the instructions, the convention used is that the gold edge connector is the <u>bottom</u> of the board. Orientation identification is molded into the plastic of the sockets and is illustrated below:

This orientation mark identifies where pin #1 of the integrated circuit is to be positioned when inserted into the socket. The socket should be inserted in the board so that the orientation mark is in the <u>lower right</u> hand corner.

Orientation of the transistors, tantulum capacitors, diodes and voltage regulator is specified in the component layout drawing. It is advisable to study this drawing and the 8 x 10 glossy photograph carefully before building the kit. Refer to both during parts installation.

SOLDERING AND SOLDER IRONS

The most desirable soldering iron for complex electronic kits is a constant temperature soldering iron with an element regulated at 650° F. The tip should be fine so that it can be brought in intimate contact with the pads of the circuit board. Both Unger and Weller have excellent products which fit the above requirements.

There are three important soldering requirements for building this kit.

- 1. Do not use an iron that is too cold (less than 600° F) or too hot (more than 750° F).
 - 2. Do not apply the iron to a pad for extended periods.
 - 3. Do not apply excessive amounts of solder.

The proper procedure for soldering components to the circuit board is as follows:

- 1. Bring the iron in contact with both the component lead and the pad.
- 2. Apply a small amount of solder at the point where the iron, component lead, and pad all make contact.
- 3. After the initial application of solder has been accomplished with the solder flowing to the pad and component lead, the heat of the iron will have transferred to both the pad and the lead. Apply a small amount of additional solder to cover the joint between the pad and the lead. DO NOT PILE SOLDER ON THE JOINT! EXCESSIVE HEAT AND SOLDER CAUSE PADS AND LEADS TO LEFT FROM THE CIRCUIT BOARD. EXCESSIVE SOLDER IS THE PRIMARY CAUSE FOR BOARD SHORTS AND BRIDGED CONNECTIONS.

PARTS INSTALLATION

Before installing parts, bend the leads of the resistors, diodes, and tantulum capacitors to the lengths shown in the right hand column. After a series of parts have been installed in the board, bend the leads slightly to hold them in place, solder the leads and trim the excess lead lengths before proceeding to the next step.

Install:

CR1	1N914/482	0-0201 signal diode	Check for orientation!	.5 in.
R1	180Ω ⅓ wa	tt resistor	brown-grey-brown	.5
R2	3.9kΩ	п	orange-white-red	.5
R3	1kΩ	п	brown-black-red	.5
R4, R7	3.3kΩ	u	orange-orange-red	.5
R5, R6	27kΩ	п	red-purple-orange	.5
R8-R13	180Ω	н .	brown-grey-brown	.5
R22	510Ω	II .	green-brown-brown	.5
R14	1.5k Ω	11	brown-green-red	.5
R15	3.3k Ω	n	orange-orange-red	.5
R16, R17	1 k Ω	п	brown-black-red	.5
R18	47k Ω	п	yellow-purple-orange	. 5
R19	4.7k Ω	п	Yellow-purple-red	.5
R23	510Ω	11	green-brown-brown	.5
R24	1 k Ω	16	brown-black-red	.5
R25, R26	510Ω	н	green-brown-brown	.5
CR2, CR3	1N4742 12	v. zener diode	Check for orientation!	.5
C3	.82 µfd t	antulum capacitor	II	.5
C5, C6	2.7 µfd	n .	II	.5
R21, 27, 2	8 240Ω ½	watt resistor	red-yellow-brown	.6
4 Mhz crys	tal		Use a trimmed resistor to secure the crystal t circuit board	

Check for orientation! .6 in. C2, C4 39 ufd tantulum capacitor н Socket 11A 8-pin low profile 11 Sockets 3A, 6A-8A 16-pin low profile Sockets 9A, 10A, 13A 14-pin low profile Sockets 4A, 5A 20-pin low profile Sockets 1B-3B, 5B, 8B, 10B, 12B, 13B п 14-pin low profile а Sockets 4B, 6B, 7B, 9B, 11B 16-pin low profile o. Sockets 1C-3C, 5C, 12C 14-pin low profile 11 Sockets 4C, 6C-8C, 13C 16-pin low profile 11 Sockets 9C-11C 20-pin low profile R20 470Ω 1 watt resistor yellow-purple-brown .8 J2 7-pin right angle header with pins oriented toward the top of the board (away from the 100 pin Check for orientation! edge connector) 50-pin right angle header with J1 pins oriented toward the top of the board (away from the 100 Ħ pin edge connector) Slide switches S1 and S2 After installation, these two switches should be switched in the same direction: to the left if the Disk Jockey is to be used with a full-size floppy disk or to the right if it is to control a 5" minifloppy disk drive. C1 .01 ufd disk capacitor By-pass capacitors (17 each). Values may vary from .01 µfd to $.1 \mu fd.$ LM340.5/7805 5 volt regulator Bend lead, insert and hand tighten the nut and bolt with the bolt running through the bottom of the board through the heat sink and through the regulator. The nut should be hand tightened over the regulator. Solder the leads. If heat sink grease is available, apply a thin film between the board, heat sink, and regualtor. Finally, tighten the nut firmly.

POWER-UP AND SYSTEM CHECK OUT

Power Supply/Voltage Regulator Check Out

Voltage requirements:	(reference to grou	nd - pins 50 and 100)
Pins 1 and 51	no less than 7 volts not more than 10 volts	approx. 700 ma
Pin 2	not less than 13 volts not more than 22 volts	approx. 25 ma
Pin 52	not less than -22 volts not more than -13 volts	approx. 100 ma

Before installing any of the integrated circuits, apply power to pins 1 and 51, pin 2 and pin 52 (ground at pins 50 and 100) as specified above. Perform the following measurements with a volt meter:

(1)	Pin 4 of 11A (741)	-12 volts
(2)	Pin 7 of 11A	+12 volts
(3)	Pin 14 of 13A	+5 volts
(4)	Pin 14 of 13B	+5 volts
(5)	Pin 16 of 13C	+5 volts

If the voltage at any of the check points differs from the required value, return the unit for trouble shooting and repair.

Power-Up Check Out

Install the integrated circuits as per the layout on the board (note: IF the legend at 5A reads 74LS373, change to 74LS273). When inserting these parts, be careful not to bend pins under the package -- a pin which is bent under the integrated circuit may appear to be inserted in the socket. BENT PINS ARE THE MOST COMMON REASON FOR MALFUNCTIONING BOARDS!

After all the parts have been installed, voltages checked, and integrated circuits installed, reconnect the power supplies and power up the board again. This can be done stand-alone or in the mainframe of your computer. Check again the voltages called out in the previous section. If there are any differences from the required values, please return the board.

8"/5" Hardware Selection

Be sure that the two slide switches S1 and S2 are to the left for 8" drives and to the right for 5" drives. BOTH SWITCHES MUST FACE THE SAME DIRECTION -- BOTH TO THE LEFT OR BOTH TO THE RIGHT.

SYSTEM CHECK OUT

Introduction

The Disk Jockey interface occupies 1024 bytes of memory 340:000Q through 343:377Q (E000H through E3FFH). These addresses can be customized for special applications but in the check out procedures which follow, standard addressing will be assumed. Users of customized boards should apply appropriate address offsets.

ROM

- (1) Examine location 340:000Q and verify that it contains 303Q (location E000H should contain C3).
- (2) Examine location 340:003Q and verify that 303Q is also present (E003 should contain C3).
- (3) Examine location 341:377Q and verify that it contains 311Q (E1FF should contain C9).

RAM

- (1) Examine location 342:0000 (E200H).
- (2) Verify that the following values can be deposited at the stated addresses.

0c	tal	Hex	×
Address	Value	Address	Value
342:000	000	E200	00
342:001	001	E201	01
342:002	002	E202	02
342:003	004	E203	04
342:004	010	E204	08
342:005	020	E205	10
342:006	040	E206	20
342:007	100	E207	40
342:010	200	E208	80
342:011	377	E209	FF

I/0

Disk Pause Logic

Enter the followping program in the Disk Jockey RAM:

0	ctal			Hex			
342:000 342:003				3A 01 E3 C3 00 E2	LOOP	LDA JMP	MARK Loop

Start this program at the label LOOP. The system should hang with PREADY low. Turn the computer off and then on in order to generate a \overline{POC} (power-on-reset). Note: many S-100 systems generate a POC signal when the reset switch is pressed and in this case it is not necessary to turn the system off and then on again. In the sequel, the above procedure will be referred to as "generate a \overline{POC} signal."

Enter the following program in the Disk Jockey RAM:

0	ctal	Hex			
	072 000 3 303 000 3	 3A 00 E3 C3 00 E2	L00P	LDA JMP	DATA LOOP

Start the program. This time the system should not hang. Stop the computer. This completes the test of the disk pause logic.

Load Head Command

Generate a \overline{POC} signal. Verify that pins 1 and 15 of IC 6A are at a logic 1 (approx. 4-5 volts). Enter the following program in the Disk Jockey RAM.

Octa1		Hex			
342:000 072 342:003 303	 	3A 03 E3 C3 00 E2	LOOP	LDA JMP	HEAD LOOP

Start this program and after a moment, stop the program. Verify that pins 1 and 15 of IC 6A are now at a logic zero (approx. 0 volts). For the next test, it is necessary that pins 1 and 15 are at a logic zero. This completes the test of the Load Head Command.

Disk Function Register

This test is to be performed just after the Load Head Command test. Before the Disk Function Register is tested, pins 1 and 15 of IC 6A should be at a logic zero level. Enter the following program:

342:003 343:006	062 011 072 002 303 000	343 E 342 E	203 3 206 C	2 09 E2 A 02 E3 3 00 E2		STA JMP	TEST FUNCTION LOOP
343:011	XXX	E	209 X	X	TEST	DATA	

The above program is to be run with various values for DATA deposited in location TEST. With each value of DATA, pins 16, 26, 28, 30, 32, 34 and 36 of J1 are to be probed. For each different value of DATA, exactly one of these pins should be at a logic zero and all the rest at a logic one. The table below details what the different values for DATA should be and what pin of J1 should be and what pin of J1 should be at logic zero for the given value of DATA.

Value Deposit	ed at location TEST	J1 Pin Which Should be at Logic Zero
Octa1	Hex	.
001	01	28
002	02	32
004	04	34
020	10	36
040	20	30
100	40	16
200	80	26

Each time the value of TEST is changed, the short program at the beginning of this section should be run. While it is running, the various pins of J1 detailed above should be probed. After the pins have been checked, the program should be stopped and the next value in the table should be deposited at TEST, etc. This completes the test of the disk function register.

Disk Write Logic

If a logic probe or oscilliscope is available, attach it to pin 38 of J1. Momentarily ground pin 1 of IC 13B. There should be a stream of negative pulses 250 ns wide, four microseconds apart. In any event, pin 1 of IC 3C should be at a logic 1 as well as pin 40 of J1. Enter the following program in the Disk Jockey RAM:

Octal	Hex		
342:000 041 000	E203 36 FF	LOOP LXI	H,DATA
342:003 066 377		MVI	M,377Q (FFH)
342:005 303 003		JMP	LOOP+3

Start the program. While the program is running, verify that pin 40 of J1 is now at a logic Ø. If a logic probe or oscilliscope is available, verify that pin 38 of J1 has a stream of negative pulses 250 ns wide, two microseconds apart. Also verify that pin 1 of IC 3C has a stream of negative pulses 250 ns wide, four microseconds apart. Stop the program. This completes the test of the disk write logic.

Disk Status Register

Generate a \overline{POC} signal. Read the following list of values from location 343:002Q (E302H). Associated with each value is a pin # of J1. In each case, ground only this pin while examining 343:002Q.

Value a	at 343:002Q	Pin Number of J1 to
Octal	Hex	be Grounded
001	01	44
010	08	42
020	10	20
040	20	10
100	40	24
200	80	22

Serial Input Port

Generate a \overline{POC} signal. With a jumper connect J2 pin 1 to 11A pin 4. Examine location 343:002Q with this jumper in place and verify that 343:002 now contains the value 2. Disconnect the jumper. Examine location 343:003. Verify that the value of 343:002 is zero. This completes the test of the serial input port.

Serial Output Port

Generate a \overline{POC} signal. Verify that pin 2 of J2 is at or near -12 volts. Enter the following program in the Disk Jockey RAM:

0	ctal	Hex		
342:003	072 011 342 062 003 343 303 000 342	E200 3A 09 E2 E203 32 03 E3 E206 C3 00 E2		LDA TEST STA SERIAL JMP LOOP
342:011	• •	E209 XX	TEST	DATA

DATA is to have two values: 400 (20H) and zero. Run the program with DATA = \emptyset and verify that pin 2 of J2 is at or near +12 volts. When DATA = 400, pin 2 of Jw should be at or near -12 volts.

This completes the preliminary check out of the Disk Jockey controller. Further check out will require that a drive be connected to the board.

CABLE ORIENTATION

The cable should be so that the ribbon cable leaves the connector toward the back of the computer and away from the component side of the Disk Jockey circuit board. Care should be exercised in fabricating a disk cable. The connectors at the cable ends must be oriented so that pin 2 of J1 matches pin 2 of the drive connector. A cable for Shugart drives is available from Thinker Toys.

DISK DRIVE CHECK OUT

The following tests can be conducted after the appropriate disk drive is connected to the Disk Jockey (via the 50 conductor cable) and proper power is furnished to the drive.

Track Zero Seek

Set up the disk drive so that it is drive #1. Make sure that the head is away from the track zero sensor. Enter the following program into the memory of the host computer.

Octal	Нех		
000:076 210	00 3E 88	MVI	A,210Q
002:072 163 342	02 3A 73 E2	STA	DRIVE
005:072 002 343	05 3A 02 E3	STA	FUNCTN
010:061 150 342	08 31 58 E2	LXI	SP,STACK
013:315 011 340	OB CD O9 EO	CALL	TKZĖRO
016:303 016 000	OE C3 OE OO	JMP	STALL

Insert a write protected diskette into the drive and close the door securely. Start the program. The head should load and move outward until it reaches track $\emptyset\emptyset$. After sixteen revolutions, the head should unload. Do not turn off power to the computer and drive between this test and the next. Do not change the stack pointer.

Seek and Disk Read

In the next test, be sure that a write protected disk which is cleanly initialized is inserted in the drive. Enter the following program:

0cta	:1				
100 102 105 107 112 115 120 123	016 002 315 014 340 016 001 315 017 340 041 000 001 042 164 342 315 025 340 303 123 000	START STOP	MVI CALL MVI CALL LXI SHLD CALL JMP	C,2 TSEEK C,1 SECTOR H,1:000 DMA DISKR STOP	SEEK TRACK 2 SET SECTOR 1 SET DMA TO PAGE 1 READ DATA
Hex 40 42 45 47 4A 4D 50 53	OE 02 CD 0C E0 OE 01 CD OF E0 21 00 01 22 74 E2 CD 15 E0 C3 53 00	START STOP	MVI CALL MVI CALL ŁXI SHLD CALL JMP	C,2 TSEEK C,1 SECTOR H,100 DMA DISKR STOP	

Power-Up and System Check Out

Start the program. The head should load and move to track 2 and unload after sixteen revolutions. Verify that the value 345 is contained in all memory locations between 1:000Q and 1:177Q.

This completes the check out of the Disk Jockey controller.

HARDWARE DESCRIPTION SECTION

Power Supply

The board has a worst case requirement for

```
800 ma of +8V (between 7.5V and 11V)
100 ma of -16V (between -12.5V and -20V)
30 ma of +16V (between +12.5V and +20V)
```

Plus 5 volt power (Vcc) is derived from the +8V supply in the computer backplane by means of the 5 volt 7805/LM340.5 regulator with C2 and the small disk capacitors labeled on the circuit board legend providing power supply bypass capacitance.

The +12V supply is derived from the backplane +16V supply by zener CR3 and R28, filtered by C6.

The -12V supply is derived similarly from the backplane -16V supply by means of CR2, R27, and C5.

Floppy Disk Section System Clock

The basic timing for the DJ is derived from a crystal stabilized 74LS124 oscillating at 4 MHz. This clock signal is divided by two to produce 2 MHz SYS CLK frequency appropriate for a standard floppy at pin 9A-5. This signal is divided by two again to produce the 1 MHz SYS CLK appropriate for a mini-floppy at pin 9A-9. The appropriate frequency is selected by switch SW2 to drive the SYS CLK line. An inverted version of the system clock is generated at pin 10B-10.

Device Address Decoding

The Disk Jockey contains 256 bytes of RAM, 512 bytes of ROM, and I/O registers. The memory address where each of these modules appears is controlled by the MMI6301 256x4 PROM in slot 8C. This PROM monitors the highorder address lines A8-A15. Connections to the chip select pins of SINTA, SINP, and SOUT make sure that the PROM is only active when the address present on the S-100 bus is a valid memory address during either a bus memory read or write cycle.

Take, as an example, the RAM ENBL output of the PROM. Normally high, this output should go low when the address on A8-A15 corresponds to the high order address byte that should select the RAM. The PROM internally has 256 cells, each one selected by a different combination of the high order address bits. All of these cells have 1 stored in them except for the one corresponding to the address where the RAM should be located which has a 0 stored in it. Thus, when the CPU refers to any location within the 256 byte page of memory where the RAM should be located, the cell with the zero in it is selected by A8-A15 and the RAM ENBL output goes low, selecting the RAM (see System RAM below).

The firmware ROM, being 512x8, has two consecutively addressed bits of zero in the address decoding ROM pattern. Thus, the signal ROM ENBL is generated for either 256 byte page of ROM. Address line A8 also goes to these ROMs to allow them to decide which 256 byte page within them is being called for (see System ROM below).

Likewise, the signal $\overline{\text{I/O ENBL}}$ is generated when the CPU refers to memory addresses within the 256 byte page assigned to the I/O device registers.

The fourth output, 8C-9, is programmed to go low if any of the other outputs of the address decoding ROM is low, that is, if the CPU is referring to any address within the DJ.

Internal Data Bus

Data flow in the floppy disk controller is organized around an eightbit wide internal data bus DATAØ - DATA7.

When the CPU is writing to the RAM or an I/O register, the S-100 bus signal MWRITE is high or \overline{WR} is low. The \overline{WRITE} line (5B-11) therefore goes low. Data flows from the S-100 bus data out lines (DOØ - DO7) through the 81LS97 in slot 9C to the internal data bus. Depending on what DJ register is being loaded, other control logic causes the data to be loaded from the internal data bus to the appropriate destination.

When the CPU is reading from the ROM, RAM, or I/O registers, other logic, detailed below, causes the selected data byte to be gated onto the internal data bus. When the line PDBIN goes high, indicating that the CPU is reading and the PROM in slot 8C detects that the CPU is addressing a region of memory assigned to the DJ and therefore grounds pin 8C-9, the signal INPUT ENBL goes low. Since pin 10C-1 is now low and pin 10C-19 high, the 74LS241 in slot 10C gates the data byte from the internal data bus to the S-100 bus data lines (DIØ - D17).

Processor Wait States

Whenever a reference is made to the I/O registers of the Disk Jockey, the line PREADY is grounded to add two wait states to the CPU's cycle. The disk data transfer operations can require up to two processor cycles before they can obtain synchronization and ground PREADY to indicate that they require wait states. Two wait cycles are therefore added to every address reference within the Disk Jockey I/O registers to make sure the CPU does not leave the wait state before the disk data transfer control logic can ground PREADY.

The two flip-flops in 1B-9 and 1B-5 act as a shift register to supply the two machine cycle delay required. These flip-flops are always cleared by PSYNC, connected in an inverted version to 1B-13 and 1B-1. 1B-5 will then stay low until the second negative edge of PHI-2 after PSYNC goes low. At this time, the high level connected to 1B-12 will appear at 1B-5 having been shifted through the two flip-flops. During this interval when

1B-5 is low, 3B-11 will be held low, so 3B-8 will be held high. This will happen whether or not the CPU is referencing the Disk Jockey. However, only if the CPU is referencing Disk Jockey I/O registers will the signal I/O ENBL be low, enabling the 74367 in slot 3Å to drive PREADY low. If necessary, the other inputs to gate 3B (9 and 10) are driven low by the disk data transfer control logic before 3B-11 goes high to add additional wait states (see below).

Due to the high speed of the disk, the use of processor wait states is essential for synchronization. If a flag bit were provided which the processor could test to determine if the floppy disk was ready for data transfer and branch accordingly, the execution of this loop would take up too much time to allow transfer of data from memory to disk. The use of processor wait states, suggested by Eugene Fisher of Lawrence Livermore Laboratory, synchronizes the processor instruction execution to the disk data transfer and ensures sufficient time for each byte of data to be transferred from device registers to memory before the next byte must be transferred.

System RAM

The Disk Jockey contains two 2112 256 word x 4 bit RAMs located in slots 6B and 7B. When the CPU during a read or write operation puts the address assigned to this RAM onto the high order S-100 bus address lines (A8-A15), RAM ENBL (8C-11) goes low. This signal activates the chip select pin (13) of both RAMs. Whether the RAM does a read or a write operation is controlled by the WRITE line which connects to the R/W pin (14) of both RAMs. During a write, data from the CPU flows from the internal data bus into the bidirectional I/O pins of the RAMs. During a read, data from the RAMs flows out onto the internal data bus and from there onto the S-100 bus.

System ROM

The Disk Jockey is provided with two MMI6306 512 word x 4 bit ROMs in slots 6C and 7C. These ROMs contain 512 bytes of disk utility software. When the address of the area of memory assigned to the ROM appears on the S-100 bus address lines (A8-A15), the signal $\overline{\text{ROM}}$ ENBL (8C-12) goes low. This enables the two ROMs by grounding their $\overline{\text{CS}}$ pin (13). The nine low order S-100 bus address bits (AØ - A8) specify to the ROMs which word is to be read. The ROM drives the byte of software stored at this location onto the internal data bus.

I/O Registers

When the address decoding PROM in slot 8C detects an S-100 bus memory reference to an address within the 256 byte page assigned to the I/O registers, it grounds 8C-10 (I/O ENBL). By grounding 4C-2 and 4C-14, this activates the 74LS155 decoder in slot 4C whose function is to decide which

of the eight I/O registers is being referred to. The four readable registers and the four writeable registers are each assigned a section of the decoder. The two low order S-100 bus address bits select one of the four registers in each group. If the CPU is engaged in a write operation, the line WRITE will be low and the selected output of the "write section" decoder will go low (4C-9, 10, 11 or 12). If the CPU is engaged in a read operation, the line INPUT ENBL will be high, and the selected output of the read section of the decoder will go low (4C-4, 5, 6, or 7). Note that the definitions of WRITE and INPUT ENBL are such that either a write register or a read register will be selected, never both.

The uses of the READ DATA, WRITE DATA, READ MARK, and WRITE MARK signals are discussed below under the disk data flow section.

The READ STATUS signal, when low, enables the 81LS97 octal buffer in slot 4A, allowing the status byte (six bits of status from the floppy disk drive cable, the TTL level serial input, and the status of the head flip-flop) to flow onto the internal data bus.

The LOAD HEAD signal, when low, resets the 74LS161 hex counter in slot 4B and sets the 74LS74 flip-flop in slot 10A. This flip-flop drives the LOAD HEAD signal to the disk drives and enables the 74368 driver in slot 6A. The 74368 driver at 6A drives the disk select lines over the flat cable to the disk drive(s).

The WRITE FNCTN signal, when low, clocks the contents of the internal data bus into the 74LS273 octal register in slot 5A. Seven of the data bits set control lines to the floppy disk and are driven onto the flat cable from the 74LS273 outputs by the 74368 drivers in slots 3A and 6A. Bit 3 of the octal latch controls whether the read/write head will automatically unload after sixteen revolutions of the disk. The latch output at 5A-9 is connected to 4B-7, the enable input of a 74LS161 counter. This counter is reset to zero when the heads are loaded. The flip-flop in slot 10A is also set to 1 enabling the 74368 to drive the current setting of the drive select lines and the load head signal onto the flat cable. If bit 3 specifies that automatic head unloading is not to take place, the counter is not enabled, and no counting takes place. If the bit enables the counter, then every time the index point of the disk is passed (assuming the correct setting of SW1), the pulse on 4B-2 will cause the counter to count up by 1. When the counter finally counts to 15, the carry output at 4B-15 will go high. This will clock the 74LS74 at 10A-11, which will load a Ø from the data input 10A-12. The head load flip-flop will therefore be reset, disabling the driver in slot 6A, deselecting all drives and causing the LOAD HEAD signal to become inactive. Note that there is no explicit command to unload the heads; it can only be done by this timeout mechanism.

The WRITE SERIAL signal clocks the 74LS74 serial out flip-flop at 10A-3 whose data input is bit 5 of the internal data bus. The flip-flop output 10A-5 drives the translator from TTL to RS232 levels for the serial output; 10A-6 controls the 20 ma TTY driver. These level translation circuits are discussed in more detail below.

Serial Communications Level Translators

Serial Input:

A terminal device is connected to either the RS-232 or the TTY input pin.

If the RS-232 input pin is at a "SPACE" level (greater than 3 and less than 12 volts), transistor Q2 is cut off, so transistor Q1 is cut off, and the SERIAL INPUT line is pulled up to Vcc through R19 and the string of R5, R6, and R20. When the Read Status Register is read, the SERIAL INPUT line is gated onto bit 1 of the internal data bus through the 81LS97 in slot 4A, and a Ø appears in bit 1 of the word read. If the RS-232 input pin is at the "MARK" level (less than -3 but greater than -12 volts), current flows through transistor Q2, causing Q1 to pull the SERIAL INPUT line to ground. A 1 will therefore appear in bit 1 of the register when it is read.

If current is allowed to flow between the TTY+ input, through the terminal device, and back into the TTY- input (-12 volt return), representing the "MARK" condition, the voltage level at the SERIAL INPUT line is pulled down to approximately ground through R5 and R6. If no current is flowing ("SPACE"), the voltage level of the SERIAL INPUT line is pulled up to Vcc through R20, R5, and R6 (as well as R19).

Serial Output:

When the program writes a byte to the Write Serial Register, the signal WRITE SERIAL at pin 10A-3 clocks the value of bit 5 of the internal data bus into flip-flop 10A-5. The Q output 10A-5 is compared to a reference value of 1.6 volts developed by R15 and R14 at the 741 op amp. The output of the op amp will be saturated at approximately -12 volts for a "MARK" and approximately +12 volts for a "SPACE". This output drives the RS-232 level output through resistor R4. Capacitor C1 is required to lengthen the rise and fall time of the RS-232 output, as required by specifications.

Meanwhile, the $\overline{\mathbb{Q}}$ output 10A-6 is driving transistor Q3 which acts as a current source for the TTY 20 ma output. If 10A-5 is set corresponding to the "MARK" condition, 10A-6 is low, current flows from Vcc, out the base of Q3, through R16, and into 10A-6. Q3 is turned on and approximately 20 ma of current will flow from Vcc, out the TTY+ output, in the TTY- output, and through R21 to ground. If the flip-flop is reset, corresponding to the "SPACE" condition, 10A-6 is sufficiently close to Vcc that the voltage at the base of Q3 will not turn on Q3. No current will flow out the TTY+ output in this condition.

When the power on clear signal \overline{POC} goes low, the flip-flop 10A-5 is set, putting both interface lines into the "MARK" state. This is the appropriate quiescent state for these lines.

Disk Data Flow

Before discussing the Write Data, Write Mark, Read Data, and Read Mark operations, it is necessary to consider the way that data is organized on the disk.

Each data bit is recorded on a small region of magnetic material that moves under the read/write head. This region is known as a bit cell. Since the disk is rotating under the read/write head at a fixed velocity, the bit cell also corresponds to a region of time.

Bit cells can contain one or two pulses, depending on whether the data recorded is a zero or a one. Bit cells start with a pulse known as the clock pulse. A fixed amount of time later will be a data pulse if the bit cell has a one recorded in it, or no data pulse if a zero is recorded.

A byte of data is recorded on the disk at eight consecutive bit cells. The most significant bit of the byte is the first bit recorded.

For more effective use of the disk surface, information stored on a track of the disk is traditionally divided into fixed size blocks of data called sectors. Sectors consist of a header field identifying track and sector number, a data area, and a checksum. A number of sectors follow each other in series around the circumference of the track. Gaps between the sectors and between the header field and the data area within each sector allow time for the write head to turn on or off so that one of these fields can be written without erasing the following field or sector.

There are two techniques in use to identify where on the track sectors begin. These techniques are referred to as hard sectoring and soft sectoring. Each of them will be discussed below.

Every diskette has a hole punched in it hear the hub. This hole, sensed by a phototransistor, identifies the beginning of the track. In the hard sectoring format, a similar hole is punched for each sector, indicating that the sector is about to begin.

Given the small physical size of information on the disk, it is impossible to accurately align a mechanical hole with the data recorded on the disk. The hole therefore only indicates that the track or sector is about to begin. An address mark is recorded just before the data on the disk. After the hole is sensed, it is necessary to wait until this recorded mark is detected, indicating the precise start of a sector. The soft sectored format relies solely on these recorded marks, making sector holes unnecessary.

There are actually four such marks. The Index Mark is used to indicate the beginning of each track. The Sector Mark is used to indicate the beginning of the header field of each sector. The Data Mark is used to indicate the start of the data area of each sector. The Deleted Data Mark can be used in place of the Data Mark to indicate an alternate type of data; in practice it is rarely used. The meaning of this last mark is up to the user.

What makes a mark immediately recognizable is that it is missing certain clock pulses. If we encode the clock bits from each bit cell into a byte of hexadecimal in the same manner as the data bits with a present clock pulse being considered a 1 and an absent clock pulse a Ø, we get the following pattern for the four marks:

	Data Bits	Clock Bits
Index Mark	FC	D7
Sector Mark	FE	С7
Data Mark	FB	C7
Deleted Data Mark	F8	C7

The Index Mark and Sector Marks are normally written onto a track only when the disk is initialized. The Deleted Data Mark is not used in standard format disks. The Data Mark is written each time the data area of a sector is updated.

The marks also serve another function: synchronization. Each separately writeable field on the disk must have a gap before and after it in which no meaningful information is recorded to allow time for the write head to turn on before writing and turn off after writing. These gaps are invariably filled with random magnetizations that would make any system lose track of where the byte divisions in the serially recorded bit stream are (to say nothing of which pulses are clock pulses and which pulses are data pulses). By a technique discussed below, the marks, which precede every field written on the disk, can be used to decide unambiguously where the clock pulse of the first bit cell of the first byte of the field begins.

The four disk data transfer operations have certain hardware in common. The basic timing for each bit cell is generated by four D flip-flops in slots 12B and 13B. They are clocked by SYS CLK connected to pins 11 and 3. Arranged as a shift register, the flip-flops cycle through a sequence of eight states before repeating themselves as illustrated in Figure 1.

The first state represents the time period in which a clock pulse, if any, will occur. There are then three states which correspond to the interval of three clock periods before the data pulse. Then there is the state corresponding to the data pulse, followed by three more states representing the time from the data pulse to the end of the bit cell.

The counter in slot 13C is used to count the eight bit cells that comprise one byte. It is clocked by \overline{C} CLK attached to pin 2, so that it changes state at the start of each bit cell. Although it is a divide by sixteen counter, it is used as a divide by eight counter (no connection is made to Q-D). In certain circumstances, it is necessary to set this counter to a certain state for synchronization purposes. This is accomplished by the signal \overline{LD} , connected to pin 9, which loads the value 1100 binary into the counter on the next clock pulse when it is low. This value comes from the data in pins 6, 5, 4 and 3 which are strapped high or low as appropriate. Pin 6 could actually have been connected to Vcc or ground since Q_D is not used.

Figure 2 shows the states of this counter, starting with a load operation.

See Timing Diagram under Schematics

Two additional signals are derived from this counter: \overline{EOC} (End of Character) is generated by gate 13A-6 during the C CLK period of the bit cell where the low order three bits of the counter have the binary value 110; \overline{EOW} (End of Word) is generated by 13A-8 during the A CLK period of the following bit cell.

Let us now consider the actual data transfer operation.

WRITE DATA: When the WRITE DATA line goes low at pin 3C-9, 3C-8 goes high, clocking a Ø into flip-flop 2B-5. 3B-8 goes high, driving PREADY low. This halts the CPU until proper synchronization is obtained and the byte of data can be accepted from the S-100 data bus into the Disk Jockey.

1/O ENBL connected to 3A-15 provides insurance that the board can only halt the CPU when one of its registers is actually selected.

Since the bit cell and byte counters will be at an unpredictable point when the Write Data command is issued, it is necessary to wait until they come to the point indicating the start of a byte. This is done by flipflop 1C-5. When the byte counter reaches the EOC state and A CLK goes high, the 1 present on 2B-6 since the write command was issued (WRITE SYNC) is loaded onto 1C-5, and WRITE GATE goes high. In the next bit cell, EOW goes low, setting 2B-5, so WRITE SYNC goes low. Thus, on the positive edge of A CLK in the next bit cell where EOC is high, WRITE GATE will go low. WRITE GATE will therefore be high for eight consecutive bit cells between two EOCs when a Write Data is being performed.

When WRITE GATE goes high, it drives 10B-1 low. As a result, 11B-9 and 11B-3 stay high. This ensures that the bit cell counter cycles through its normal sequence of states without being set or reset.

Shift register 11C converts the data to be written from its parallel format on the internal data bus to the serial format that must be sent to the disk. It is clocked by the A CLK towards the end of each bit cell to set up the data for the next bit cell. During the first EOC after WRITE SYNC goes high, it is necessary to load the parallel byte of data written by the CPU into the shift register. Pin 5C-8 goes high during this time, driving 11C-9 high. This indicates the load function to the 74LS299 in slot 11-C. Accordingly, when the A CLK goes high, the byte of data from the internal data bus is loaded into the shift register from its bidirectional data pins. These pins are inputs since READ ATTN is low, so 11C-2 is high.

On the next bit cell, the $\overline{\text{EOW}}$ signal becomes active. This resets WRITE SYNC. As a result, 2B-5 goes high, 3B-8 goes low, and PREADY goes high. The CPU is now free to leave the wait state; appropriate since the byte of data to be written to disk has been successfully loaded into shift register 11C.

Output Q-H of 11-C now has the first bit to be written to the disk. Since WRITE SYNC is no longer high, 11C-9 is low, conditioning the 74LS299 for a right shift operation. In subsequent bit cells, EOC will not be high, so 11C-9 will stay low, even if WRITE SYNC comes high again as a result of a subsequent write operation before the termination of the current one. Since 11C is now a shift register, when A CLK goes high at the end of each bit cell of the byte, the shift register shifts one position, presenting the next bit to be written serially to the disk at output Q-H. Gate 12C-8 combines the data bit with clocks to produce the WRITE DATA signal. This signal goes low during the appropriate time of a bit cell if a 1 is to be written to disk; it stays high if a zero is to be written.

Meanwhile, WRITE CLOCK pulses are being generated for each bit cell by 12C-6 at the appropriate time in each bit cell. 12C-2 is always a one since this is not a Write Mark operation and shift register 9B is therefore conditioned to shift in all ones from its serial input 9B-10. A clock pulse is therefore generated for every bit cell.

WRITE DATA and WRITE CLOCK are or'ed together by 3C-3 and driven onto the WRITE DISK DATA line to the write head of the disk by the 74368 in slot 3A. The WRITE GATE signal is also driven onto this cable to indicate to the floppy disk drive the time period during which the WRITE DISK DATA is meaningful.

When the last data bit has been shifted out of 11C, the byte has been successfully written to disk. If the CPU has issued another Write Data command while the first byte was being shifted out, this new byte will be loaded into 11C during the EOC period and the process will repeat itself. WRITE GATE will stay high since WRITE SYNC will have been set again. If no further data bytes were written, WRITE SYNC will be low, and WRITE GATE will go low during EOC on the positive edge of A CLK (the end of the last bit cell of the character).

A Write Mark operation is identical to a Write Data operation except that the byte of data is written with certain clock pulses missing. Shift register 9B is responsible for generating the correct pattern of missing clock pulses during the eight bit cells. This register is clocked by A CLK. When $\overline{\text{WRITE}}$ MARK is low, during the EOC period when data is being loaded into the data shift register, 8B-3 goes low and the pattern present at the data inputs to 9B is loaded into it. This pattern is then shifted out during the next eight bit cells analagously to the data bits, with a one providing a clock pulse for the bit cell and a zero providing a missing clock pulse.

The pattern of missing clocks can be a hexadecimal C7 or D7, depending on which mark is to be written. Gates 10B-13 and 5B-6 decode from the data pattern of the mark which clock pattern is appropriate and provide the appropriate level to 9B-3 for parallel loading into the shift register.

READ MARK: The Read Mark operation performs two functions: (1) to synchronize the data separator logic (flip-flops 12B and 13B) and the byte counter (13C); and, (2) to read the data pattern of the mark so that software can determine which of the four types of mark it is (see above).

When the $\overline{\text{READ MARK}}$ signal goes low, 5C-6 goes high, clocking a one into 2C-5 ($\overline{\text{LD}}$ and AØ are both high at this time, so 2C-1 is high). The output of OR gate 8B-6 is therefore high, putting the set input 2B-10 into its inactive state.

Meanwhile READ MARK going low causes 3C-6 (READ ATTN) to go high, clocking a Ø into 2B-9. This causes 3B-9 to go low, 3B-8 to go high, so PREADY goes low, putting the CPU into a wait state until the mark is detected.

DISK DATA -- the serial stream of clock and data pulses from the disk enters gate 10B-3 from the disk drive cable. Since this is a read operation, WRITE GATE will be low and DISK DATA will appear at 10B-1. A pulse on this line will be routed to DATA or to reset (RESET) flip-flops 12-B and 13-B

depending on whether A CLK is high or low respectively. A CLK therefore defines the time window in which a pulse from DISK DATA is interpreted as a clock or a data pulse. If A CLK is high, the pulse is a data pulse; if it is low, it is a clock pulse. Nominally, data pulses should occur when the flip-flops are all set and clock pulses should occur when the flip-flops are all reset. The pulse from the disk by setting the flip-flops to all zeros will force the flip-flops to this exact synchronization if the pulse occurs any time within the A CLK window.

Register 7A wired to act like a shift register has the function of counting missing clock pulses. It is clocked on the positive edge of A' CLK just after the clock pulse nominally occurs. Since the head is loaded, HEAD is high and ONE comes high. The register is reset by RESET which goes low whenever a clock pulse is detected from the disk. In normal operation, ONE will go high between clock pulses, but be reset by each clock pulse. If a clock pulse is missing, however, the register will not be reset and TWO will come high just after the first missing clock pulse should have occurred. If three missing clock pulses occur in a row, FOUR will come high. As will be discussed in a moment, this condition occurs when an address mark with three missing clock pulses is read. Four missing clock pulses cause OUT OF PHASE to go high. There are only two possible explanations of four missing clock pulses: garbage is being read off the disk or the system has its timing backwards and is interpreting data pulses as clock pulses and vice versa. If the missing "clock pulses" is actually supposed to be a data pulse, the correct response is to reset the system timing and interpret the next pulse as a clock pulse.

As can be seen on the prints, OUT OF PHASE going high forces 10B-4 low and 11B-14 high, routing the next pulse from DISK DATA to the RESET line, i.e., interpreting it as a clock pulse. While this may not help much if garbage is being read off the disk, this technique guarantees that the system will be correctly distinguishing clock and data pulses within several bit cells after the start of a field of data zeros on the disk (since the missing data pulses of the zeros will continually trigger the missing clock detection logic if the system is out of sync). The recording standards for floppy disks require that six bytes of zeros be written before any address mark for this reason. This function is automatically carried out by Disk Jockey firmware when standard subroutines are used to write data sectors. As a result, one can guarantee that clock and data pulses are being correctly interpreted by the start of the address mark. Now it only remains to figure out which bit cell starts a byte of recorded data.

The System Timing Diagram shows the response of the circuitry to the last bit cell of the zero field followed by an address mark. The first two bit cells of the address mark are ones with normal clock pulses. The next three bit cells have missing clock pulses, denoted by dotted lines. As a result, the line FOUR will be high after the third missing clock pulse. On the next positive edge of \overline{C} CLK, \overline{LD} will be low, since READ MARK, FOUR and \overline{EOW} will be high (\overline{EOW} must be high since A' CLK is low at the positive edge of \overline{C} CLK). Byte counter 13C will therefore load the value 1100 from its data inputs (13C-6, 5, 4, and 3), establishing synchronization of bit

cells, reaching the value 111 at the start of the next byte. As can be seen, \overline{EOC} will be low during the last bit cell of the byte, and \overline{EOW} will be low during the first bit cell of the next byte, which is the correct synchronization for these signals.

The $\overline{\text{LD}}$ signal also goes to 5B-1. Since AØ is high (the address of the Read Mark Register being odd), 2C-1 goes low, 2C-5 goes low, and 8B-4 goes low. This indicates that synchronization of the byte counter has been obtained. As a result, the next time $\overline{\text{EOW}}$ goes low, slightly into the first bit cell of the following byte, 8B-6 goes low, 2B-10 goes low, and 2B-9 goes high. The CPU is therefore allowed to leave the wait state, terminating the READ MARK operation.

Parallel to this activity, the data bits of the mark have been collected. Flip-flop 2C-9 is reset by B CLK during the clock period of each bit cell. It will then be set by DATA if there is a pulse during the data pulse period of the bit cell (corresponding to a recorded 1 bit) or 2C-9 will stay at Ø if there is no DATA pulse corresponding to a recorded Ø.

The READ DATA signal generated by this flip-flop goes to the serial input of shift register 11C. 11C is conditioned to act as a shift register since WRITE SYNC is low at 3B-5 so 11C-19 must be \underline{low} . The data bit is shifted into the register on the positive edge of \overline{A} CLK, shortly after it is detected.

READ ATTN being high at 11B-13 puts 11C-2 low, enabling the byte stored in the shift register onto the internal data bus. Thus, when the CPU is released from the wait state at the end of a read operation, the last eight data bits collected will be available on the S-100 bus data in lines. The CPU will read this byte before the next positive edge of A CLK reads the next bit into the register.

READ DATA: Read Data is simpler than Read Mark since all counters are already running in sync. The normal state of shift register 11C is to collect the read data coming in from the disk. When the READ DATA line goes low at 3C-5, READ ATTN goes high, enabling the contents of 11C onto the internal data bus and loading flip-flop 2B-9 with a Ø, putting the CPU into the wait state. When the byte currently being read has been completely loaded into shift register 11C, EOW goes low. Since the Read Data Register is at an even address, AØ is low, so 2C-1 is low, 2C-5 is low, and 8B-4 is low. 8B-6 therefore goes low, setting 2B-9 to 1 and the CPU leaves the wait state and reads in the collected byte.

Note that while both READ MARK and READ DATA release the CPU from the wait state slightly into the byte after the one being read, this still leaves adequate time for the CPU to read the current byte and issue a new READ DATA instruction before the new data byte has been assembled from the disk (EOW going low). Since read data is continuously collected in shift register 11C, it is therefore no problem to read a series of bytes from the disk.

		E 0164 3				•						
	340:000	ORIGIN E		HUGART FIRMWA D:0000	KE.		0 3	315 305 340 302 055 340		CALL JNZ RET	DREAD 8LOOP	READ IN BOOTSTRAP RELOAD ON ERROR BRANCH TO BOOTSTRAP
	343:000 343:001 343:002 343:003	READDATA READMARK STATUS E	A EQUI OI K EQUI RE EQUI REA	RIGIN+300H EADDATA+1		07 07	4 1	166 166		HLT HLT		TEST INSTRUCTION TEST INSTRUCTION
	343:000 343:001	WRITEDAT WRITEMAR				07	6 (000:002		DS	2	
	343:002 343:003 342:166	DISKENCT SERIAL E	EQU 5	TATUS		10- 10	2 0	006 200 052 160 342	INPUT	MVI LHLD	B.200Q SCON	INITIALIZE BIT COUNT GET SPEED CONSTANT
	342:171 342:164	SECREG E	QU 801 EQU BUI	FFER+3 FFER-2		10 11: 11	0 (021 002 343 032 037	WAIT	LXI LDAX RAR	D,STATUS D	INITIALIZE STATUS REG ADDR GET STATUS AND TEST
	342:175 342:375		A EQU DA	ATAMARK+200Q		11 11	12 (13 ()37 332 110 340		RAR JC	WAIT	FOR A START BIT
	342:163 342:160	DRIVE E SCON E	QU BU			11 12 12		315 205 340)32)37		CALL LDAX RAR	XLOOP D	WAIT ¼ A BIT TIME GET STATUS AND TEST
	000:010 000:024	TZERO E MOVIN E	QU 240	j		12 12	3 ()37 332 100 340		rar JĈ	INPUT	THAT THE START BIT IS STILL THERE
	000:020 000:020 000:010	MVOUT E STEP E TCONST E	QU 200)			2 (DATAL	CALL LDAX	SDELAY D	WAIT ONE BIT TIME GET
	000:010 000:143 000:043		QU 14:				33 (34 (rar Rar Mov	A,B	THE SERIAL DATA IN B
		501100	.40 00			13 13	36 ()37 107		RAR MOV	B,A	ROTATE SAVE
003	303 033 340 303 100 340	DBOOT TERMIN	JMP JMP	BOOT INPUT		14	10 3	322 127 340 315 201 340		JNC CALL	DATAL SDELAY	TEST FOR BYTE COLLECTED WAIT ONE BIT TIME
011	303 155 340 303 252 341	TRMOUT TKZERÓ	JMP JMP	OUTPUT HOME			1 1			CALL MOV	SDELAY A,B	WAIT SECOND BIT TIME CHARACTER TO ACC
014 017 022	303 204 341 303 047 341 303 077 341	TRKSET SECTOR DMA	JMP JMP	SEEK SETSEC SETDMA			52 : 54 :	346 177 311		AN I RET	177Q	TRIM PARITY RETURN
025 030	303 305 340 303 214 340	READ WRITE	JMP JMP	DREAD DWRITE				021 003 343	OUTPUT	LXI	D,SERIAL	INITIALIZE SERIAL O/P REG
	061 166 342 041 041 000	BOOT	LXI LXI	SP.BUFFER H.410	INITIALIZE STACK POINTER SERIAL CONSTANT	16		016 013	0.000	ADD MV1	A C,11	APPEND START BIT TOTAL BIT COUNT
041	076 210 001 200 000		MVI FXI	A,2100 B,2000	DRIVE A SELECT CONSTANT BOOTSTRAP LOAD ADDRESS		3 (34 (O LOOP	STC RAR MOV	B,A	SET REST BITS ROTATE SAVE
046 047	305 365		PUSH PUSH	B PSW	INITIALIZE THE	16	66 2 67 (237		SBB STAX	A D	SEND PRESENT BIT TO OUTPUT
051	345 305		PUSH.	8 H	SYSTEM FOR BOOTSTRAP		73 1	315 201 340 170 ·		CALL MOV	SDELAY A,B	DELAY ONE BIT TIME GET NEXT BIT
055	062 002 343 315 252 341	BLOOP	STA CALL	STATUS HOME	LOAD MOVE HEAD TO TRACK ZERO	17 17	75 3	015 302 163 340		DCR JNZ	C OLOOP	DECREMENT BIT COUNT OUTPUT NEXT BIT
	018 001 315 047 341		MVI Call	C,1 SETSEC	INITIALIZE SECTOR	20	JU .	311		RET		

	35 016 001 17 171 30 346 037 32 310 33 10 34 306 345 36 237 37 300 30 041 171 342 31 161 31 163 31	SETSEC	MYI MOY ANI RZ MOY ADI SBB RNZ LXI MOY CALL MOY INX MOY INX MYI XRA RET	C,1 A,C 37Q C,A 345Q A H,SECREG M,C SÉTCRG M,C H M,B H M,OFBH A	SECTOR NUMBER ONE GET SECTOR NUMBER TEST IF LESS THAN ONE SAYE TRIMMED VALUE TEST FOR GREATER THAN 25 SAYE SECTOR VALUE & COMPUTE HEADER CHECK SUM SAYE CHECK SUM AT END OF SECTOR HEADER & ALSO WRITE DATA MARK	154 156 157 160 161 162 164 165 166 167	172 346 34 253 107 172 017 346 36 251 117 043 321 172 274 330 302 12 173 275 330	7 341		MOV MOVI ANI XRA MOV RRC ANIA MOV RC ANIA POV CMP CMP CMP CMP CMP	C,A A,D OEOH E B,A A,D OFOH C C,A H CRECH+3 A,E L	
	77 041 000 342 12 174 13 270 14 310 15 074 16 270 17 310 10 140 11 151 12 042 164 342 13 11	SETDMA	LXI MOV CMP RZ INR CMP RZ MOV MOV SHLD RET	H.ORIGIN+2: A.H B A B H.B L.C DMAADDR	The state of the s	204 207 210 211 213 214 216	315 35 300 171 346 17	5 341 7 3 7 342	SEEK SLOOP	CALL RNZ MOV ANI MOV ADI SBB RNZ LXI MOV	CRECH+3 LOHEAD A.C 1770 C.A 179 A H,BUFFER+1 A.M	SELECT DRIVE, LOAD HEAD A TEST READY GET TRACK VALUE TRIM & SAVE TEST FOR GREATER THAN 76
1 1 1 1 1 1 1 1	16 041 166 342 21 021 172 342 20 001 377 377 27 325 30 176 31 251 32 127 33 017 34 017 35 017 36 017 37 346 017	SETCRC CRECH	LXI LXI PUSH MOV XRA MOV RRC RRC RRC ANI	H,BUFFER D,SECREG+1 B,-1 D A,M C O,A	Self () Sel	224 225 230 233 235 236 241 243 244	271 312 33 332 24 006 026 065 303 246 006 026 064 315 323 303 223	7 341 1 341 0 4 341 4 341	MOVEIN	CMP JZ JC MVI DCR JMP MVI INR CALL JMP	C C DELAY-2 MOVEIN B.MYOUT M MOVEIN+3 B.MOVIN M TDELAY SLOOP	GET CURRENT TRACK COMPARE WITH NEW TRACK HEAD SETTLE DELAY NEW TRK GREATER THAN OLD TRK NEW TRK LESS THAN OLD TRK STEP DISK HEAD
] 1 1 1 1	11 252 42 137 43 017 44 017 46 017 46 127 47 346 037 51 250		XRA MOV RRC RRC RRC MOV ANI XRA	D,A 1FH B		255 256 260 263 264 266 271	315 359 300 006 024 315 322 032 346 010 302 256 006 020 315 322	1 2 341) 5 341	HLOOP STEPO	CALL RNZ MVI CALL LDAX ANI JNZ MVI CALL	LOHEAD B,MOYIN TOELAY D TZERO HLOOP B,MYOUT TDELAY	START DRIVE, LOAD HEAD & TEST FOR READY STEP TOWARD TRACK ZERO

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340:201	052 160 342	SDELAY	LHLD	SCON	INITIALIZE SERIAL DELAY CONSTANT DECREMENT COUNT TEST FOR COUNT EQUAL TO ZERO
	051		DAD	Н	SERIAL DELAY CONSTANT
205	053	XLOOP	OCX	Н	DECREMENT COUNT
206	174		MOV	A,H	TEST FOR
207	265		ORA	L	COUNT
210	302 205 340		JNZ	XLOOP	EQUAL TO ZERO
213	311		RET		
213	J11				
214	315 355 341	OWRITE	CALL	LDHEAD	LOAD THE HEAD
217	032	D	LDAX	Ď	CET CTATUS
				7	TEST FOR
222	300		QN7	•	WRITE PROTECT
223	052 164 242		LMLD	OMAADDR	GET DMA ADDRESS
223	027 104 247		LITE	D.DATAMARK	
226	021 1/3 346		DIICH	0	SAVE POINTER
231	323		CALL	ACCD	LOAD DISK DATA BUFFER
232	315 363 340		DACE	AFER U	DECOVER POINTER
235	341		CALL	COECH	CALCULATE CRC WORD
236	315 124 341		HOU	UNEUN	STORE CHO HOND
241	161		MOA	M s C	COC
242	043		11/1/	H	-UNOR
243	346 001 300 052 164 342 021 175 342 325 315 363 340 341 315 124 341 161 043 160 016 007 315 377 340 300 006 012		MUY	m,e	SAVE POINTER LOAD DISK DATA BUFFER RECOVER POINTER CALCULATE CRC WORD STORE -CRC -WORD NUMBER OF HEADER BYTES+1 FIND RIGHT SECTOR RETURN IF DISK ERROR DELAY GAP DUMMY READ WAIT FOR 10 BYTES
244	016 007		MVI	U,/	FIND DICHT SCOTOS
246	315 377 340		CALL	KZECI	PETHON IS DISK COOOD
251	300		RN Z	- 10	KETOKA IF DISK ENVOY
252	006 012 032		MVI	B'10	DELAT GAP
424	UJE	1 000	E B1111	Ü	DUMMI KEAU
255	005		ULK	В	WALL FOR TO DITE?
256	005 302 254 340		JNZ	FLOOP	BEFORE WRITING
261			XRA	A	PREPARE IU WRITE ZEROS
262	022	ZLOOP	STAX	A D C	DUMMY READ WAIT FOR 10 BYTES BEFORE WRITING PREPARE TO WRITE ZEROS WRITE SIX BYTES
263	015		DCR	C	SIX BYTES OF ZEROS
264	302 262 340		JNZ	ZLOOP	OF ZEROS
267	023		INX	Ð	
270	015 302 262 340 023 176 022 033 043 176 022		MOV	A,M	GET THE DATA MARK & WRITE ADJUST BACK TO DATA ADJUST POINTER GET DATA BYTE WRITE IT ON THE DISK MOVE POINTER AHEAD & TEST FOR LAST BYTE PREPARE TO WRITE A ZERO WRITE FINAL ZERO
271	022		STAX	D	DATA MARK & WRITE
272	033		DCX	D	ADJUST BACK TO DATA
273	043		INX	Н	ADJUST POINTER
274	176	WLOOP	VQM	A,M	GET DATA BYTE
275	022		STAX	D .	WRITE IT ON THE DISK
276	064		IND	L	MOVE POINTER AHEAD
277			JNZ	WLOOP	& TEST FOR LAST BYTE
302	257		XRA	A	PREPARE TO WRITE A ZERO
303			STAX	D	WRITE FINAL ZERO
304			RET		
304	J				
305	315 377 340	DREAD	CALL	RSECT	FIND CORRECT SECTOR RETURN IF DISK ERROR GAP WIDTH
310			RNZ		RETURN IF DISK ERROR
311			MVI	B,14 D	GAP WIDTH
313		RWA1T	LOAX	D	READ PAST
313			•		

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340:314	005			DCR	В	CRAP IN
	202 212	240		JNZ	PWA! T	THE GAP
315	302 313 072 001	340		1.04	OFARMARY	CRAP IN THE GAP WAIT FOR A MARK LOOK FOR ANOTHER
320		343		LUA	M DREAD L D M.A	LOOK FOR ANOTHER
323	276			CMP	M	SECTOR IF NO MARK ADVANCE POINTER GET DISK DATA STORE IN BUFFER
324	302 305	340		JNZ	DREAD	SECTOR IT NO MAKE
327	054			INR	i.	ADVANCE POINTER GET DISK DATA STORE IN BUFFER ADVANCE POINTER A TEST IF DONE BACK UP POINTER SAVE IN D-E PAIR CALCULATE CRC
330	032		RL00P	LOAX	D	GET DISK DATA
			NE OU	MOV	M 4	STORE IN BUFFER
331	167			1101	11110	ADVANCE DOINTER
332	054			INR	L	ADVANCE FOUNTER
333	302 330	340		JNZ	RLOOP	# 1521 15 DOME
336	054 302 330 055			DCR	L	BACK UP POINTER
337	363			XCHG		SAVE IN D-E PAIR
	041 175	242		1 1 1	H DATAMARK	CALCULATE CRC & TEST FOR ZERO
340	041 1/3	342		CALL	CBECH	CALCULATE COC
343	315 124	341		CALL	CRE CIT	# TECT
346	171			MOV	A,C	a (E3)
347	260			ORA	В	FUR ZERO
350	076 001			MV [Ă,1	
352	300			RN7		RETURN IF CRC ERROR
	300	240		FXI	D.DATAMARK+1	
353	021 1/6	342			DMAADDR	
356	052 164	342		LHLD	OFMADUA	GET READY FOR DMA XFER
361	076 001 300 021 176 052 164 353 033			XCHG	_	GET MENOT FOR DEM WITH
362	033			DCX	D	
363	006 200		XFER	WA1	B.200Q	
365	176			MOV	A,M	GET DISK DATA
366	000			18/9	0	
	023			STAX	Ď	STORE IN MEMORY
367	022			21///	-1	
370	043			INX	H	DECREMENT BYTE COUNT & TEST FOR DONE
371	005			DCR	В	DECREMENT BYTE COUNT
372	302 365	340		JNZ	XFER+2	& TEST FOR DONE
375	170			MOV	A.B	
376	311			RET		
3/0	311			142		
			DEFET	CALL	LOHEAD	LOAD HEAD &
377		341	RSECT	CALL		
341:002	300			RNZ		TEST FOR READY
003		342		LXJ	H,BUFFER	TEST FOR READY SECTOR HEADER MARK SECTOR HEADER COUNT ADJUST TO READ DISK DATA READ DISK DATA SET FLAGS TEST FOR ZERO READ A MARK COMPARE WITH HEADER READ REXT HEADER IF ERROR ADVANCE BUFFER POINTER
006	006 007	, -		MVI	8.7	SECTOR HEADER COUNT
010	033			DCX	n`	AOJUST TO READ
	033				ñ	DISK DATA
011			ZWATT	DEA	0	OEAD DISK DATA
012			ZWALI	LDAX	Ų	READ DISK DAIN
013	247			ANA	Α	SET FLAGS
014	302 017 072 003	341		JNZ	ZWAIT	TEST FOR ZERU
017	072 00	343		LDA	RE ADMARK	READ A MARK
022		• • • •		CMP	M	COMPARE WITH HEADER
023		2 240		JNZ	RSFCT	READ NEXT HEADER IF ERROR
		7 340	TSL00P	INX	H	ADVANCE BUFFER POINTER
028	043		12000	INA		PERSONAL PRINCE CONTRA
027	005			DCR	В	DECREMENT HEADER COUNT
030	310			RZ		RETURN IF DONE
031	032			LDAX	D M TSLOOP	READ DISK DATA
032	276			CMP	M	COMPARE WITH MEMORY
	212 02	6 241		ĴΖ	TSL00P	READ MORE IF NO ERROR
033) J12 U2	0 341		MUT	12001	TEST FOR
036	0/6 00	5		MVI	A.5	
040	220			208	8	WRONG TRACK
04	370		,3200	RM		NUMBER
043	2 303 37	7 340		JMP	RSECT	NOT A TRACK ERROR
044						

E020 E030 E032 E035 E038 E038 E03C E030	CD AA E1 OE O1 CO 27 E1 CO C5 E0 C2 20 E0 C9 76	BLOOP	CALL MVI CALL CALL JNZ RET HLT HLT	HOME C.1 SETSEC DREAD BLGOP
E03E	0002		DS	2
E040 E042 E045 E048 E049	06 80 2A 70 E2 11 02 E3 1A 1F	INPUT WAIT	MVI LHLD LXI LDAX RAR RAR	B,200Q SCON D,STATUS D
E04A E04B E04E E051 E052 E053	1F DA 48 EO CO 95 EO 1A 1F 1F		JC CALL LDAX RAR RAR	WAIT XLOOP D
E054 E057 E05A E058 E05C	DA 40 E0 CD 81 E0 1A 1F 1F	DATAL	JC CALL LDAX RAR RAR	JMPUT SOELAY D
E050 E05E E05F E060 E063	78 1F 47 02 57 E0 CD 81 E0		MOV RAR MOV JNC CALL	B,A DATAL SDELAY
E066 E069 E06A €06C	CD 81 E0 78 E6 7F C9		CALL MOV ANI RET	SDELAY A,B 177Q
E060 E070	11 03 E3 87	OUTPUT	ADO	D,SERIAL A
E071 E073 E074	. OÉ OB 37 1F	OLOOP	MVI STC RAR	C,11
E075 E076 E077	47 9F 12		MOV SBB Stax	8.A A D
E078 E078 E07C	CD 81 EO 78 OD		CALL MOV DCR	SDELAY A,B C
E070 E080	C2 73 E0 C9		JN Z RET	OL002

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E081	2A 70 E2	SDELAY	LHLD	SCON
E084	29		DAD	H
E085	28	XŁOOP	DCX	н
£086	7C .		MOY	A,H
E087	B5		ORA	L
E088	C2 85 E0		JNZ	XLOOP
£08B	Č9		RE T	
	•			
E08C	CD ED E1	DWRITE	CALL	LDHEAD
E08F	IA .		LOAX	D
E090	E6 01		ĀNI	1
E092	CO		RNZ	_
E093	2A 74 E2		LHLD	ONAA DDR
£096	11 70 E2		LXI	O DATAMARK
			PÜSH	D
E099	05		CALL	XFER
A603	CD F3 E0			
E090	E1		POP	H Crech
E09E	CD 54 E1		CALL	
EOA1	71		MOV	M,C
EOAZ	23		INX	H
E0A3	70		MOV	M,B
EOA4	OE 07		MVI	C.7
EGA6	CD FF EO		CALL	RSECT
EOA9	CO		RNZ	
EDAA	06 OA		MAE	B,10
EOAC	1A	FLOOP	LDAX	D.
EOAO	05		DCR	8
EOAE	C2 AC EO		JNZ	FLOOP .
E081	AF		XRA	A
EOB2	12	ZLOOP	STAX	Ð
	√ QD		DCR	C
E084	C2 82 E0		JNZ	ZLOOP
E087	13		INX	D .
EOBB	žĚ		MOV	A,A
E CB9	12		STAX	Đ
EOBA	18		DCX	Ď
EOBB	23		INX	Ķ
EOBC	7E	WLOOP	MOV	Ä,M
E080	12	WL COP	STAX	D
EOBE	2C		INR	Ĺ
EOBF	C2 BC EO		JNZ	WLOOP
			XRA	A
EOC2	AF		STAX	Ô
EDC3	12		RET	U
EOC4	C9		RE I	
EOC5	CD FF EO	DREAD	CALL	RSECT
EOCS	CO	DITTE	RNZ	PARGI
EOC9	06 DE		MVI	B,14
	IA	RWALT	LDAX	0
EOCB		RMALI	OCR	B
EOCC	05			
EOCD	C2 CB EO		JNZ	RWAIT
EODO	3A 01 E3		LDA	READMARK
EOD3	· BE		CMP	M
EQD4	C2 C5 E0		JNZ	DREAD

341:276 277 301 304 306 310 311 311 311	346 010 312 271 341 257 6 041 166 342 0 066 376 2 054 3 312 045 341 6 167	1L00P	LDAX ANI JZ XRA LXI MVI INR JZ MOV JMP	D TZERO STEPO A H,BUFFER M,OFEH L SETSEC-2 M,A 1LOOP	GET DISK STATUS & TEST FOR TRACK ZERO SET UP DATA BUFFER FOR TRACK ZERO & SECTOR ZERO
32: 32: 33: 33: 33: 33: 33: 34: 34: 35: 35: 35: 35: 35: 35: 36: 37: 37: 38: 38: 38: 38: 38: 38: 38: 38: 38: 38	5 365 5 346 353 0 250 1 022 2 356 020 4 022 5 361 6 022 7 006 010 1 076 143 3 177 4 075 5 302 343 341 0 005 1 302 341 341	TDELAY DELAY	LDA PUSH ANI XRA STAX STAX POP STAX MVI MOV DCR JNZ RET	DRIVE PSW 353Q B D STEP D STEP D B,TCONST A,MSEC A,A A CELAY+2 B DELAY	MERGE DIRECTION & STEP SEND TO DRIVE FINISH STEP PULSE SEND TO DRIVE HEAD MOTION CONSTANT DELAY ONE MILLESECOND TEST FOR DELAY ONE
37	0 032 11 346 004 3 072 003 343 6 006 043 10 314 341 341 13 032 14 27 15 237 16 074	LDHEAD	LXI LDAX ANI LDA MYI CZ LDAX RAL SBB INR RET	D,STATUS D 4 LOADHEAD B,SETTLE DELAY D A A	DISK STATUS MEM LOC GET STATUS BYTE STRIP OFF HEAD LOAD BIT LOAD HEAD & SELECT ORIVE HEAD LOAD SETTLE TIME CONDITIONALLY WAIT FOR SETTLE GET STATUS SET ACC TO 3770 OR ZERO THE ACC FOR READY FLAG TO SYSTEM

Disk Jockey Shugart Firmware

	E000 E300 E301 E302 E303	READMARK STATUS E	EQU ORI EQU READI QU READI	
	E300 E301 E302 E303 E276 E279 E274 E27D E2FD E273 E270	SERIAL E BUFFER E SECREG E DMAADDR DATAMARK	K EQU RI EQU STA QU LOADA QU ORIGE QU BUFFI EQU BUFI EQU BUFI EQU BUI U BUFFEI	EADMARK ATUS HEAD IN+2:166Q ER+3 FER-2 FFER+7 TAMARK+200Q R-3
	0008 0014 0010 0010 0008 0063 0023	TZERO MOVIN MVOUT STEP TCONST MSEC SETTLE	EQU EQU EQU EQU EQU EQU	10Q 24Q 20Q 20Q 10Q 143Q 35
E003 E006 E009 E00C E00F E012 E015	C3 18 E0 C3 40 E0 C3 6D E0 C3 AA E1 C3 84 E1 C3 27 E1 C3 3F E1 C3 C5 E0 C3 8C E0	DBOOT TERMIN TRMOUT TKZERO TRKSET SECTOR DMA READ WRITE	AMC AMC AMC AMC AMC AMC AMC AMC	800T INPUT OUTPUT HOME SEEK SETSEC SETDMA OREAD DWRITE
E01B E01E E021 E023 E026 E027 E028 E029 E029	21 21 00 3E 88 01 80 00 C5 F5 E5	воот	LXI LXI MVI EXI PUSH PUSH PUSH STA	SP,BUFFER H,41Q A,210Q B,200Q B PSW H B STATUS

16A	4F		MOV	C.A
TON				
168	7A		MOV	A,D
160	E6 E0		AN I	0E0H
16E	AB		XRA	Ε
16F	47		MOY	B,A
170 171	7A		MOV	A,D
171	OF		RRC	
172	E6 F0		ANI	O FOH
174	A9		XRA	C
175	4F		MOV	Č,A
175 176	23		INX	H
177	01		POP	Ď
177 178	7A		MOV	Ă,D
179	ВC		CMP	H
178	DB		RC	11
17A 17B				CDCCU13
178	C2 57 E1		JNZ	CRECH+3
1 7E 17F	7B		MOV	A,E
17F	BD		CMP	Ŀ
180 181	D8		RC .	
101	C3 57 E1		JMP	CRECH+3
184	CD ED E1	SEEK	ÇALL	LDHEAD
187	CO		RNZ	
188	79		MOV	A.C
189	E6 7F		AN!	1770
189 188	4F		MOV	C,A
18C	Č5 B3		ADI	179
18E	9F		588	Ā
18F	CO		RNZ	,,
100	21 77 E2		LXI	H,BUFFER+1
190 193	7E	SLOOP	MOV	A.M
194		SLOOP	CMP	
195	89 54 BC C1			C
130	CA DF E1		JZ	DELAY-2
198 19B	DA A1 E1		JC	MOVEIN
TAR	06 10		MVI	B_MVOUT
190 19E	35		DCR	M
	C3 A4 E1		JMP	MOVE IN+3
141	05 14	MOVE IN	MVI	B,MOVIN
1A3	34		1NR	М
1A3 1A4	CD D2 E1		CALL	TOELAY
1A7	C3 93 E1		JMP	SLOOP
1AA	CD ED E1	HOME	ÇALL	LOHEAD
140	CO		RNZ	
1AE 180	06 14	HLOOP	IVM	B,MOVIN
180	CD D2 E1		CALL	TOELAY
183	1A		LDAX	D
184	£6 08		ANI	TZERO
186	C2 AE E1		JNZ	HLOOP
186 189	06 10	STEPO	MVI	B MYOUT
103	00 10	SILFU	U.T.I.	ו טטצייו, פ

E188 E18E E18F E1C1 E1C4 E1C5 E1C8 E1CA E1CA E1CCF	CD D2 E1 1A E6 O8 CA B9 E1 AF 21 76 E2 36 FE 2C CA 25 E1 77 C3 CA E1	ILOOP	CALL LOAX ANI JZ XRA LXI MVI INR JZ MOV JMP	TDELAY D TZERO STEPO A H,BUFFER M,OFEH L SETSEC-2 M,A ILOOP
£102 £105 £106 £108 £109 £10A £10C £10C	3A 73 E2 F5 E6 EB AB 12 EE 10 12 F1	TDELAY	LOA PUSH ANI XRA STAX XRI STAX POP STAX	DRIVE PSW 353Q B D STEP D PSW D
E10F E1E1 E1E3 E1E4 E1E5 E1E8 E1E9 E1EC	06 08 3E 63 7F 3D C2 E3 E1 05 C2 E1 E1	DELAY	MVI MVI MOV DCR JNZ DCR JNZ RET	B,TCONST A,MSEC A,A A DELAY+2 B DELAY
E1ED E1F0 E1F1 E1F3 E1F6 E1F8 E1FC E1FC E1FC E1FF	11 02 E3 1A E6 04 3A 03 E3 06 23 CC E1 E1 1A 17 9F 3C C9	L DHE AD	LXI LDAX ANI LDA MVI CZ LDAX RAL SBB INR RET	D,STATUS O 4 LDADHEAD B,SETTLE OFLAY D A A

EOD7	20		INR	Ļ
EOD8	1A	RL00P	LDAX	D
E009	77		MOY	M,A
EGOA	2C		[NR	L
EOOB	C2 D8 E0		JNZ	RLOOP
EODE	20		DCR	L
EOOF	EB		XCHG	
EOEO	21 70 E2		LXI	H, DATAMARK
E0E3	CD 54 E1		CALL	CRECH
EOE6	79		MOA	A,C
EOE 7	80		ORA	В
EQE8	3E 01		MVI	A.1
EOEA	Ç0		RNZ	D CATAMARY 1
£0EB	11 7E E2		LXI	D.DATAMARK+1
EOEE	2A 74 E2		LHLD	DMAAODR
EOF1	EB		XCHG DCX	D
EOF2	18	vern	MVI	B.200Q
EQF3	06 80	XFER	WOA	A.M
EOF5	7E 13		INX	ָרָה . ה
E0F6	13 12		STAX	Ď
EOF7 EOF8	23		INX	Ĥ
EOF9	05		DCR	В
EOFA	C2 F5 EQ		JNZ	XFER+2
EOFD	78		MOV	A,B
EOFE	Ć9		RET	014
COLE	63			
EOFF	CD ED E1	RSECT	CALL	LDREAD
E102	CO		RNZ	
E103	21 76 E2		LXI	H,BUFFER
£106	06 07		MVΙ	8,7
E108	18		DCX	D
E 109	18		DCX	D
ElOA	1A	ZWAIT	LDAX	0
£108	A7		ANA	A
E 10C	C2 OA E1		JNZ	ZWAIT
€10F	3A Q1 E3		LDA	READMARK
E112	BE		CMP	M
E113	C2 FF EO		JNZ	RSECT
E116	23	†\$LOOP	INX	H .
E117	05		DCR	В
£118			RZ	•
E119	1A		LDAX	D
EllA	BE		CMP	M TSLDOP
£118	CA 16 E1		JZ	
E11E	3E 05		MVI Sub	A,5 B
E 120	90		RM SUB	Đ
E121	FB		JMP	RSECT
E122	C3 FF EO		QP#	710601

125	OE 01		MVI	C.1
127	79	SETSEC	MOV	A,C
128	E6 1F		AN I	37Q
12A	Ç8		RZ	
12B	4F		MOV	C.A
120	C6 E\$		ADI	345Q
12C 12E 12F 130	9F		\$88	A
12F	CO		RNZ	
130	21 79 E2		FXI	H,SECREG
E133 E134	71		MOV	M,C
E 1 34	CD 4E E1		ÇALL	SETCRC
E137 E138	71		MOV	M,C
	23		INX	H
E 139 E 13A	70		MOV [NX	M,B H
E 13A	23		MVI	M.OFBH
E 13B	36 FB		XRA	A
E 130	AF		RET	^
E 13E	C9		RE 1	
E13F	21 00 E2	SETOMA	LXI	H,ORIGIN+2:000Q
E142	7C	JC TOPEN	MOV	A,H
E143	88		CMP	8
E144	C8		RZ	•
E 145	30		INR	A
E 146	88		CMP	B
E 147	Č8		RZ	-
E 148	60		MOV	H.B
E149	69		MOV	L,C
E 14A	22 74 E2		SHLO	OMAADDR
É14D	C9		RET	
E 14E	21 76 E2	SETCRC	LXI	H.BUFFER
E151	11 7A E2	3210.10	LXI	D.SECREG+1
E154	01 FF FF	CRECH	LXI	B,-1
E157	05	ONLOW	PUSH	Ď, -
£158	7E		MOV	A.M
£159	Á9		XRA	C
₹15A	57		MOV	D.A
E15B	ÓF		RRC	- •
E15C	ÕF		RRC	
£150	OF		RRÇ	
E15E	0F		RRC	
E15F	E6 OF		AN1	QFH H4Q
E161	AA		XRA	٥
£162	5F		MOV	E ,A
E 163	OF		RRC	
E164	OF		RRC	
£165	OF		RRC	
E 166	57		MOV	D.A
E 167	E6 1F		ANI	1FH
E 169	A8		XRA	8

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Acknowledgment

Like many other Thinker Toy products designed by Morrow, the Disk Jockey I controller utilizes the power of the CPU to accomplish many of its I/O tasks. However, unlike other Thinker Toy products, the key idea which makes the interface work is not the "brain child" of George Morrow. The Disk Jockey controller uses the CPU's "READY" line in a novel fashion. This unusual use of READY first appeared in an article by Eugene Fisher in the November 8, 1975 issue of <u>Electronics Design Magazine</u>. Mr. Fisher's penetrating insight into the effective use of the microcomputer to control a floppy disk has served as an inspiration in the design of the Disk Jockey controller.

Warranty

Parts are warranted to be free from defects in material and work-manship. Parts for the Disk Jockey I purchased in kit form is warranted for ninty days from invoice/purchase date. The Disk Jockey I purchased as an assembled unit or as part of the DISCUS I system is warranted for six months from invoice/purchase date. Any board purchased in kit form which is returned for testing/repair is subject to a fee of up to \$35.00. Any out-of-warranty repair of up to \$35.00 will be made without prior approval of customer.

Parts and labor warranty for the disk drive is for forty-five days from invoice/purchase date. For a period of up to one year, there is a flat \$55.00 labor charge for warranty parts replacement. After one year, charges will be made for parts and labor.

Warranty is void if in the opinion of Morrow/Thinker Toys the unit has been subject to abuse, misuse, improper assembly, or if directions have not been followed in assembly.

A COPY OF THE INVOICE OR PROOF OF PURCHASE IS REQUIRED FOR IN-WARRANTY SERVICE. A description of the problem must accompany any item returned for repair. Shipments must be made to Thinker Toys prepaid. Morrow/Thinker Toys is not responsible for any consequential damage.

The foregoing warranty is in lieu of all other warranties expressed or implied and in any event is limited to product repair or replacement.

Morrow Designs, Inc.

Thinker Toys ** 5221 Central Avenue, Richmond, CA 94804 (415) 524-2101

LIMITED WARRANTY

Morrow Designs Inc. warrants its products to be free from defects in workmanship and material for the period indicated. This warranty is limited to the repair or replacement of parts only and liability is limited to the purchase price of the product. The warranty is void if, in the sole opinion of Morrow Designs Inc., the product has been subject to abuse, misuse, unauthorized modification, improper assembly, non-conformance to assembly directions, or if the unit is used in any other manner than intended.

KITS - Parts, including the printed circuit boards, purchased in kit form are warranted for a period of ninety (90) days from the invoice/purchase date. If a board, which was purchased in kit form, is returned for testing or repair, a minimum service charge of \$35. will be assessed.

ASSEMBLED BOARDS - Parts, including the printed circuit boards, purchased as factory assembles, are warranted for a period of six (6) months from the invoice/purchase date. Out-of-Warranty boards returned for testing or repair will be assessed a minimum of \$35. service charge. If the charge to repair will exceed \$35., the customer will be notified prior to the actual repair.

ELECTROMECHANICAL PERIPHERALS - Peripheral equipment, such as floppy disk drives, hard disk drives, etc., not manufactured by Morrow Designs Inc. have warranties which vary according to the manufacturer. In most cases, Morrow Designs Inc. provides a warranty equal to or greater than the original manufacturer. Please contact the factory for individual warranty information. Warranty information for each device is included with the equipment when it is shipped.

RETURN PROCEDURE - A COPY OF THE INVOICE OR PROOF OF ORIGINAL PURCHASE IS REQUIRED AND MUST ACCOMPANY THE ITEM FOR IN-WARRANTY SERVICE. Items returned without proof of original purchase will be sent back, shipping charges collect. A description of the problem must accompany the returned item. Shipment must be made prepaid to Morrow Designs Inc. Repaired items will be shipped via U.P.S. surface. Shipment by air requires payment of the additional charges. Morrow Designs Inc. is not responsible for any consequential damages or for damage incurred in transit.

The foregoing warranty is in lieu of all other warranties either expressed or implied and, in any event, is limited to product repair or replacement.

Effective February 1, 1980

Specifications, terms, and pricing are subject to change without notice.

LIMITED WARRANTY

DISCUS 1 and DISCUS 2D Systems

This addendum to Morrow Designs Inc. Limited Warranty applies to the Shugart Associates Model 800/801 Floppy Disk Drives as used in the DISCUS 1 and 2D Disk systems.

Parts and labor for a floppy disk drive purchased from Morrow Designs Inc. are warranted for a period of forty-five (45) days from the invoice/purchase date. For a period of one (1) year from the invoice/purchase date, parts are warranted. A fixed fee of \$55. will be charged for labor. After one (1) year current rates for parts and labor will be charged.

